CS 5416 Recitation 5 multi-threading II

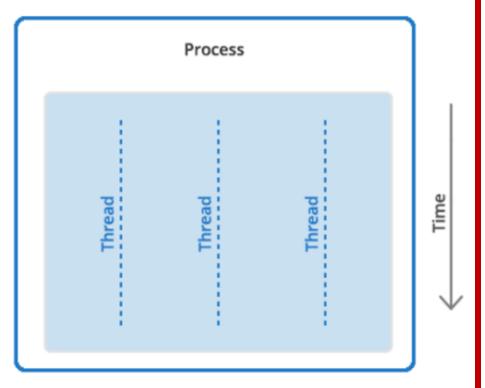
09/26/2025

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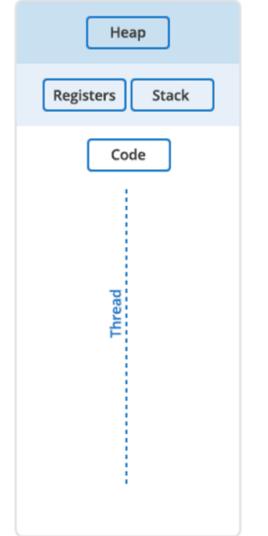
Recap

- Multithreading
- Race condition

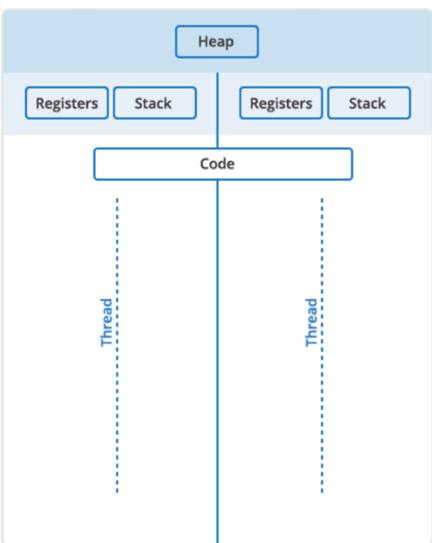
Concurrency



Single Thread



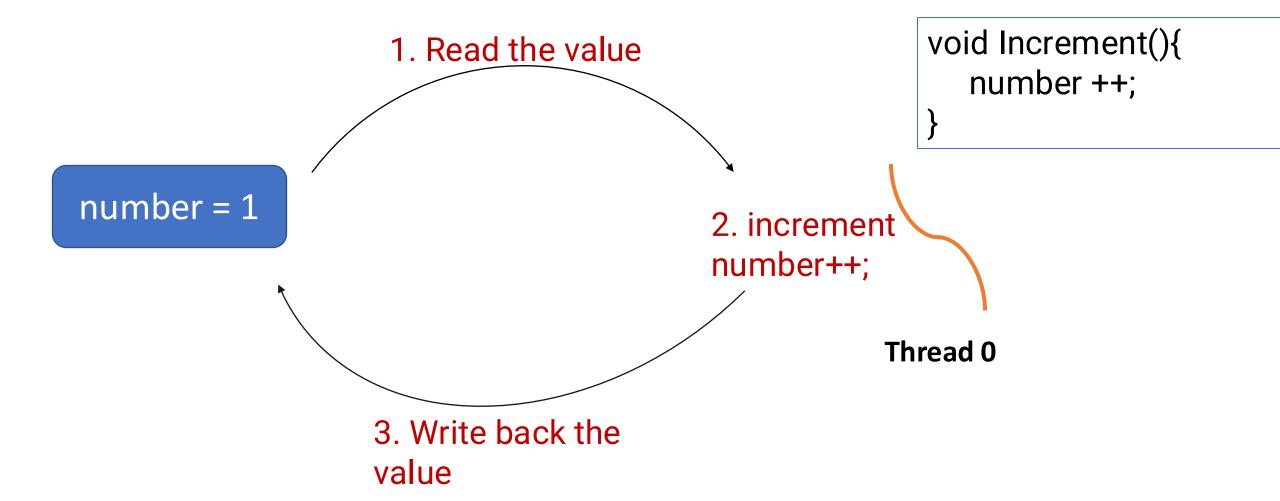
Multi Threaded

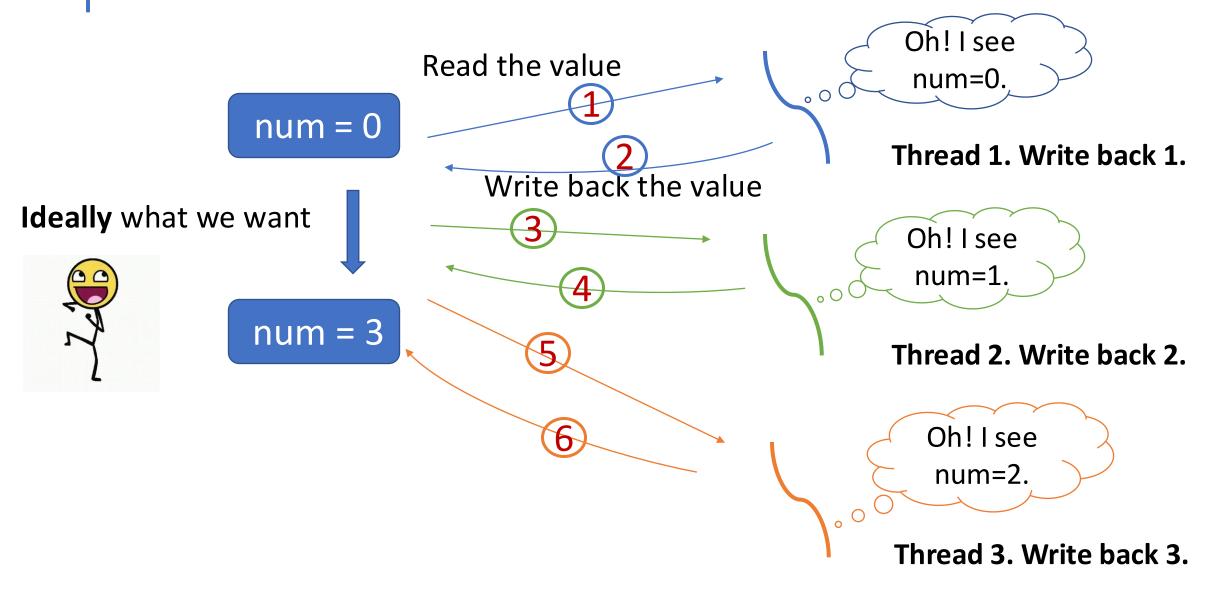


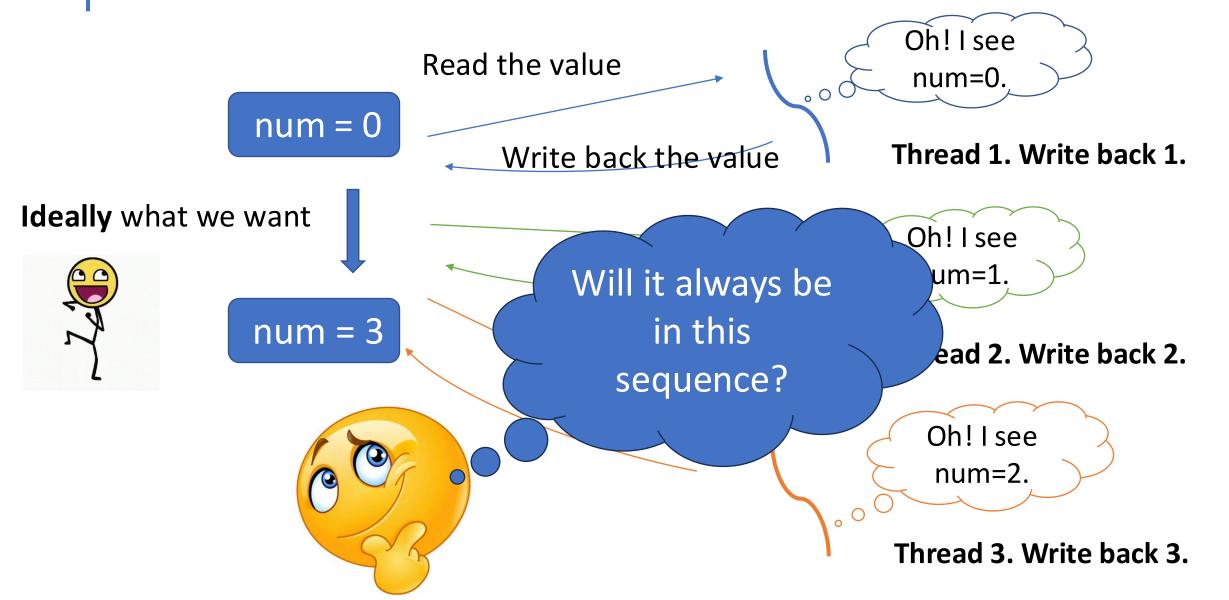
