

27 Aug 2025

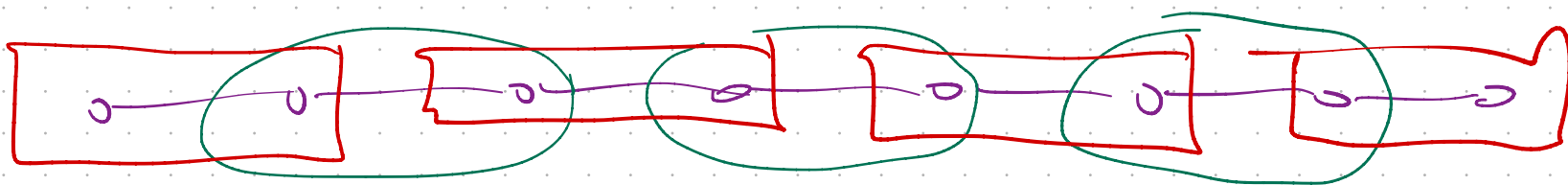
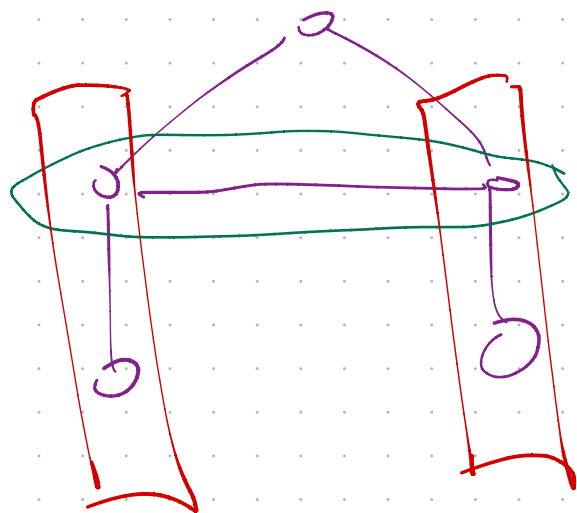
# Computing a Maximum Matching

Recap:

Matching = set of edges with no vertices in common

Free = not part of any edge in the matching

Matched = not free



Def. An augmenting path in  $G$  with respect to  $M$  (" $M$ -augmenting path") is a set of edges forming a (simple) path in  $G$ , starting & ending at free vertices, whose edges alternate btw belonging to  $M$  and not belonging.

Standard.

"path" means "simple path"  
(no repeated vertices)

"walk" means "non-simple path"  
(may repeat vertices)

Lemma. For a graph  $G$  and a matching  $M$ , the following are equivalent.

(A)  $M$  is a maximum matching in  $G$

(B)  $G$  contains no  $M$ -augmenting path.

