



CS519: Computer Networks

Lecture 9: May 03, 2004

Media over Internet



Media over the Internet



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- Media = Voice and Video
- Key characteristic of media:
 - Realtime
 - Which we've chosen to define in terms of *playback*, not “latency over the network”
- A digitized sample of media must be “played back” at a precise time (relative to the previous sample)



Media samples

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- Media is sampled (and played out) at uniform time period
- CD quality audio: 44100 samples per seconds with 16 bits per sample, stereo sound
 - $44100 * 16 * 2 = 1.411 \text{ Mbps}$
- Telephone quality voice: 8K samples per second, 8 bits per sample
 - $8000 * 8 = 64 \text{ Kbps}$



Media samples

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- Video
 - For 320*240 images with 24-bit colors
 - $320 \times 240 \times 24 = 230\text{KB/image}$
 - 15 frames/sec: $15 \times 230\text{KB} = 3.456\text{MBps} = 27.6 \text{ Mbps}$



MPEG “compression”



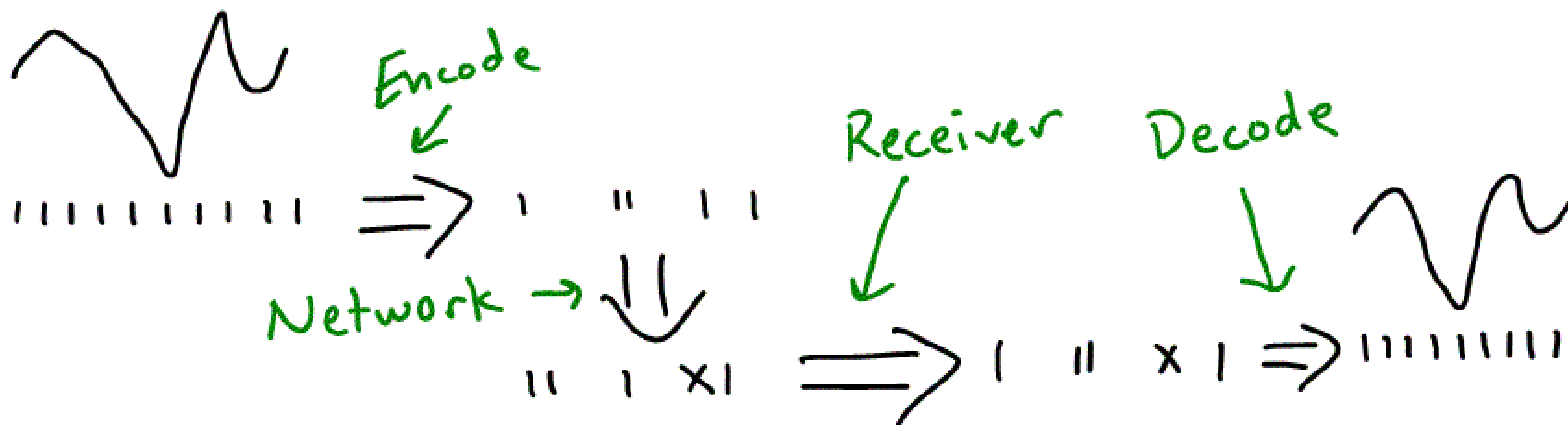
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- MP3 audio compression
 - Typical rates are 96kbps, 128kbps, 160kbps
 - From 1.4Mbps: 14.6x, 10.9x, and 8.75x reduction respectively
 - With very little perceived degradation!
- MPEG1 and MPEG2 video compression
 - 1.5Mbps – 6Mbps
 - From 27.6Mbps: 18.4x – 4.6x reduction

What does this compression mean to us?

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- Compressing periodic, fixed-size samples produces:
 - non-periodic, variable-size “units”





It's all about receive buffer...



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- Receiver must reproduce timing of original compressed packets
 - Timing was screwed up by the network (jitter and delay)
- The more we buffer at the receiver, the more jitter we can tolerate
 - Best case: download entire file before playing any of it
 - Worst case: conversational voice
- We mentioned this in QoS lecture . . .