- Race condition:
 - The situation where the outcome depends on the relative ordering of execution of operations on two or more threads;
 the threads race to perform their respective operations.

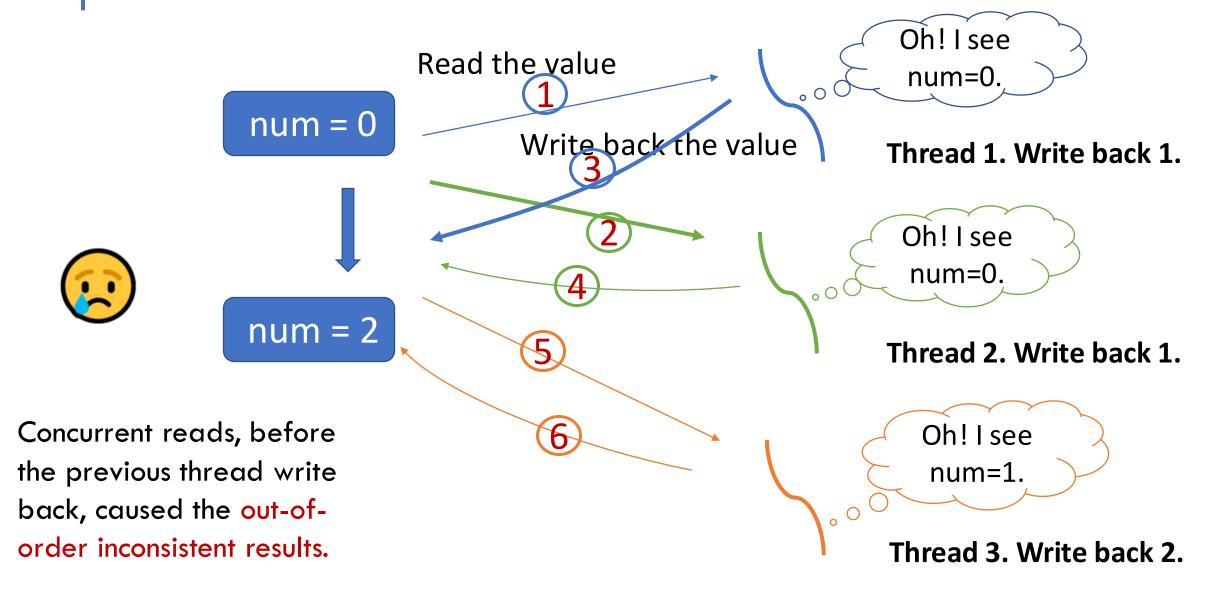
- Example: Concurrent increments of a shared integer variable.
 - Each thread shares an integer called count initialized to 0, increments it 1 million times concurrently without any synchronization