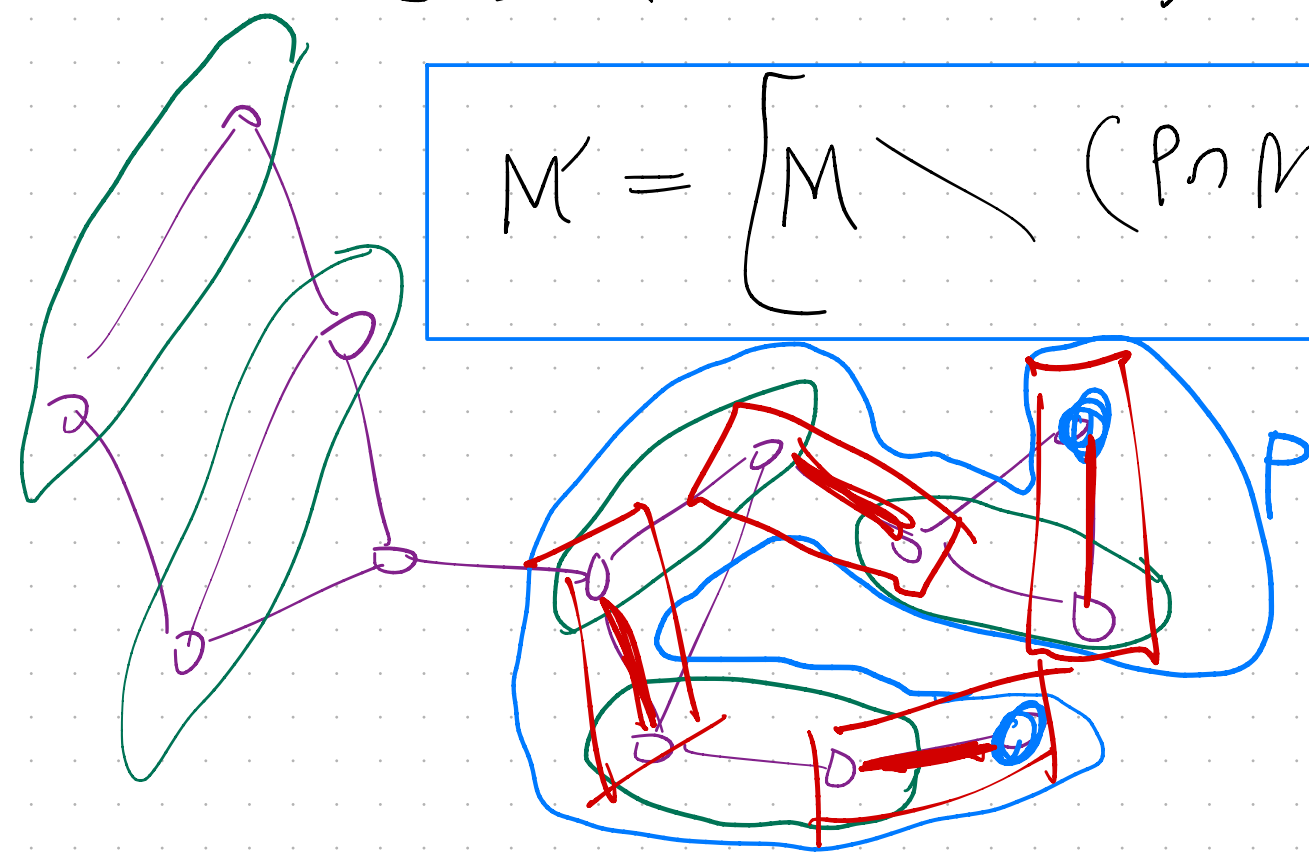
Proof  $(\neg B) \Rightarrow (\neg A)$

If  $P$  is a  $M$ -augmenting path construct new edge set  $M'$  by

$$M' = [M \setminus (P \cap M)] \cup (P \setminus M)$$

$$|M'| = |M| + 1.$$

$$M' = M \oplus P$$



$$(\neg A) \Rightarrow (\neg B)$$

"If  $\exists$  a matching  $M'$  with more edges than  $M$ , then  $\exists$  a  $M$ -augmenting path."

Look at  $M \oplus M'$ . It has maximum degree 2, so it is a union of disjoint paths and cycles.

The edges of each of these paths/cycles alternate between  $M$  and  $M'$ .

$M'$  has more edges than  $M$ .

$\Rightarrow$  At least one of these paths/cycles has more

$M'$  edges than  $M$  edges.

$\Rightarrow$  It must be a path  $P$  that starts and ends with  $M'$  edges.

This  $P$  is an  $M$ -augmenting path.

Endpoints of  $P$  belong to an edge of  $M \cap P$ . They can't belong to any other edge of  $M$ .

They also don't belong to any other edge of  $M \oplus M'$ . This implies they can't belong to any edge of  $M \implies$  they are free w.r.t.  $M$ .

Naive Iterative Max Matching

NIMM( $G$ ):

Start with  $M = \emptyset$ .

while  $\exists$  an  $M$ -aug path  $P$

$M \leftarrow M \oplus P$

endwhile

output  $M$ .

Lemma above  $\Rightarrow$  if this terminates  
its output is a max matching.

# while-loop iterations is  $\leq$   
# edges in  $G$

so termination is assured.

How fast can we test if  
 $\exists$  an  $M$ -augmenting path?

