

# **A Quick Introduction to How Blockchains Work**

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(some slides due to former postdoc  
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# Blockchain's Promise and Limits

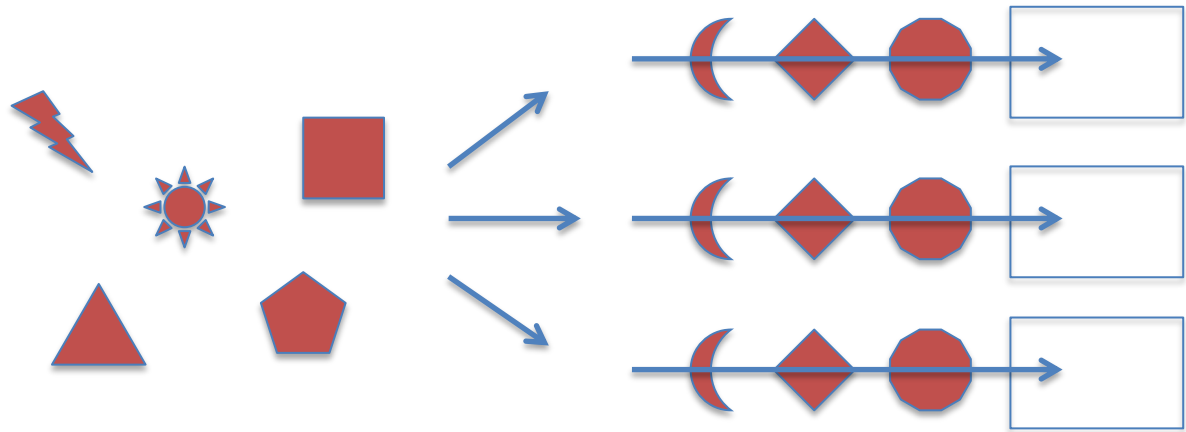
## Promises

- Global currency
- *Smart contracts*
- Notarization
- Accountability
- ...

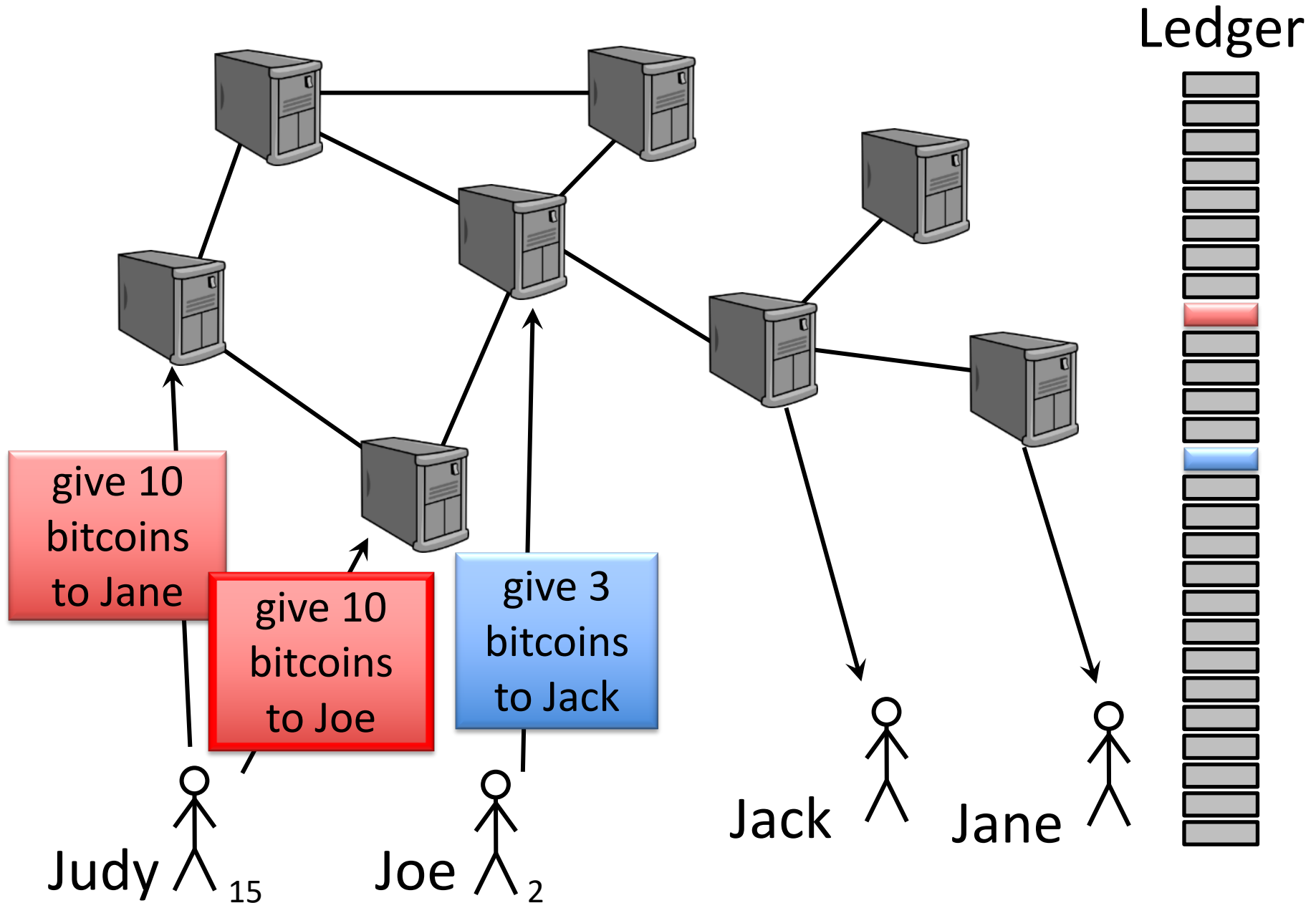


# State Machine Replication (Lamport'78)

- A generic way to **tolerate failures**
- Simply start multiple **replicas** (copies) of a state machine, and keep them in sync by **agreeing** on the transitions (operations) and the **order** in which to apply them