Number of threads	Final value
1	1000000
2	1059696
3	1155035
4	1369165









Thread Safety

• A function, a piece of code, or an object is **thread-safe** when it can be **invoked** or **accessed concurrently** by **multiple threads without** causing unexpected behavior, race conditions, or data corruption.

Thread safe

• Entities in C++ standard library and their thread-safety guarantees

Thread safe?

- Is integer type inherently thread-safe?
 - No, as we showed just now





std::atomic

• A template that defines an atomic type.



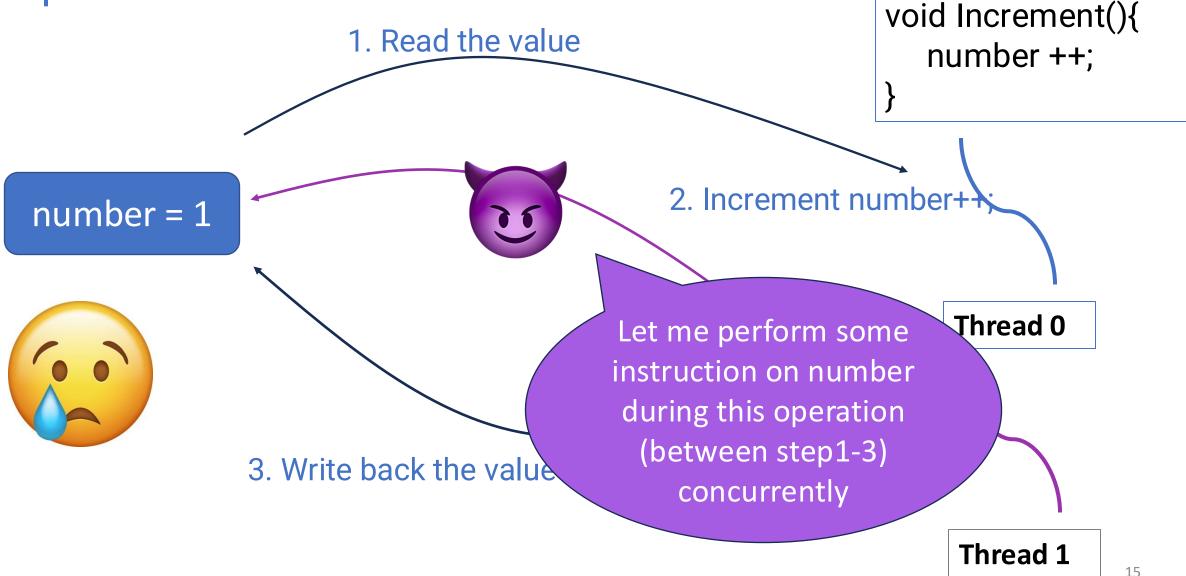
	<pre>template< class T > struct atomic;</pre>	(1)	(since C++11)
	<pre>template< class U > struct atomic<u*>;</u*></pre>	(2)	(since C++11)
* ere at	<pre>template< class U > struct atomic<std::shared_ptr<u>>;</std::shared_ptr<u></pre>	(3)	(since C++20)
end of ition if time)	<pre>template< class U > struct atomic<std::weak_ptr<u>>;</std::weak_ptr<u></pre>	(4)	(since C++20)

(more at the end of recitation if have time)