Receive buffer considerations

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- Conversational voice: we can tolerate maybe 250ms latency
 - 150ms or less is better
 - After network delay, 150ms – 200ms buffering
- “Live” media: a few seconds latency ok
- Non-live streaming media: don’t want to wait too long for start of playback



Other realtime considerations

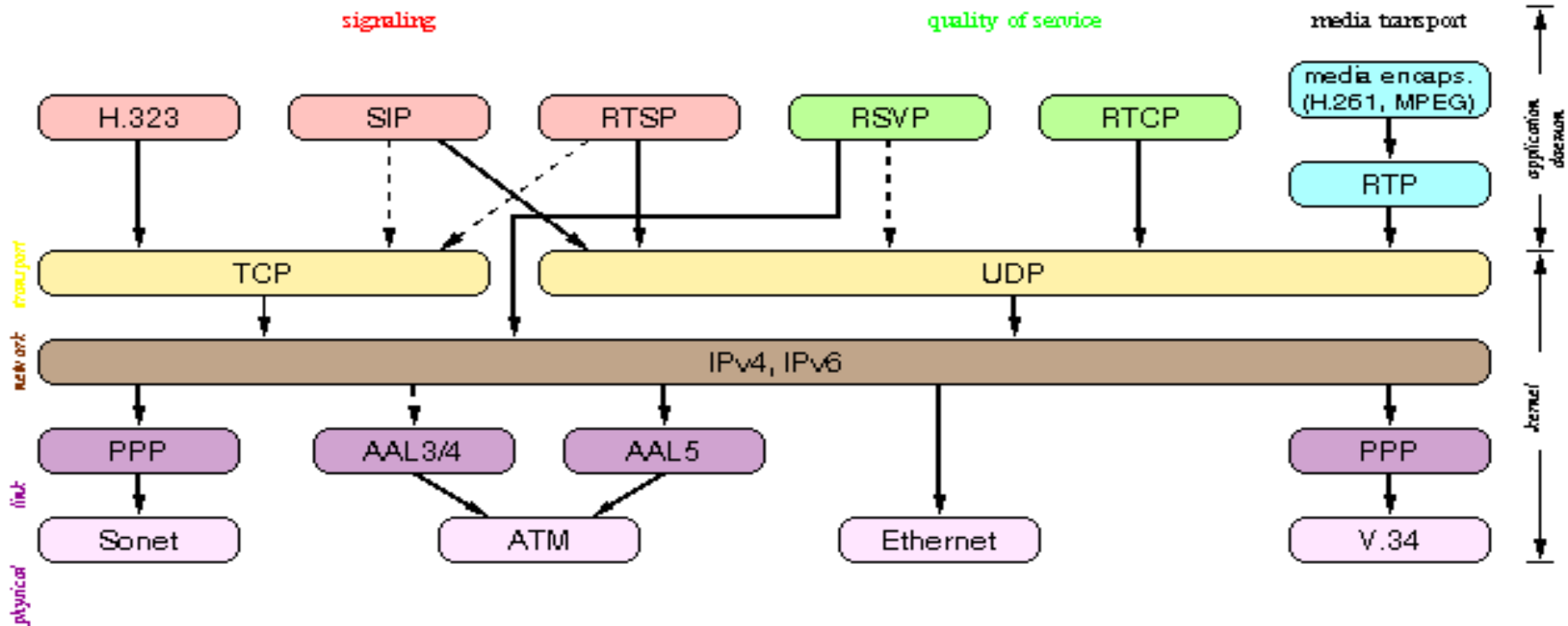


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- In addition to timing and variable size of compression units
- Encoding schemes have different loss tolerance
 - Can use FEC (Forward Error Correction) to an extent
- Some packets better to lose than others
- Encoding schemes may be able to slow down
 - At the expense of quality

Media-related protocols

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Real Time Protocol (RTP)

RFC 3550

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- Attempt to provide common transport for many types of media
- In addition to already-stated realtime requirements:
 - Must run over multicast
 - Must allow for “mixing” of streams (i.e. for conferencing)
 - Must be able to combine multiple streams
 - Multi-media, or layered encoding over multiple multicast groups