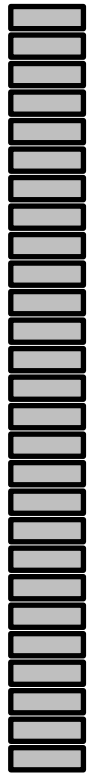


# A Replicated Ledger of Transactions

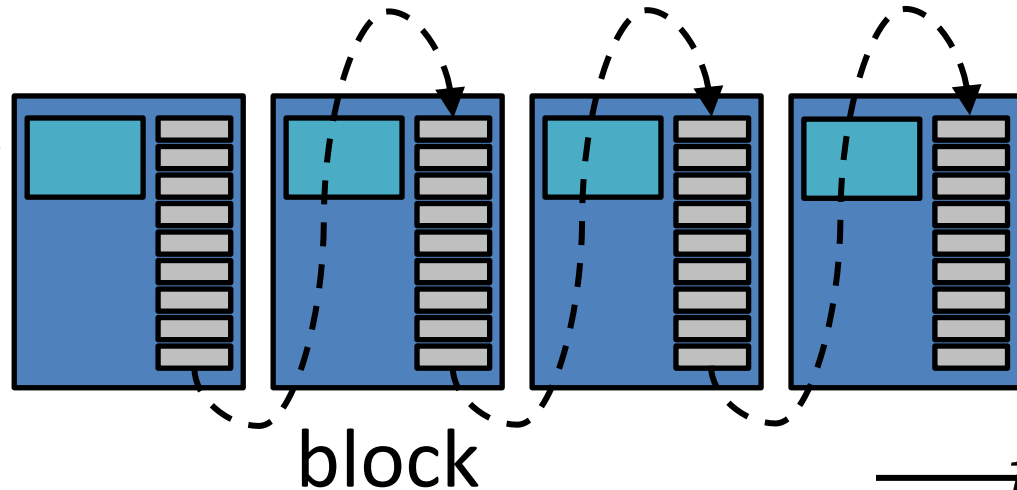


# The Blockchain

Ledger



header



$t$

*Purpose of blocks:*

***Batching transactions for efficiency!***

# Cryptographic One-Way Hash Function

$\text{hash}(X) = Y$

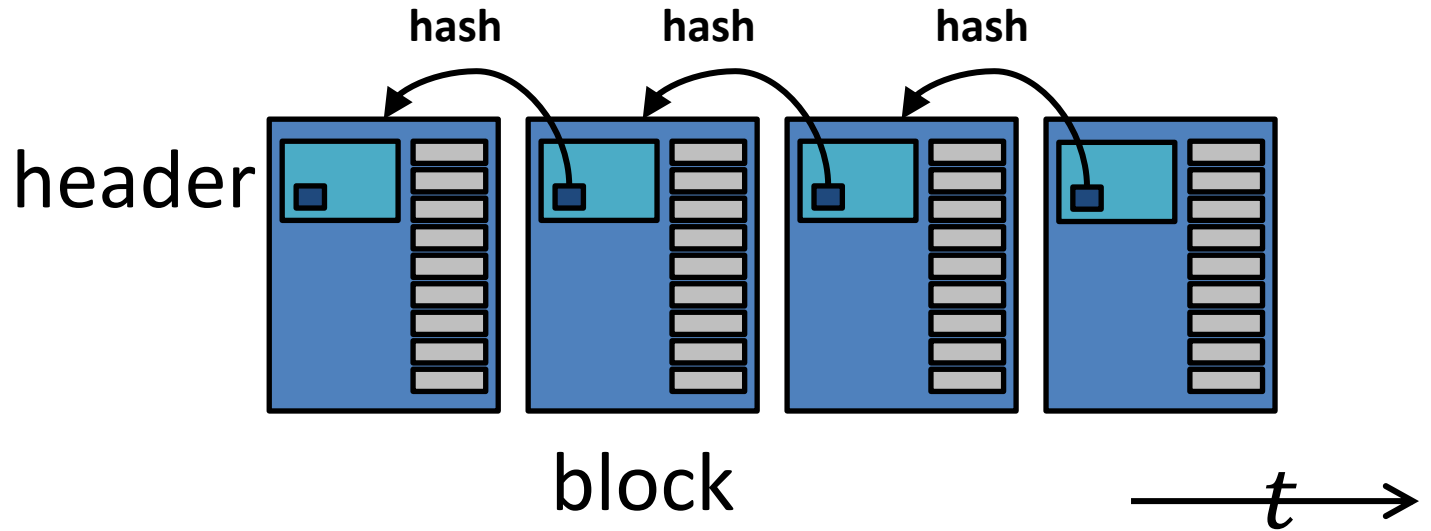
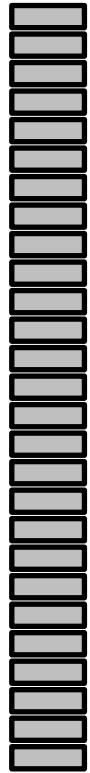
- Given  $X$  it is easy to compute  $Y$  (the *digest*)
- Given  $Y$  it is computationally infeasible to find  $X$ 
  - *unless you already know  $X$ , of course*
- It is computationally infeasible to find  $X_1$  and  $X_2$ ,  $X_1 \neq X_2$ , such that  $\text{hash}(X_1) = \text{hash}(X_2)$
- In some sense,  $Y$  *identifies*  $X$

Examples: SHA-256, SHA-3

*Note: unlike an ordinary hash function where you typically have fewer buckets than objects and thus multiple objects per bucket, with cryptographic hash functions you typically have many more “virtual buckets” than objects, and at most one object in a bucket*

# The Blockchain

Ledger



*each hash identifies the entire prefix of the ledger*

# Blockchain Desirables

- Performance:
  - High Throughput, Low Latency
  - Energy-Efficient
- Security:
  - Always available for reading and appending
  - Fair
  - Tamperproof (Integrity)
  - Possibly confidentiality as well
- No Single Administrative Domain
- Open membership (or not)



# Open Membership is Hard

- Traditional replication is based on voting
- Problem: “Sybil” or impersonation attacks
  - a participant may try to vote multiple times
  - with open membership, anybody can create identities and vote many times