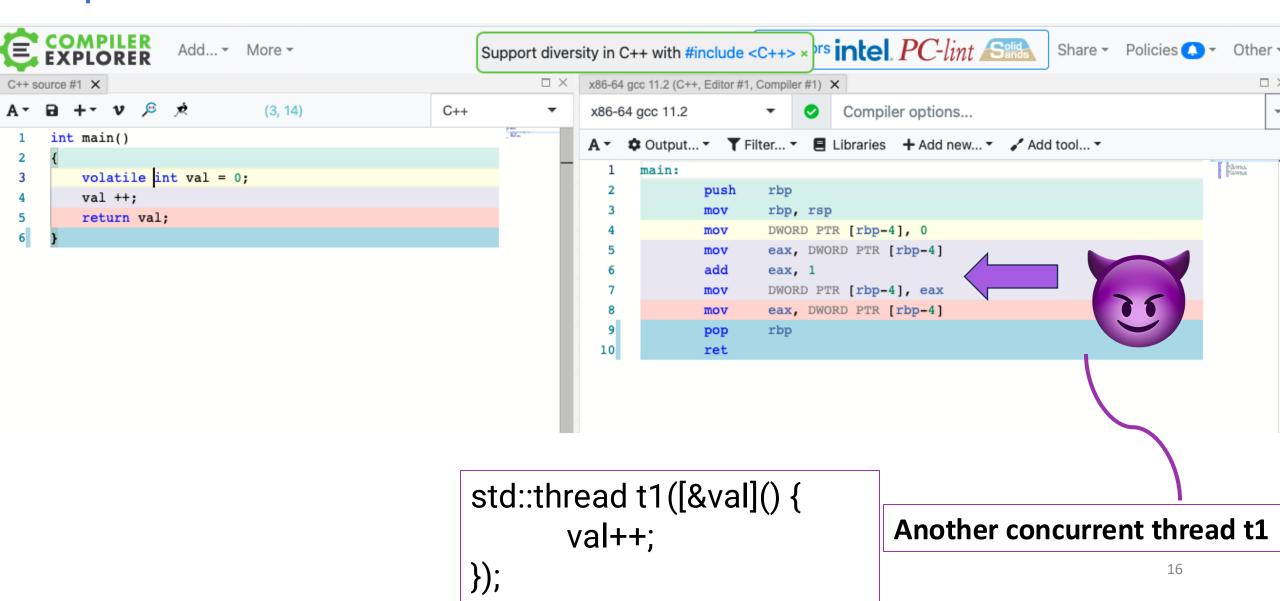
Atomic

- An atomic operation is an indivisible operation.
- The operation is **either done or not done**. Such an operation would **never be half-done** from any thread in the system.

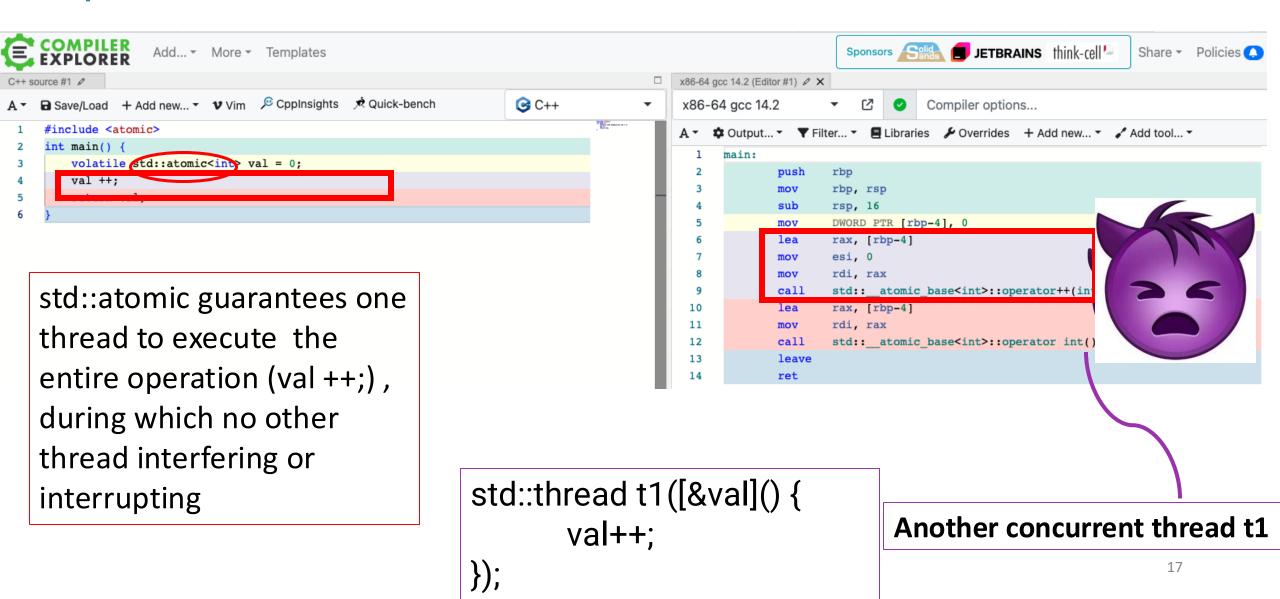
Data race condition: non-atomic access pattern



Data race condition: non-atomic access pattern



Atomic access



Atomic

- An atomic operation is an indivisible operation.
- std::atomic are implemented using hardware supports provided by modern CPU:
 - Examples of atomic instructions:
 - Compare-and-Swap (CAS)
 - Load-Linked/Store-conditional (LL/SC)
 - fetch_and_add (FAA)
 - **Different CPUs** provide **different sets** of **atomic instructions**. The implementation of std::atomic varies from architecture to architecture

Atomic member functions

```
Atomic type: std::atomic<type>
```

- Constructor
 std::atomic<bool> x(true);
 std::atomic<uint32_t> y(0);
- store() x.store(false); y.store(1, std::memory_order_relaxed);

More atomic member functions

- load()
- exchange()
- operator=
- operator+=, operator -=
- operator++, operator--

```
bool z = x.load();
uint32_t m = y.exchange(100);
y = 2;
y += 1; y.fetch_add(1); (since C++20)
y ++;
```

What about
$$y = y + 1$$
?

More atomic member functions

```
load()
exchange()
operator=
operator+=, operator -=
operator++, operator--
bool z = x.load();
uint32_t m = y.exchange(100);
y = 2;
y += 1; y.fetch_add(1);
y ++;
```

What about y = y + 1?

When multithreading, leads to **race condition**, because it involves multiple operations (read x, +1 and then assignment operation)

Thread safe

- std::atomic
- std::shared_ptr

std::vector

Does std::vector guarantee thread-safety?