RTP design approach

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- Provide general header with broad capabilities
- Provide separate control protocol for managing RTP stream
 - RTCP: Real Time Control Protocol
- Each encoding type individually specifies how to use RTP



Some RTP usage profiles



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- 2029 RTP Payload Format of Sun's CellB Video Encoding.
- 2032 RTP Payload Format for H.261 Video Streams.
- 2035 RTP Payload Format for JPEG-compressed Video.
- 2038 RTP Payload Format for MPEG1/MPEG2 Video.
- 2190 RTP Payload Format for H.263 Video Streams.
- 2198 RTP Payload for Redundant Audio Data.
- 2250 RTP Payload Format for MPEG1/MPEG2 Video.
- 2343 RTP Payload Format for Bundled MPEG.
- 2429 RTP Payload Format for the 1998 Version of ITU-T Rec. H.263 ...
- 2431 RTP Payload Format for BT.656 Video Encoding.



More RTP usage profiles

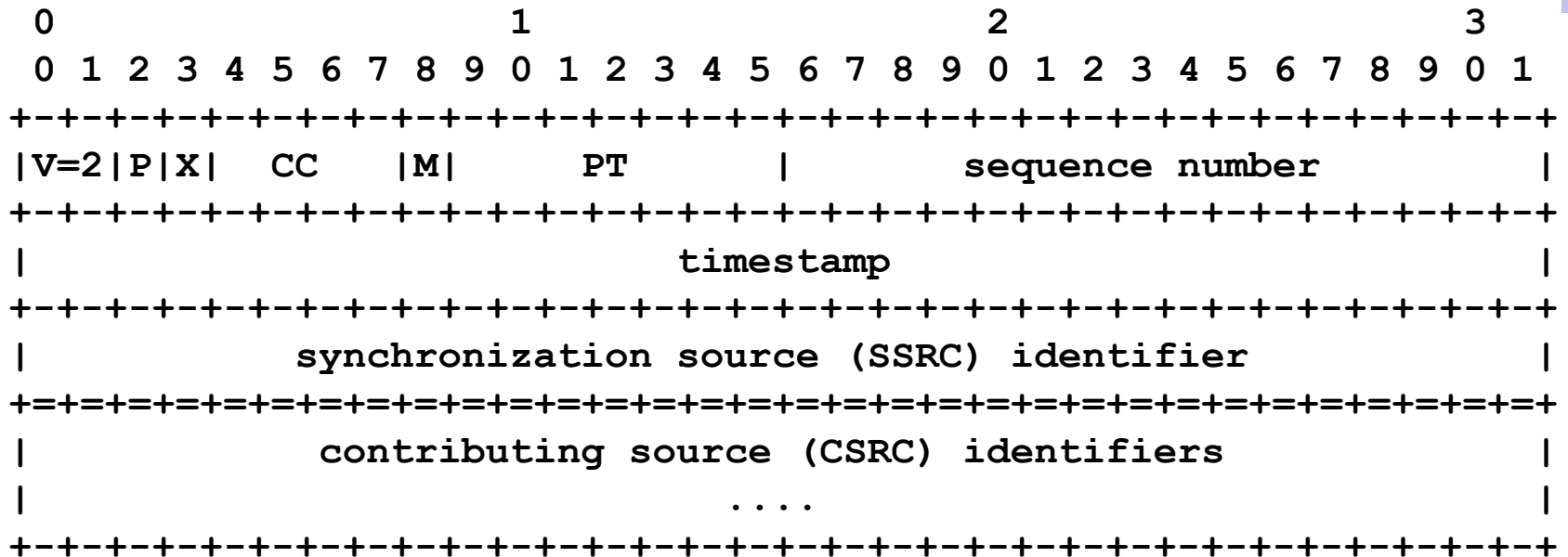


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- 2435 RTP Payload Format for JPEG-compressed Video.
- 2658 RTP Payload Format for PureVoice(tm) Audio.
- 2733 An RTP Payload Format for Generic Forward Error Correction.
- 2793 RTP Payload for Text Conversation.
- 2833 RTP Payload for DTMF Digits, Telephony Tones and Telephony...
- 3016 RTP Payload Format for MPEG-4 Audio/Visual Streams.
- 3047 RTP Payload Format for ITU-T Recommendation G.722.1.
- 3119 A More Loss-Tolerant RTP Payload Format for MP3 Audio.
- 3189 RTP Payload Format for DV (IEC 61834) Video.
- 3190 RTP Payload Format for 12-bit DAT Audio and 20- and 24-bit...
- 3389 Real-time Transport Protocol (RTP) Payload for Comfort Noise