# Permissionless vs Permissioned Blockchains

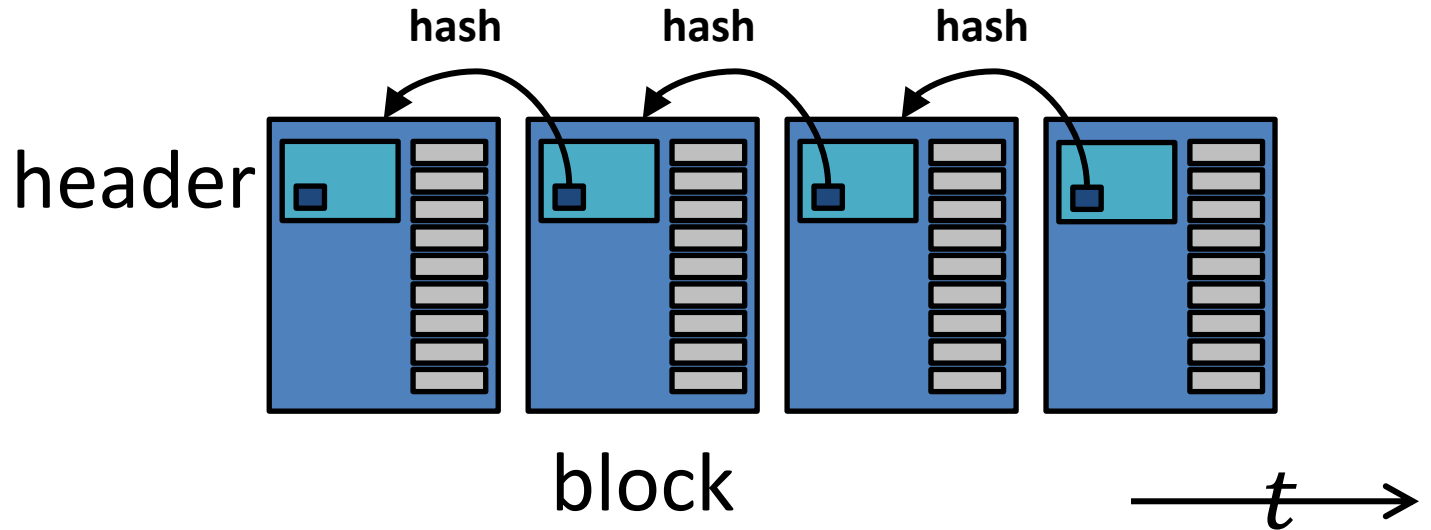
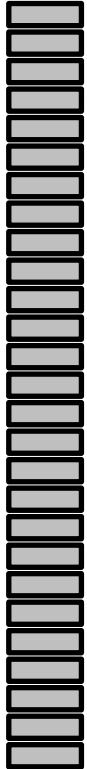
	Permissionless	Permissioned
Approach	Competitive	Cooperative
Basic technique	Proof-of-Resource	Byzantine Consensus
Trust requirements	Crypto (+ peers...)	Peers (+ crypto...)
Membership	Open	Closed
Energy-efficiency	Often terrible	Excellent
Transaction rate	At best hundreds / sec	Many thousands per second
Txn latency	As high as many minutes	Less than a second

# Bitcoin Blockchain

- Permissionless, open membership
- Proof-of-Work
- There are thousands of Bitcoin miners
  - often use cheap, dirty coal-based energy
  - they use ASIC hardware to compute SHA256 hashes
  - use about as much energy as an average European country
- Overall rate is a few transactions per second

# The Blockchain

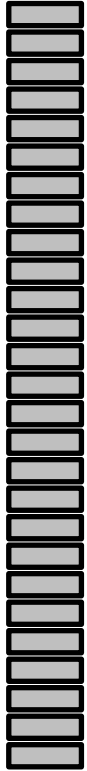
Ledger



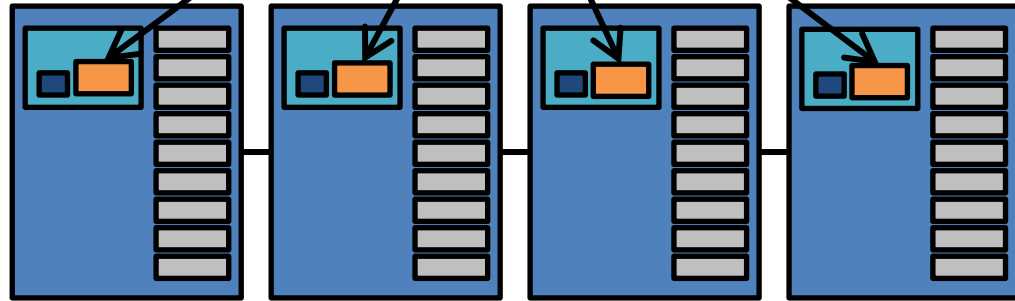
each hash identifies the entire prefix of the log

# The Blockchain

Ledger



nonce

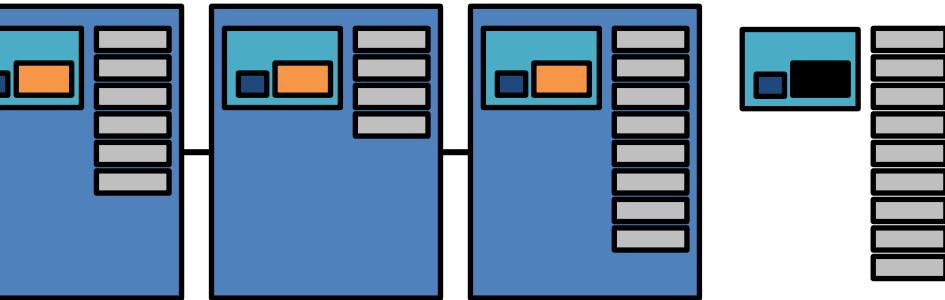


$\xrightarrow{t}$

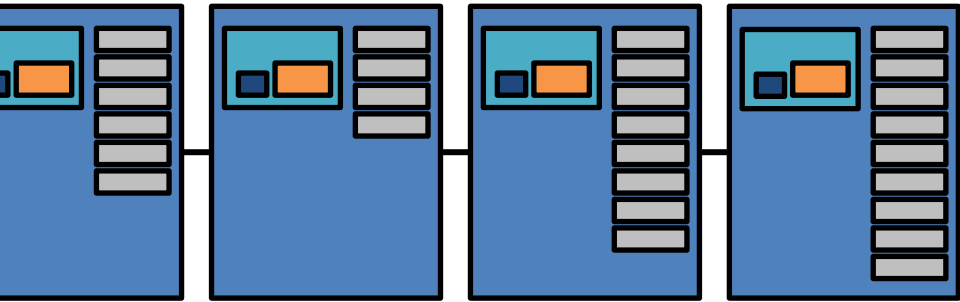
$$\text{HASH}(\text{block}) < \textit{target}$$