Multithreads' data sharing with std::vector

- When is std::vector thread-safe?
 - Each thread has its own instance of std::vector (no concurrency)
 - Read-only access
- When is std::vector not thread-safe?
 - Simultaneous Read and Write
 - Concurrent modification
 - Reallocation access on reallocation or modification

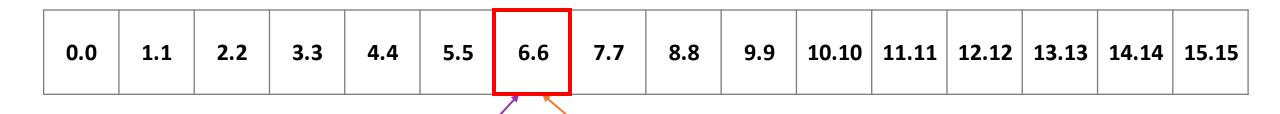
Read-only-access of std::vector

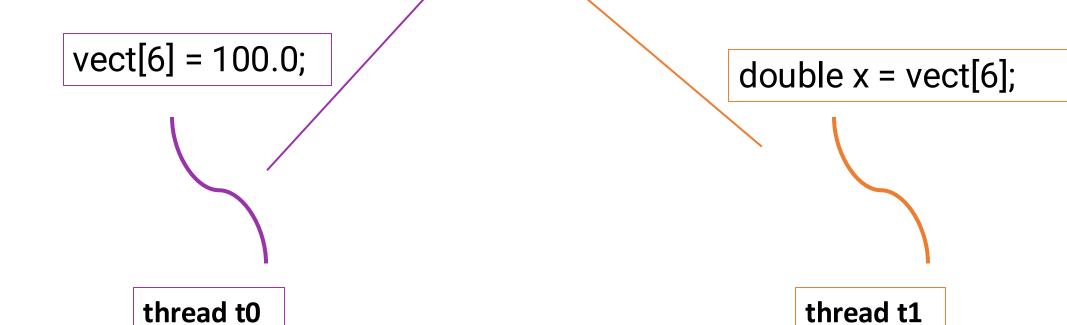


```
int main() {
  std::vector<double> vec(100, 1.00);
  double t1_sum;
  double t2_sum;
  std::thread t1(read_vector,std::ref(vec), 1, std::ref(t1_sum));
  std::thread t2(read_vector,std::ref(vec), 2, std::ref(t2_sum));
  t1.join();
  t2.join();
  std::cout << "t1_sum="<< t1_sum << ",t2_sum=" << t2_sum;
  ...}</pre>
```

Thread safe, because only concurrent reads

Simultaneous read and write





Simultaneous read and write

Concurrent Read+write to the **SAME** element is **NOT** thread-safe



vect[6] = 100.0;

thread t0

double x = vect[6];

x could be 6.6 or 100.0 after this.

thread t1

Concurrent access with reallocation

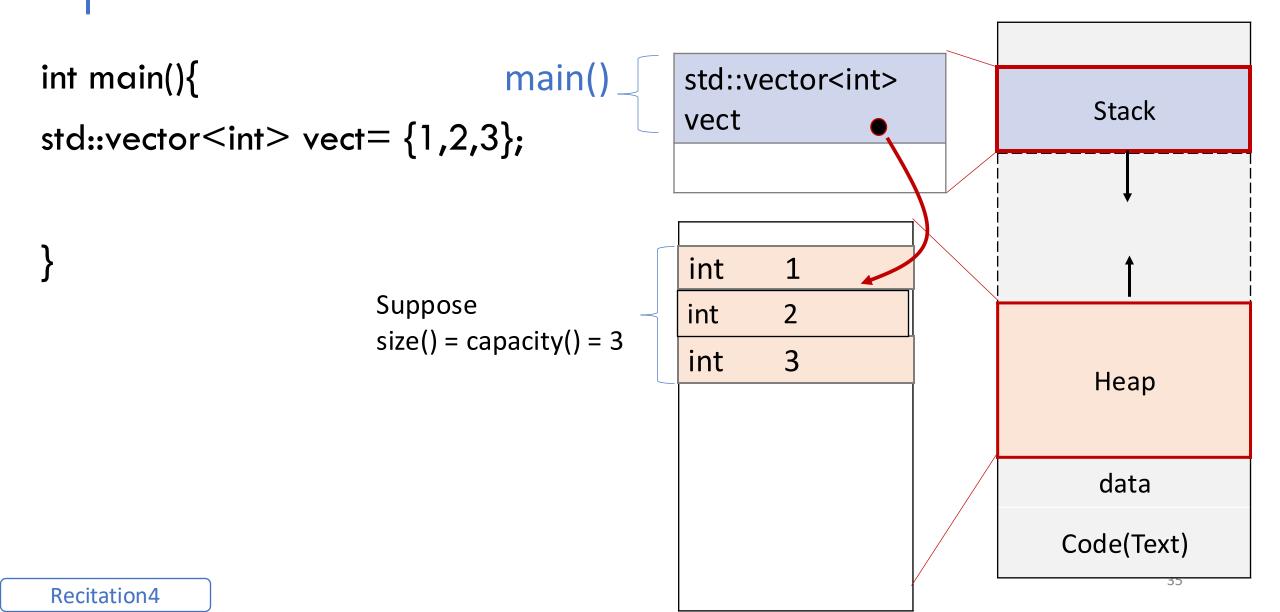


Is this code thread-safe?

```
void add_elements(std::vector<int>& vec, int thread_id) {
   for (int i = 0; i < 10; ++i) {
     vec.push_back(i);
   }
}</pre>
```

```
int main() {
    std::vector<int> vec = {1, 2, 3};
    std::thread t1(writer, std::ref(vec));
    std::thread t2([&vec]() {
        std::cout << "value: " << vec.back(); << std::endl;});
    t1.join();
    t2.join();
...}</pre>
```

