

17 Sep 2025

Finishing Duality + Simplex Alg

Announcement: Prelim covers through 9/19 lecture.

$$\begin{array}{ll}
 \text{(P)} & \text{(D)} \\
 \max C^T x & = \min b^T y \\
 \text{s.t. } Ax \preceq b & \text{s.t. } A^T y \preceq c \\
 x \succeq 0 & y \succeq 0
 \end{array}$$

Equivalent "equality form" of problem (P).

$$\begin{array}{ll}
 \max C^T x & x \in \mathbb{R}^n \quad w \in \mathbb{R}^m \\
 \text{s.t. } Ax + w = b \\
 x, w \succeq 0
 \end{array}$$

If the set of (x, w) where the max is attained is non-empty, it contains a point where n of the variables $\{x_1, \dots, x_n, w_1, \dots, w_m\}$ are zero.

Of the points x, w where max is attained \exists one with as many 0 words as possible. (x^*, w^*)

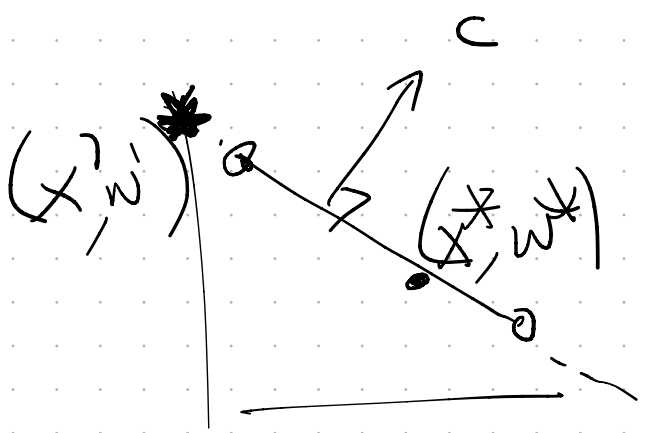
Suppose $n-1$ or fewer words are zero.
 Augment $Ax + w = b$ with $n-1$ or fewer

equations of the form $x_i = 0$
 or $w_j = 0$ corresponding to
 the coords that are zero at
 the ostensibly opt point.

The remaining coords of (x^*, w^*)
strictly positive.

$\therefore (x^*, w^*)$ belongs to interval of a
 line segment in $\mathbb{R}_{\geq 0}^{m+n}$ satisfying
 all $m+n-1$ (or fewer) of these.

$c^T x$ must be const along this segment.



Extend segment to an
 "exit point" (x', w') where
 it exits the orthant
 $\mathbb{R}_{\geq 0}^n$.

$$c^T x' = c^T x^* \quad (x', w') \in \mathbb{R}_{\geq 0}^{n+m} \quad Ax' + w' = b$$

\therefore and x' has at least one more
 coordinate equal to zero. \swarrow

If $S \subset [m+n]$ denotes the set of n coordinates of (x^*, w^*) that are $\neq 0$,

then
$$\mathbb{I}_S = \begin{matrix} n \\ \begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 \end{bmatrix} \end{matrix}$$
 e.g. $S = \{4, 3, 5, 6\}$.

("coordinate projection matrix")

$$Ax^* + w^* = b$$

The S coord's of (x^*, w^*) equal zero

$$\underbrace{\begin{bmatrix} A & \mathbb{I}^{m \times m} \\ \hline & \mathbb{I}_S \end{bmatrix}}_{M_S} \begin{bmatrix} x^* \\ w^* \end{bmatrix} = \begin{bmatrix} b \\ 0 \end{bmatrix}$$

We can find S s.t. $M_S \begin{bmatrix} x^* \\ w^* \end{bmatrix} = \begin{bmatrix} b \\ 0 \end{bmatrix}$

and M_S is invertible.

~~Then~~
$$A \begin{bmatrix} x^* \\ w^* \end{bmatrix} = M_S^{-1} \begin{bmatrix} b \\ 0 \end{bmatrix}$$

(Note this reduces LP to searching over $S \subset [n+m]$.)

Every other feasible (x, w) is of the form $M_S^{-1} \begin{bmatrix} b \\ v \end{bmatrix}$ for some $v \geq 0$.

Reason: if (x, w) is feasible,

$$M_S \begin{bmatrix} x \\ w \end{bmatrix} = \begin{bmatrix} Ax + w \\ \mathbb{1}_S(x) \end{bmatrix} = \begin{bmatrix} b \\ \geq 0 \end{bmatrix}$$

$$\begin{bmatrix} x \\ w \end{bmatrix} = M_S^{-1} \begin{bmatrix} b \\ v \end{bmatrix} \text{ for } v \geq 0.$$

The objective function $c^T x$
can be expressed as

$$\begin{bmatrix} c^T & 0^T \end{bmatrix} \begin{bmatrix} x \\ w \end{bmatrix} = \begin{bmatrix} c^T & 0^T \end{bmatrix} M_S^{-1} \begin{bmatrix} b \\ v \end{bmatrix}$$

As a function of v , the RHS
is affine (constant plus linear).

At $v = 0$ it evaluates to $c^T x^* = \text{OPT}$.

The RHS can be written as

$$\text{OPT} - z^T x - y^T w.$$

This equals $c^T x$ for all
 x, w satisfying $Ax + w = b$.

$$\forall x, w \quad Ax + w = b \Rightarrow C^T x = \text{OPT} - z^T x - y^T w$$

$$w = b - Ax$$

$\forall x$

$$C^T x = \text{OPT} - z^T x - y^T (b - Ax)$$

$$= \text{OPT} - z^T x - y^T b + y^T Ax$$

$$= (\text{OPT} - b^T y) + (y^T A - z^T) x$$

$$\therefore b^T y = \text{OPT}$$

$$c = A^T y - z$$

$$A^T y = c + z.$$

CLAIM. $y \geq 0$ && $z \leq 0$.

This is implied by local optimality of (x^*, w^*) .

(There's an issue with degeneracy ... see typeset notes.)