*“cryptopuzzle”*

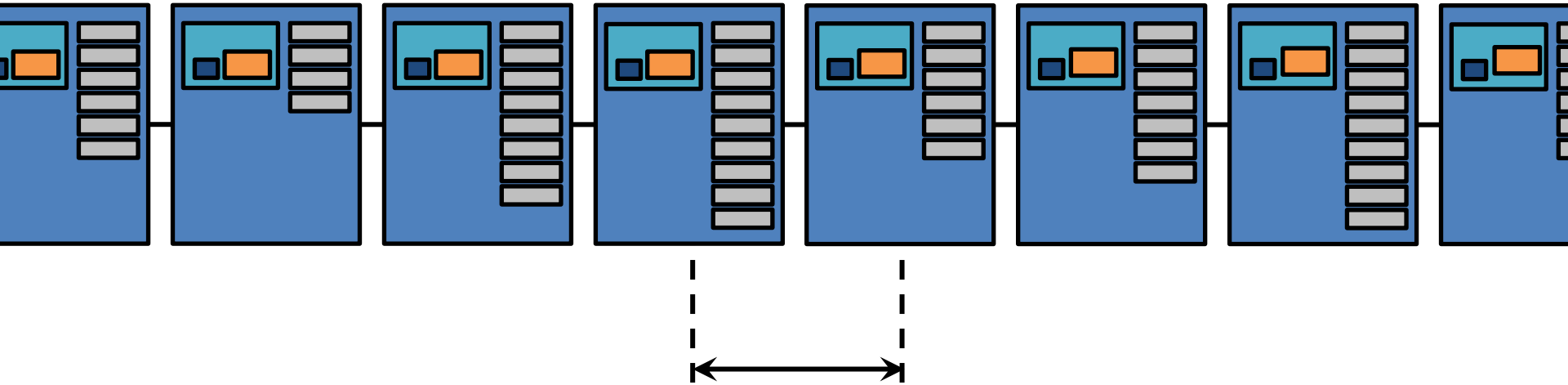
# The Blockchain



# The Blockchain



# The Blockchain



Exponentially distributed, with  
constant mean interval

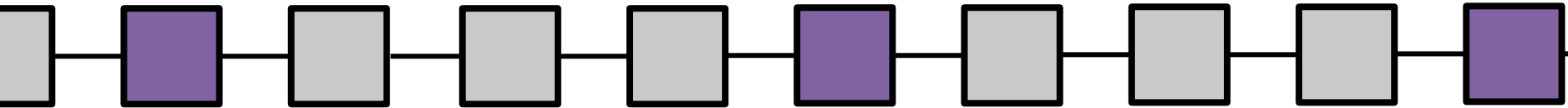
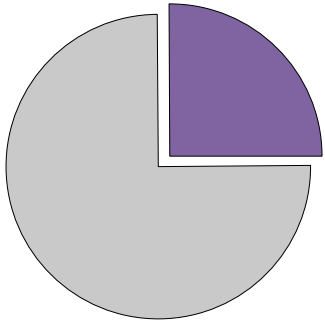
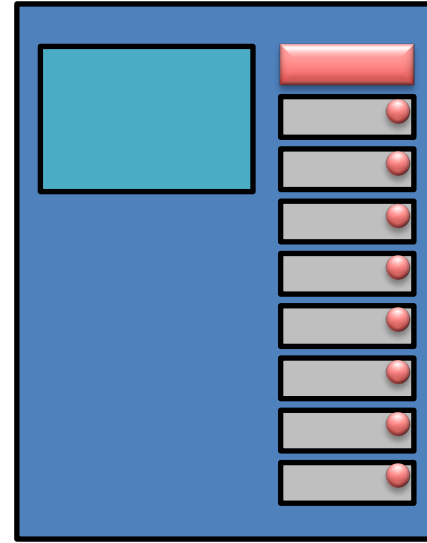
***target*** automatically adjusted every 2016  
blocks so that mean interval is **10 minutes**



# Incentives for Mining

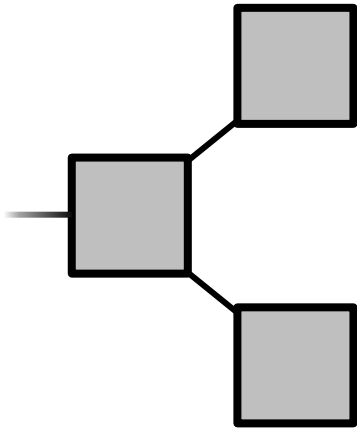
Prize:

- “Minting”
- Transaction Fees



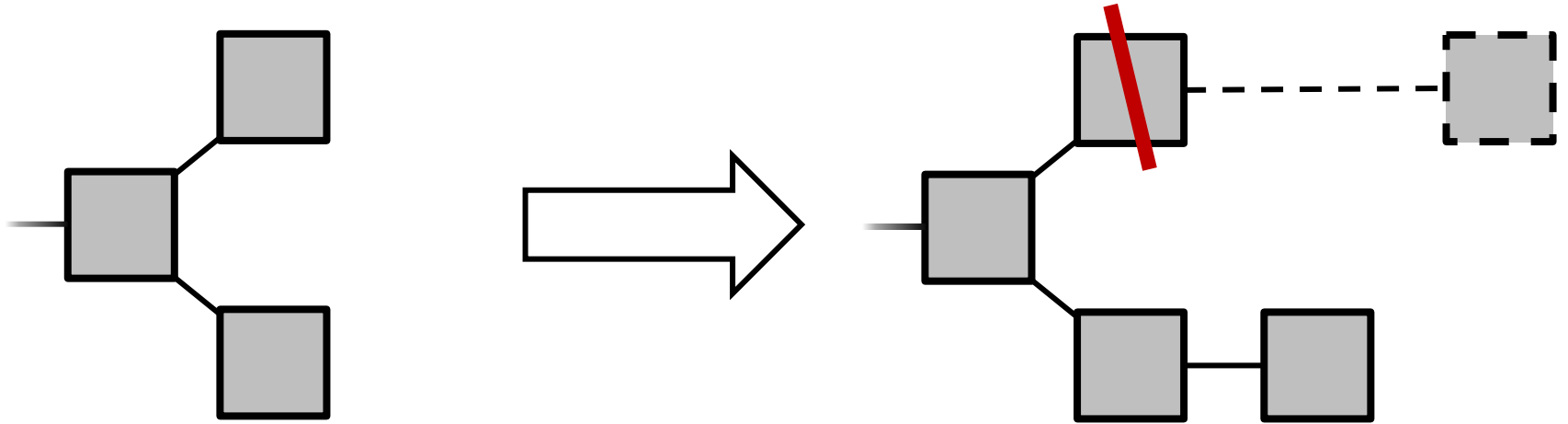
***Wins proportional to computation power***