How is std::vector allocated in memory

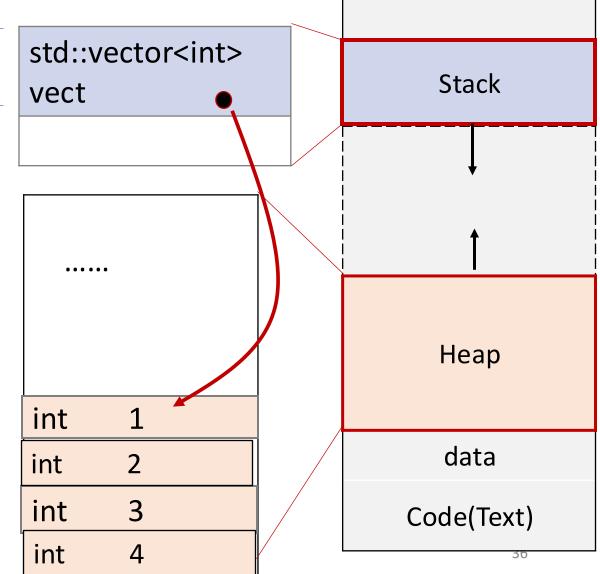


How is std::vector allocated in memory

```
int main(){
    std::vector<int> vect= {1,2,3};
    vect.push_back(1);
}
```

after the operation the new size() is greater than old capacity(),

- a reallocation takes place
- all iterators and all references to the elements are invalidated.



Concurrent access with reallocation

```
void add_elements(std::vector<int>& vec, int thread_id) {
   for (int i = 0; i < 10; ++i) {
     vec.push_back(i);
   }
}</pre>
```

Not thread safe:

one thread is modifying the std::vector (push_back), while another thread reads from it (back()), there's a risk of data races.

```
int main() {
    std::vector<int> vec = {1, 2, 3};
    std::thread t1(writer, std::ref(vec));
    std::thread t2([&vec]() {
        std::cout << "value: " << vec.back(); << std::endl;});
    t1.join();
    t2.join();
...}</pre>
```

Locking

Locking

---protecting data with mutex



- How does mutex work?
 - Before accessing a shared data structure, you lock the mutex associated with that data
 - When finished accessing the data structure, you unlock the mutex.



std::mutex



exclusive, non-recursive ownership

- A thread owns the mutex from the time when it call lock() until it calls unlock()
- The Thread Library then ensures that once one thread has locked a specific mutex, all other threads that try to lock the same mutex have to wait until the thread that successfully locked the mutex unlocks it.



Locking

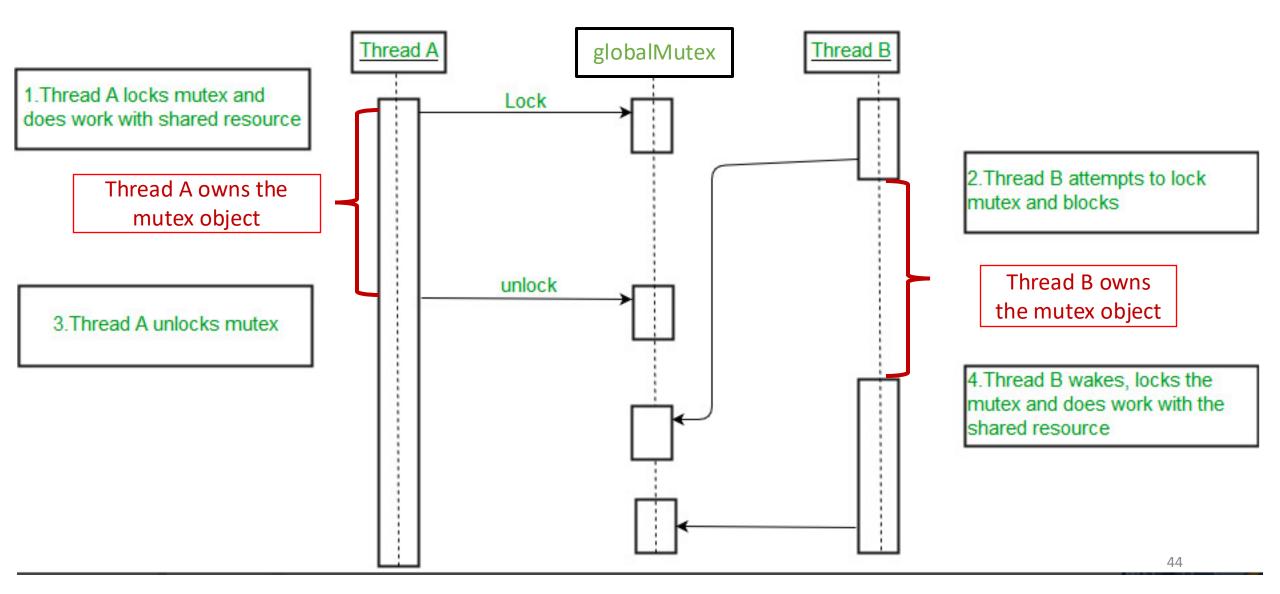
---std::mutex::lock(), unlock()

```
global_num = 0;
    std::mutex
                  globalMutex;
3
    void incre(int num){
            globalMutex.lock();
4
            global_num = global_num + 1;
5
            globalMutex.unlock();
6
    int main(){
8
9
            std::thread threadA(incre, 10);
            std::thread threadB(incre, 10);
10
            threadA.join();
11
            threadB.join();
12
```