RTP header

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- version (V)
- padding (P)
- extension (X)

- CSRC count (CC)
- marker (M)
- payload type (PT)



RTP Header

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- SSRC identifier
 - Random 32-bit value assigned by the sender
 - Per media stream
 - In multicast, used to distinguish multiple senders
 - RTCP can be used to detect colliding SSRCs
 - Also used to synchronize multi-media streams (image and sound)
 - RTCP announces when SSRCs can be combined



RTP Header



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- CSRC: Contributing Source
 - Identifies which sources were combined by a mixer
- Marker: Defined by profile.
 - For example, can indicate frame boundary
- Payload type: some well-known, some defined by profile
 - Indicates type of encoding (MPEG2, MPEG3, etc.)
- Extension: profiles can define their own extension headers



RTP Header: Sequence Number and Timestamp



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- Timestamp indicates when the media should be played back
 - Expressed in units of time defined by the profile
 - e.g., 20 ms block size of 8,000 Hz audio → 160 timestamp units per packet
 - Not absolute time, not “synchronized”
 - Rather, time since initial timestamp
 - Initial timestamp set randomly



RTP Header: Sequence Number and Timestamp

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- Sequence number used to indicate loss and ordering
- Why not use timestamp for this???



Timestamp and talk spurts



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- Receiver does not have to play out packet at exact timestamp time
- In the case of voice (with gaps in between talk spurts)
- Start of talk spurt may vary a little
 - But within a talk spurt, timing must be right
 - Think of a constant C added or subtracted from timestamp during talk spurt
- Why would we do this???



Receive buffer and jitter

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- Because of jitter, receive buffer must delay playback of voice a little
 - 10's of ms
 - More-or-less depending on RTT
 - and on amount of jitter measure over time
- Allows proper playback time even when some packets delayed



Receive buffer and jitter



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- Receiver tries to keep a certain amount of voice buffered
 - Enough to recover from jitter
 - But not so much as to introduce too much delay
- If the sender is delayed, the buffer empties a bit
- If the sender is speeded up, the buffer fills a bit
- Either way, the buffer must be brought back to the appropriate size



Receive buffer and jitter

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- The receiver can manipulate the buffer by shortening or lengthening the silences between talk spurts
 - As I said, by adding or subtracting a small constant to the timestamp
- If voice and video, must chop out some video to keep lip synch



RTCP: Real Time Control Protocol

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- Runs alongside RTP to control it in various ways
- RTP and RTCP (used to) always run on consecutive port numbers
 - But this was often screwed up by NAT, so SIP allows these numbers to be negotiated individually