# Forks



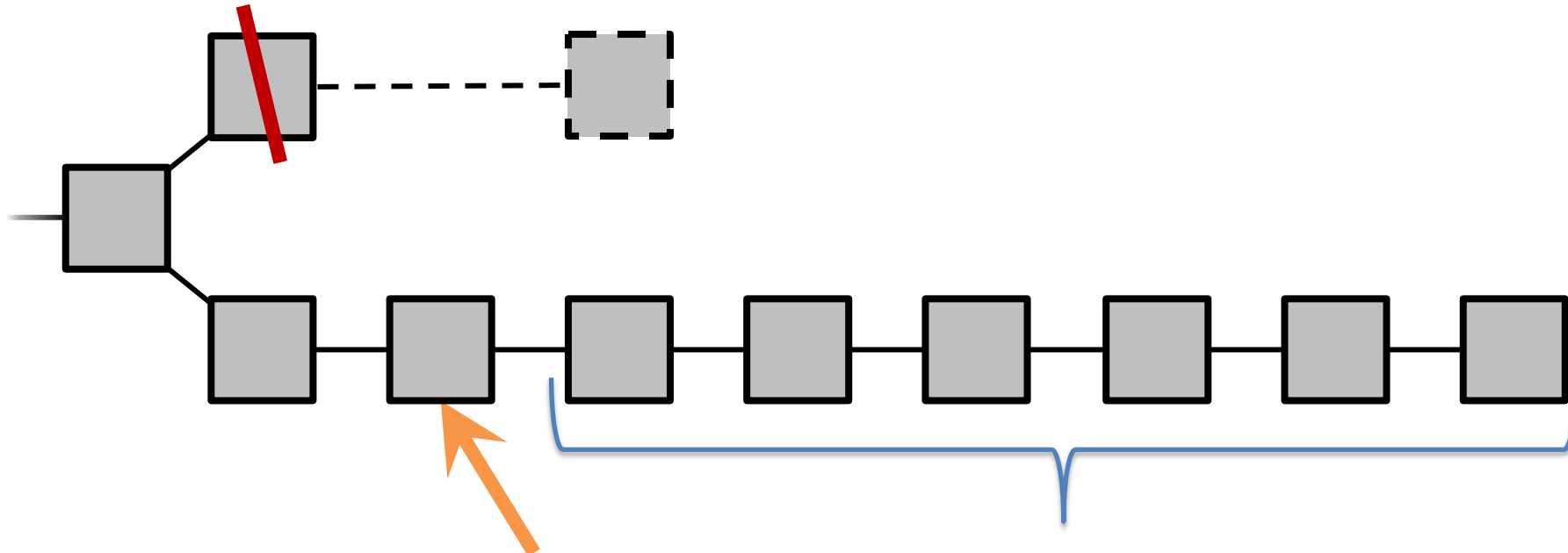
Two blocks “mined” at approximately the same time by two different miners

# Fork Resolution



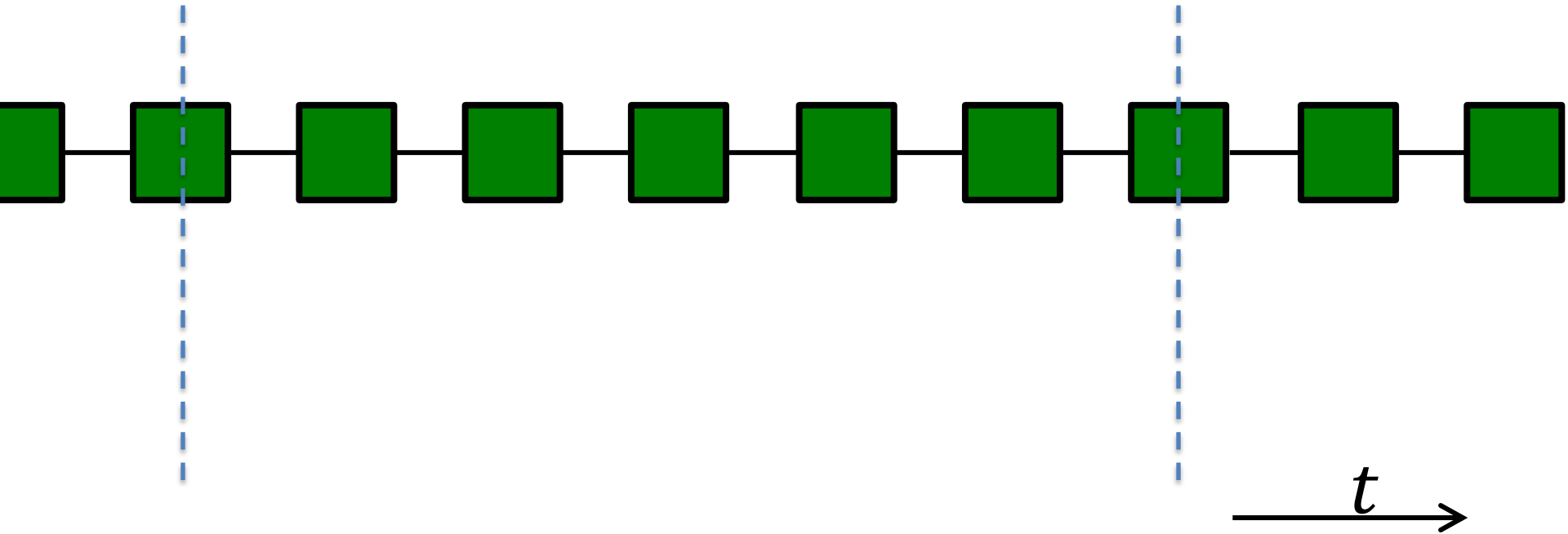
- **Longest** chain wins
- Transactions on short chain are reverted