Only one thread could enter line 5 at a time

Mutex and Lock in C++







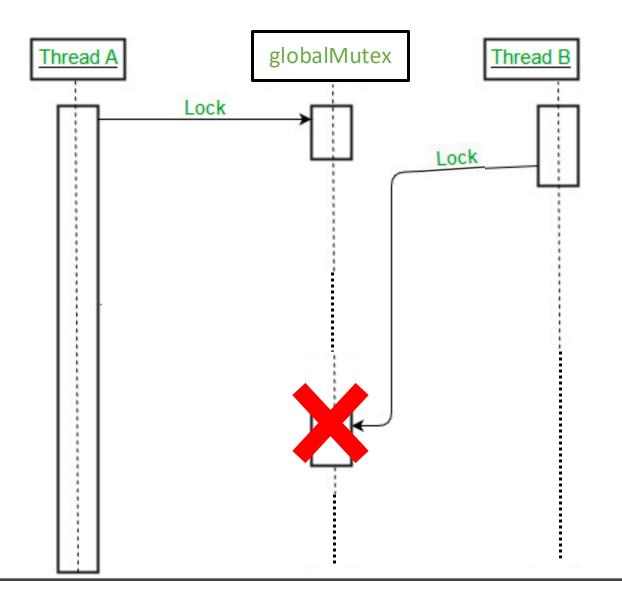
---std::mutex::lock(), unlock()

```
global_num = 0;
int
std::mutex
                 globalMutex;
                                                           Now, what will
void incre(int num){
                                                        happen, if I forget to
         globalMutex.lock();
                                                        call mutex.unlock()?
        global_num = global_num + 1;
         globalMutex.unlock();
int main(){
        std::thread threadA(incre, 10);
        std::thread threadB(incre, 10);
        threadA.join();
        threadB.join();
```

Mutex and Lock in C++



Thread A locks mutex and does work with shared resource



2.Thread B attempts to lock mutex and blocks

Thread B is unable to acquire the lock if Thread A doesn't unlock it.

Mutex and Lock in C++



- A Mutex is a lock that we set before using a shared resource and release after using it.
- When the lock is set by one thread, then **no other thread** can access the locked region of code.
- Mutex lock could only be released by the thread who locked it.

---std::mutex::lock(), unlock()

- std::mutex::lock(), unlock()
 - It is not recommended practice to call lock(), unlock() directly,
 because this means that you have to remember to call unlock() on
 every code path out of a function that called lock(), including those
 due to exceptions.

- Resource acquisition must succeed for initialization to succeed:
 - In RAII, holding a resource is a class invariant is tied to object lifetime: resource allocation is done during object creation, by the constructor; while resource deallocation is done during object destruction, by the destructor.

- If there are no object leaks, there are no resource leaks.
 - The resource is guaranteed to be held between when initialization finishes and finalization starts, and to be held only when the object is alive.

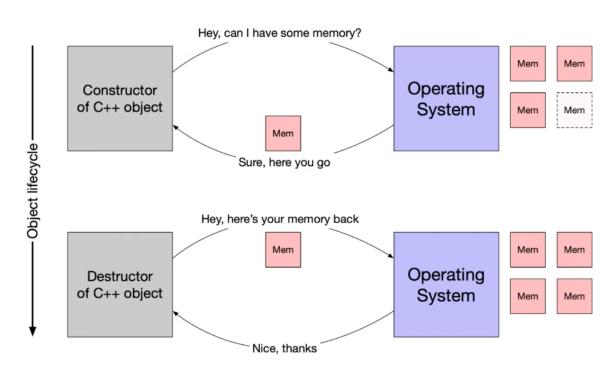
```
// problem #1
  int *arr = new int[10];
   // arr goes out of scope but we didn't delete it, we now have a memory leak 😥
// problem #2
   std::thread t1( [] () {
         // do some operations
   });
                       // thread t1 is created but not joined, if it goes out of scope, std::terminate is
                       called, this implementation doesn't properly handle the thread's life cycle 😥
// problem #3
Std::mutex globalMutex;
Void func() {
   globalMutex.lock();
      // if we never unlocked the mutex(or exception occurred before unlock),
                                                                                                     50
   it will cause a deadlock when other thread tries to acquire this lock 😥
```

```
// problem #1's fix
  int *arr = new int[10];
  delete[] arr;
// problem #2's fix
   std::thread t1( [] () {
         // do some operations
  });
  t1.join();
// problem #3's fix
Std::mutex globalMutex;
Void func() {
   globalMutex.lock(); ....
   globalMutex.unlock();
                                                                                                     51
```

RAII

• When acquire resources in a constructor, also need to release them in the corresponding destructor

- Resources:
 - Heap memory,
 - files,
 - sockets,
 - mutexes





---std::mutex::lock(), unlock()

```
global_num = 0;
int
std::mutex
                 globalMutex;
void incre(int num){
         globalMutex.lock();
        global_num = global_num + 1;
         globalMutex.unlock();
int main(){
        std::thread threadA(incre, 10);
        std::thread threadB(incre, 10);
        threadA.join();
        threadB.join();
```

Is there a better ways to manage the mutex that can automatically unlock it when not used?