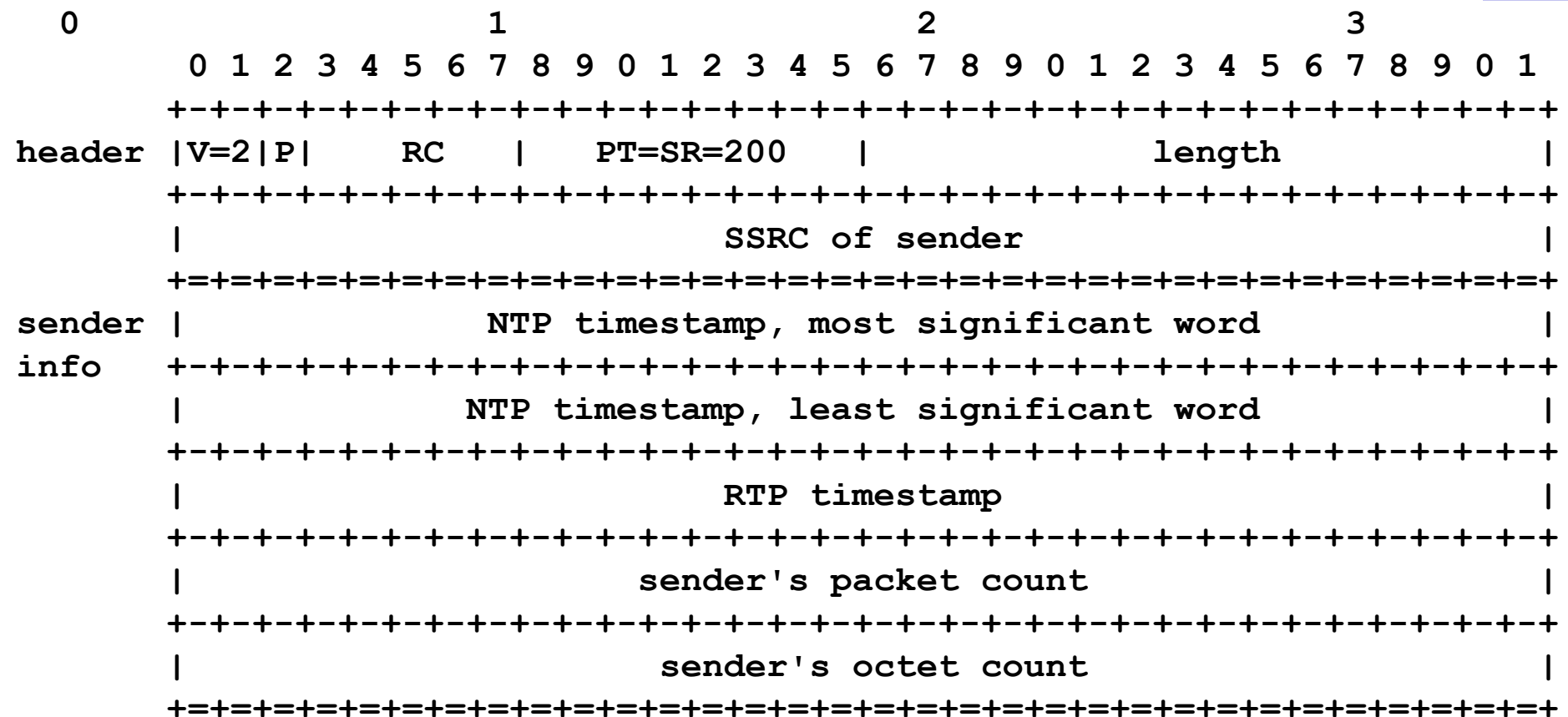
RTCP packet types

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- SR: Sender report, for transmission and reception statistics from participants that are active senders.
- RR: Receiver report, for reception statistics from participants that are not active senders.
- SDES: Source description items, including CNAME.
- BYE: Indicates end of participation.
- APP: Application specific functions.

Sender Report RTCP Packet (first part)