# Fork Resolution

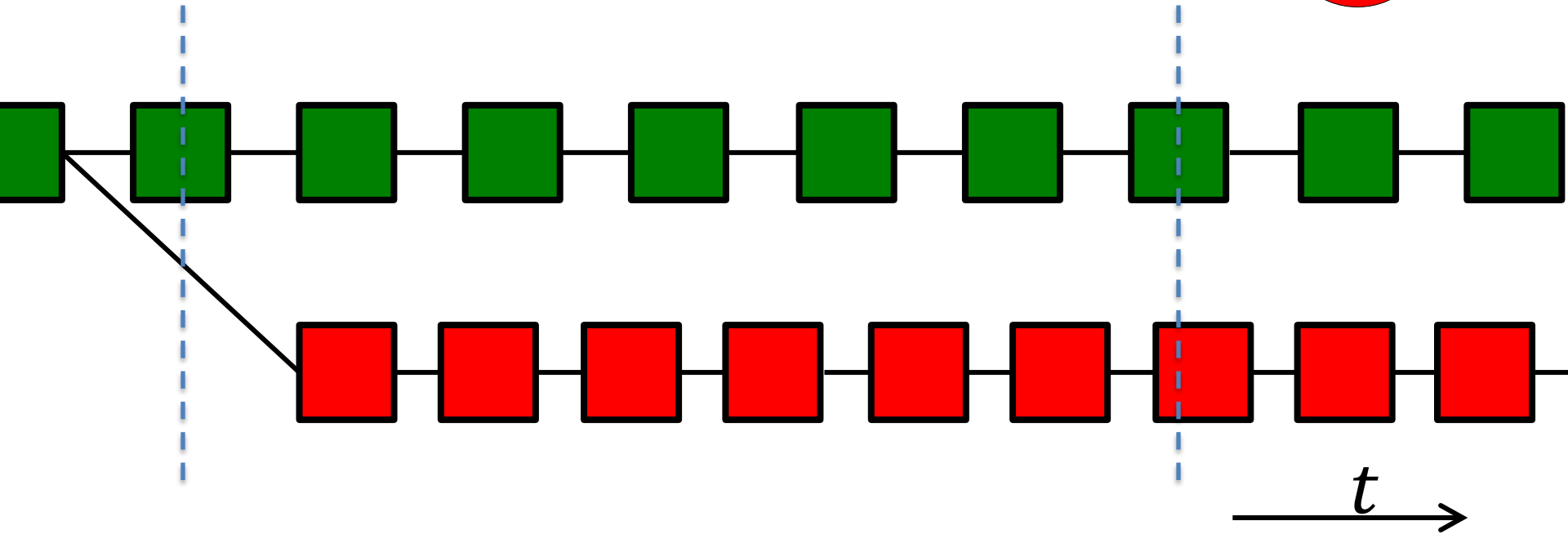
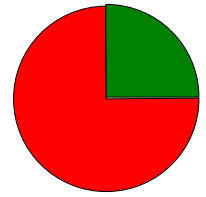


A transaction is **confirmed** when  
it is **buried** “deep enough”  
(typically 6 blocks – i.e., one hour)

# Security Threat!

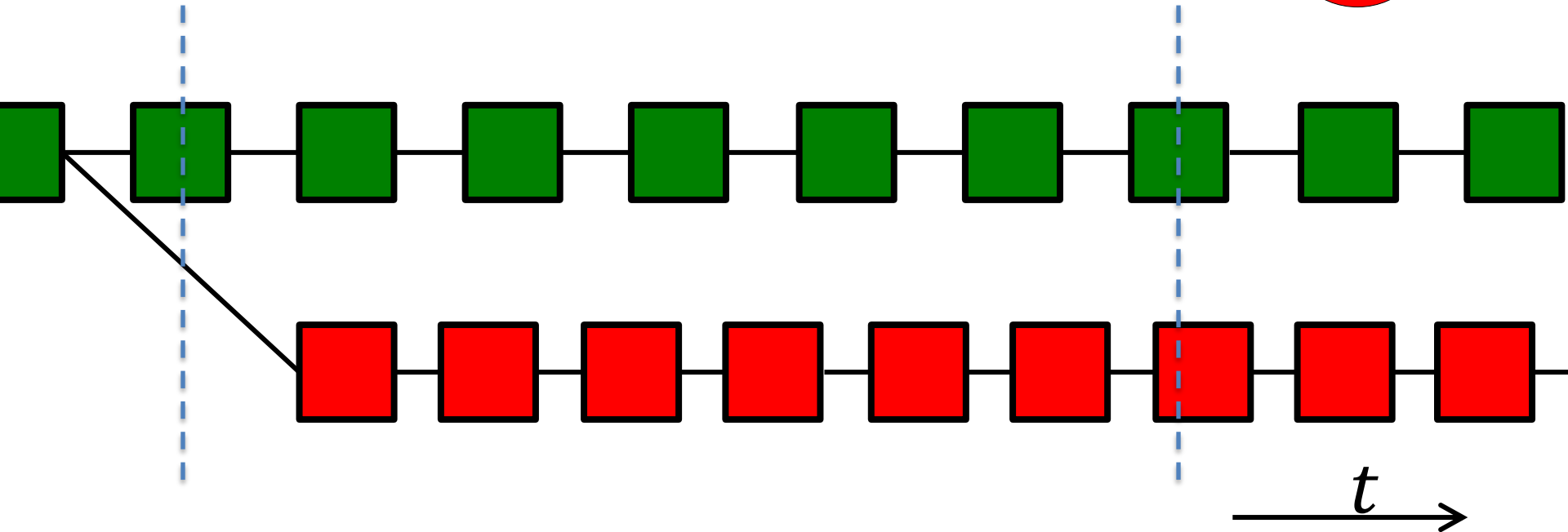
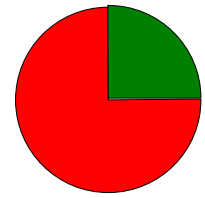


# Security Threat!

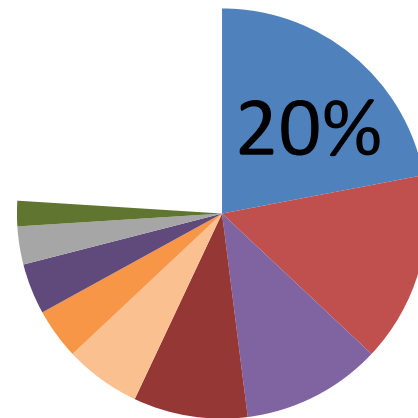


Threat: attacker outruns good miners

# Security Threat!



Threat: attacker outruns good miners  
→ **Security Assumption:** *good miners own  $>.5$  of the total compute power*



[blockchain.info,  
April 2015]

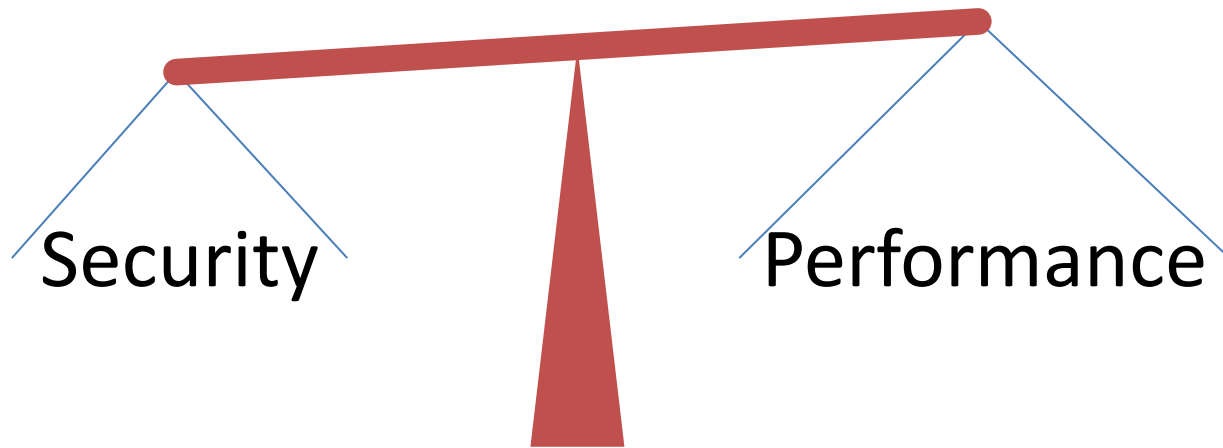
# Bitcoin Parameters

- block size = 1 MByte
- *target* set such that  $E(\text{block interval}) = 10 \text{ minutes}$ 
  - the block size is small enough and the block interval is large enough such that all miners can learn about a new block (through “gossiping”) before the next block is mined
- Results in fewer than 10 transactions per second
- *Why not a larger block or a smaller interval?*