Mutex and RAII locks



- std::unique_lock
- std::scoped_lock
- std::shared_lock

```
{
  std::unique_lock<std::mutex> lck(my_mutex);
    ... ...
}
```

```
std::shared_mutex shared_mutex;
{
    std::shared_lock<std::mutex> lck(shared_mutex);
    ... ...
}
```

- A unique lock is an **object** that **manages a mutex object** with unique ownership in both states: locked and unlocked.
- RAII: When creating a local variable of type std::unique_lock passing the mutex as parameter.
 - On construction, the object acquires a mutex object, for whose locking and unlocking operations becomes responsible.
 - This class guarantees an unlocked status on destruction (even if not called explicitly).

Features:

• Deferred locking, Timeout locks, adoption of mutexes, movable(transfer of ownership)

---unique_lock

```
global_num = 0;
    int
    std::mutex
                 globalMutex;
2
    void incre(int num){
3
           std::unique_lock<std::mutex> u_lock(globalMutex);
           global_num = global_num + 1;
6
    int main(){
8
           std::thread t1(incre, 1);
9
           std::thread t2(incre, 3);
10
           t1.join();
11
           t2.join();
12
```

Only one thread could enter line 5-7 at a time

---unique_lock

Unique_lock feature: Deferred locking

```
std::mutex mtx;
void conditional_locking(bool should_lock) {
    // Create lock but do not acquire it
   std::unique_lock<std::mutex> lock(mtx, std::defer_lock);
if (should_lock) {
    lock.lock();
                 // Conditionally acquire the lock
    std::cout << "Lock acquired." << std::endl;
  } else {
    std::cout << "Lock not acquired." << std::endl;
```

```
int main() {
    std::thread t1(conditional_locking, true);
    std::thread t2(conditional_locking, false);
    t1.join();
    t2.join();
    return 0;
}
```

• Scoped_lock: a mutex wrapper which obtains access to (locks) the provided mutex, and ensures it is unlocked when the scoped lock goes out of scope

```
global_num = 0;
     int
                      qlobalMutex;
     std::mutex
3
     void incre(int num){
4
                      std::scoped_lock s_lock(globalMutex);
6
                      global_num = global_num + 1;
8
             global_num = global_num + 1;
9
10
              ...
11
```

• std::shared_lock allows for shared ownership of mutexes.

```
std::shared_mutex mtx;
int global_val;
void print_val (int n, char c) {
  std::shared_lock<std::shared_mutex > lck (mtx);
  std::cout << global_val << std::endl;</pre>
int main () {
   std::thread th1 (print_val);
   std::thread th2 (print_val);
   th1.join();
   th2.join();
```

```
// problem #1
{
  int *arr = new int[10];
}  // arr goes out of scope but we didn't delete it, we now have a memory leak ②
```

RAII

Better fixes



```
// problem #1's fix
{
     std::unique_ptr<int[]> arr(new int[10]);
.....
}
```

```
// problem #3's fix
Std::mutex globalMutex;
Void func() {
    std::unique_lock<std::mutex> lock(globalMutex);
....
}
```

- scoped_lock()
- unique_lock()
- shared_lock()



std::scoped_lock

a mutex wrapper which obtains access to (locks) the provided mutex, and ensures it is unlocked when the scoped lock goes out of scope

When does s_lock get released?

std::scoped_lock

```
std::vector<int> my_vec;
std::mutex
                my_mutex;
void add_to_list(int new_value) {
        std::scoped_lock<std::mutex> lck(my_mutex);
        my_vec.push_back(new_value);
bool list_contains(int value_to_find) {
        std::scoped_lock<std::mutex> lck(my_mutex);
        return std::find(my_vec.begin(), my_vec.end(),value_to_find) != my_vec.end();
```

- scoped_lock()
- unique_lock()
- shared_lock()

std::unique_lock

- A unique lock is an object that manages a mutex object with **unique** ownership in both states: locked and unlocked.
- RAII: When creating a local variable of type std::unique_lock passing the mutex as parameter.
 - On construction, the object acquires a mutex object, for whose locking and unlocking operations becomes responsible.
 - This class guarantees an unlocked status on destruction (even if not called explicitly).

Features:

• Deferred locking, Timeout locks, adoption of mutexes, movable(transfer of ownership)

- scoped_lock()
- unique_lock()
- shared_lock()

std::shared_lock

Shared_lock allows for shared ownership of mutex. More than one thread could hold the mutex at the same time.

```
std::shared_mutex mtx;
int global_val;
void print_val (int n, char c) {
  std::shared_lock<std::shared_mutex > lck (mtx);
  std::cout << global_val << std::endl;
int main () {
  std::thread th1 (print_val);
  std::thread th2 (print_val);
  th1.join();
  th2.join();
```

Exercise

- How can I use the RAII class locks to implement R/W lock?
 - R/W locks allow multiple readers at the same time
 - But if there is writer, then there should be no readers, and only one writers.

Where to find the resources?

Concurrency programing:

- Book: C++Concurrency in Action Practice Multithreading
- https://learn.microsoft.com/en-us/archive/blogs/ericlippert/what-is-this-thing-you-call-thread-safe

Notes:

- Atomic built-in: https://gcc.gnu.org/onlinedocs/gcc-4.4.3/gcc/Atomic-Builtins.html
- Memory order: https://cplusplus.com/reference/atomic/memory order/#google vignette