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		1										2										3															
		0	1	2	3	4	5	6	7	8	9	0	1	2	3	4	5	6	7	8	9	0	1	2	3	4	5	6	7	8	9	0	1				
Report Block 1		<div> <div></div> <div>SSRC_1 (SSRC of first source)</div> <div></div> </div>																																			
		<div> <div></div> <div>fraction lost</div> <div></div> </div>																																			
		<div> <div></div> <div>cumulative number of packets lost</div> <div></div> </div>																																			
		<div> <div></div> <div>extended highest sequence number received</div> <div></div> </div>																																			
		<div> <div></div> <div>interarrival jitter</div> <div></div> </div>																																			
		<div> <div></div> <div>last SR (LSR)</div> <div></div> </div>																																			
		<div> <div></div> <div>delay since last SR (DLSR)</div> <div></div> </div>																																			
		<div> <div></div> <div>SSRC_2 (SSRC of second source)</div> <div></div> </div>																																			
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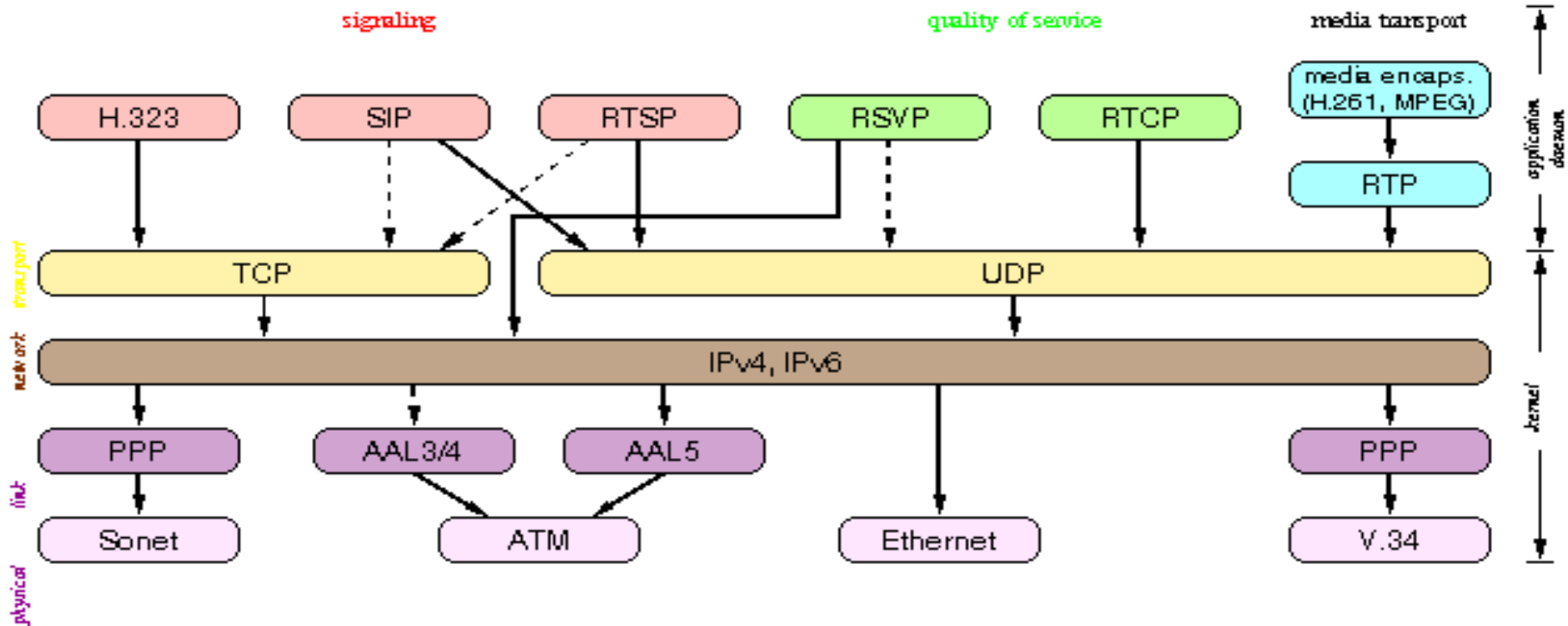
Session Initiation Protocol (SIP)

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- We've seen how RTP supports a media stream between two or more hosts
- But how did those hosts know to talk in the first place?
 - What ports to use
 - What media stream to use
 - What IP addresses to use
- SIP is one answer

Media-related protocols

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What is SIP?

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- A (formerly) lightweight signaling protocol for IP networks
 - Allows two or more hosts to tell each other what they want to do
 - Way more powerful than simple “ports”, which require a pre-established understanding
- Required for audio/video over IP
 - Because there are many types of audio/video
 - Originally a simple, multicast-aware alternative to H.323
- But has broad applicability
 - Messaging, presence, TCP, etc.