# Security-Performance Tradeoff

Nakamoto's Blockchain exhibits a tradeoff:  
[Sompolinsky+'15, Lewenberg+'15]



# Metrics

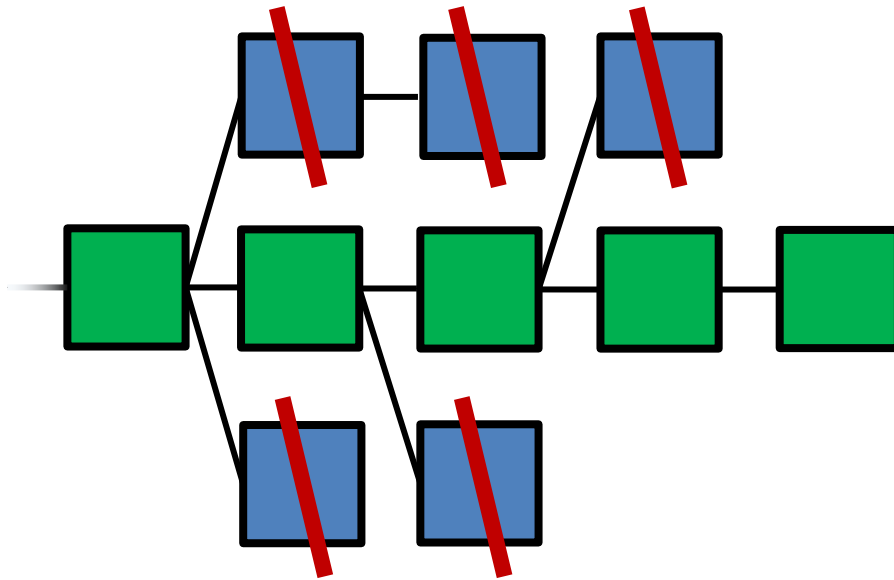
## **Performance:**

- Throughput
- Latency

## **Security:**

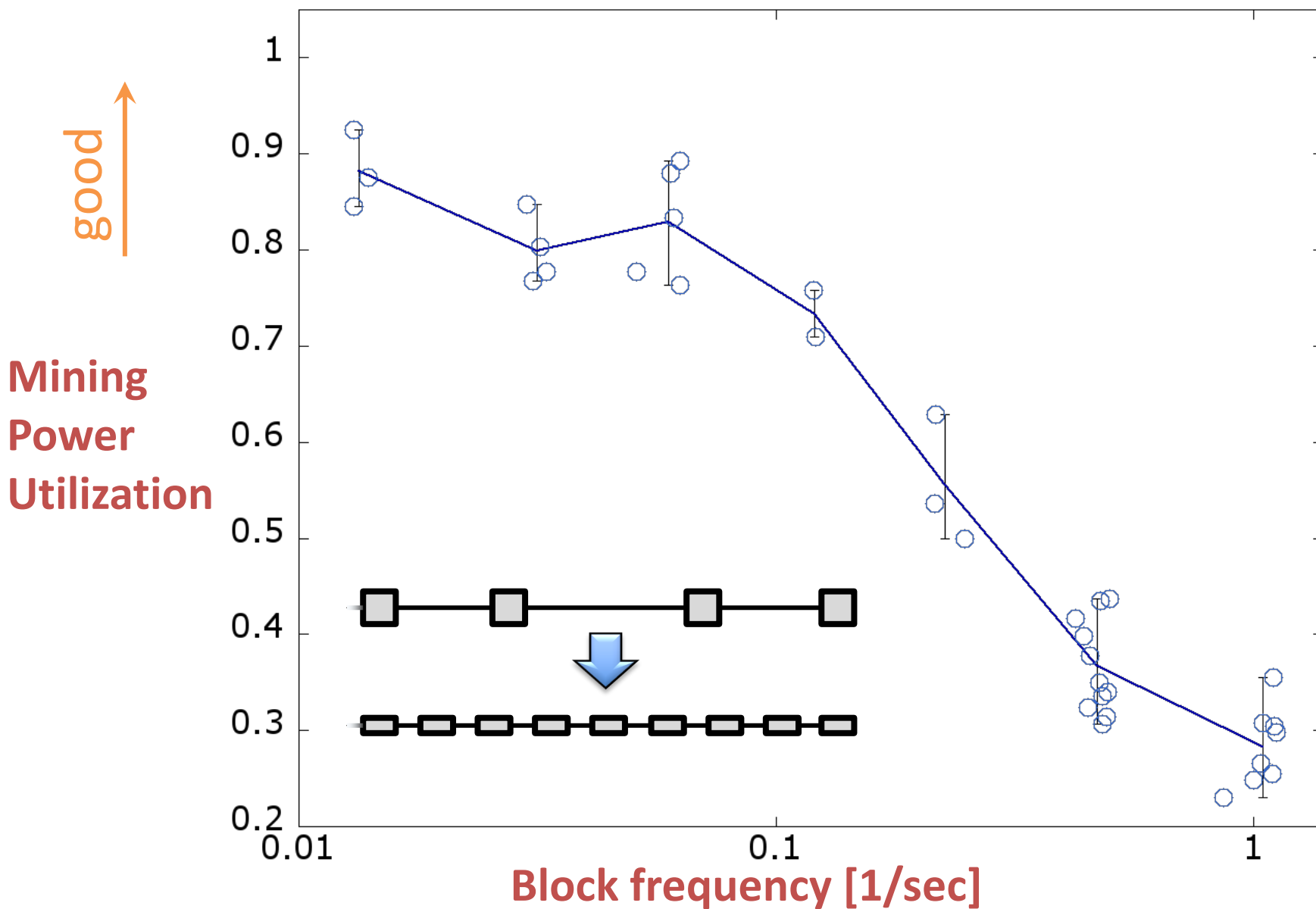
- Mining power utilization
- Fairness

# Mining Power Utilization

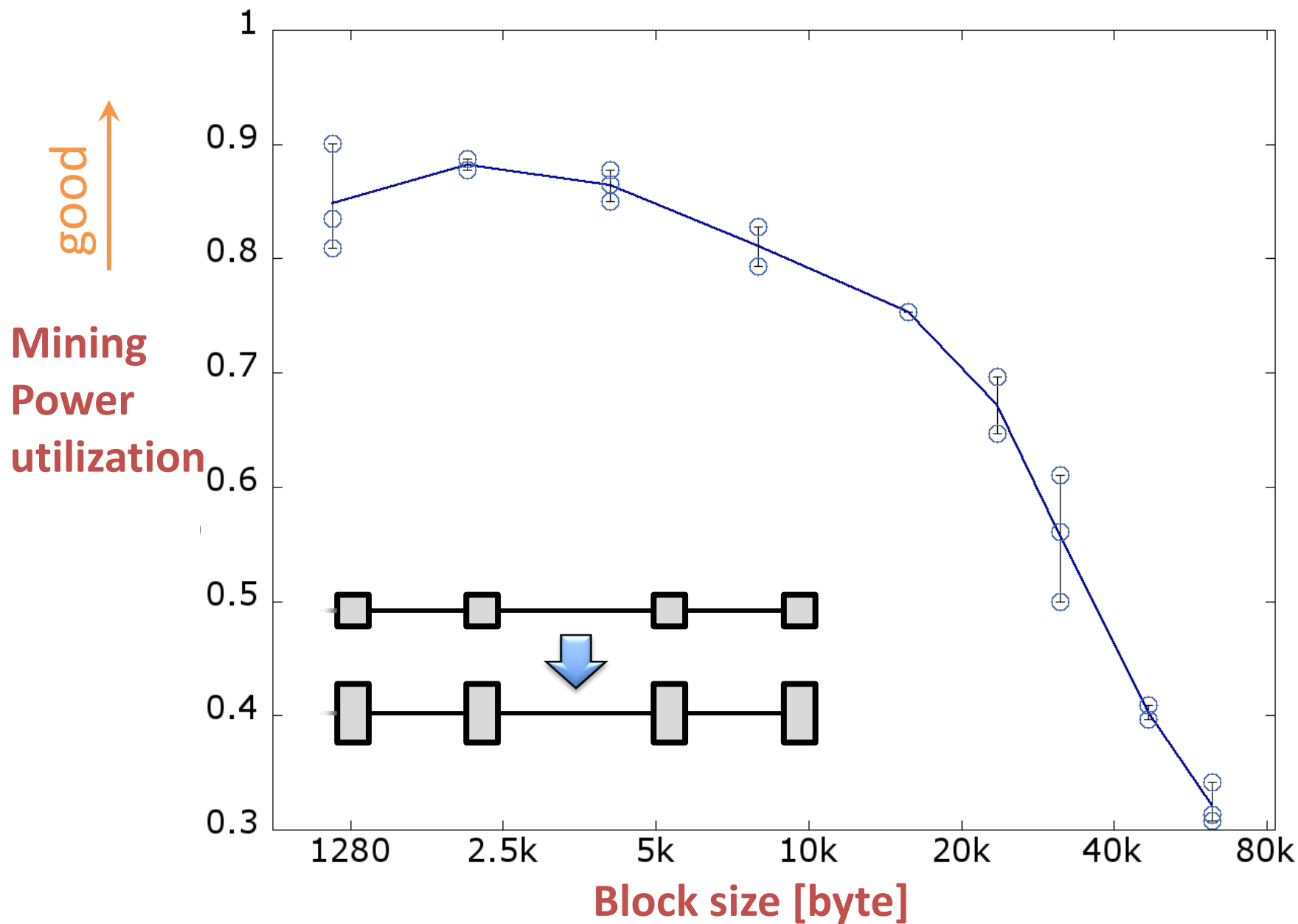


**Attacker only has to out-run the main chain**

# Block Frequency



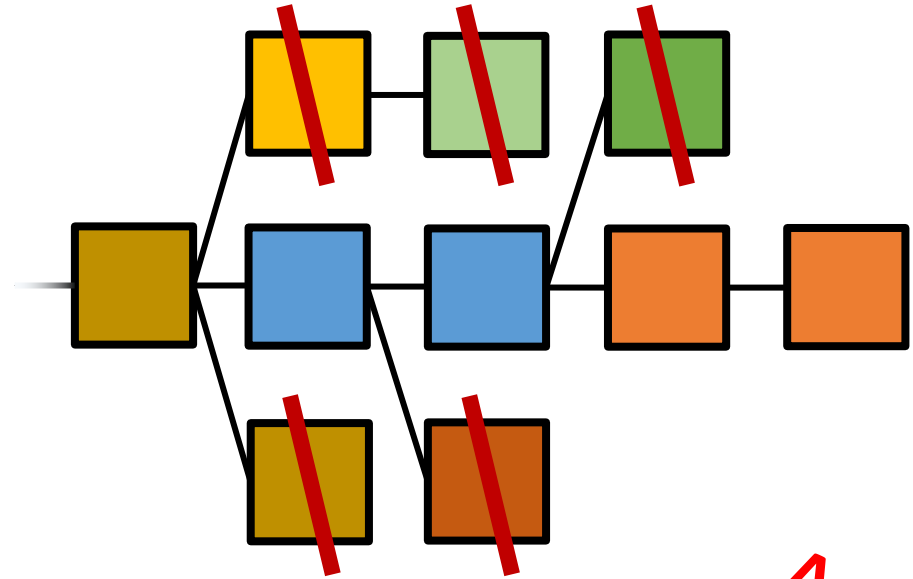
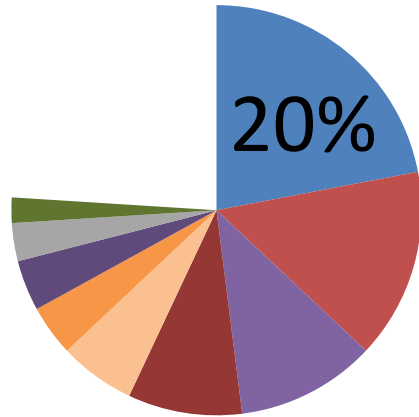
# Block Size



# Fairness

## Known Miner Sizes

[blockchain.info, April 2015]



Presence:

$$\frac{\sum_{all} \neg \square}{\sum_{all}} = 80\%$$

*Fair*

$$\frac{\sum_{main} \neg \square}{\sum_{main}} = 60\%$$

*Actual*

→ tendency towards centralization

# Permissionless vs Permissioned Blockchains

	Permissionless	Permissioned
Approach	Competitive	Cooperative
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