Capabilities of SIP



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- Addressing
 - Addresses users or machines
 - [user@domain](#), or +1-234-567-8901
- User location discovery
 - Through registration
- Routing
 - SIP server discovery, redirection
- Signaling
 - Negotiate services, media type, IP type (unicast or multicast), etc.
- Presence and (instant) messaging
 - As SIP “event package” (I.e. application)



Capabilities of SIP

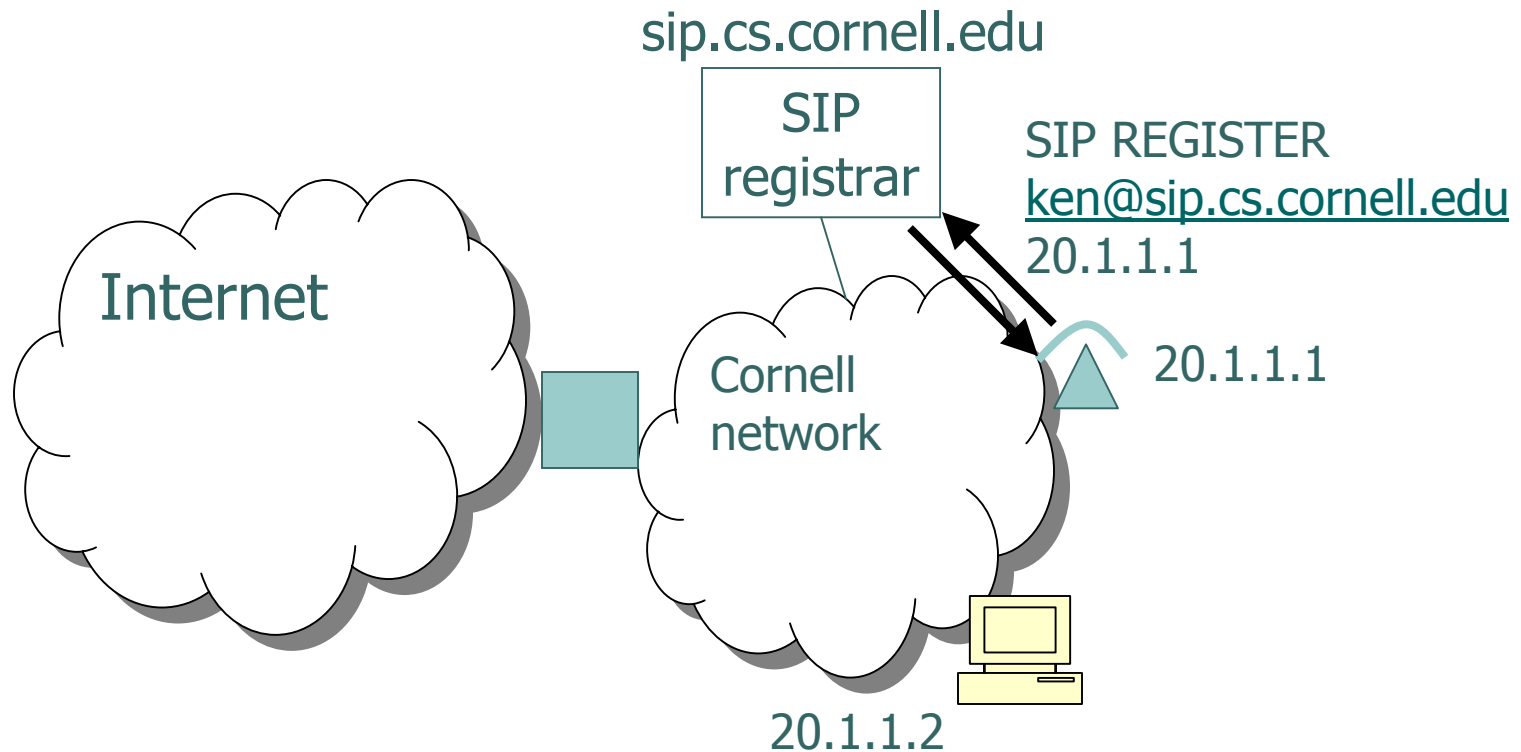


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- Secure signaling
 - Over TLS
 - Of course, can signal a secure media session, i.e. Secure RTP
- Mobility
 - Of machines across IP (re-INVITE)
 - Of users across machines (REGISTER)
- Service selection
 - Voice, email, fax, messaging, etc.
- “Call” (session) handling
 - Call forward, call transfer, 3rd party conferencing
- Interface with phone network
- NAT traversal (using STUN)

Basic SIP operation

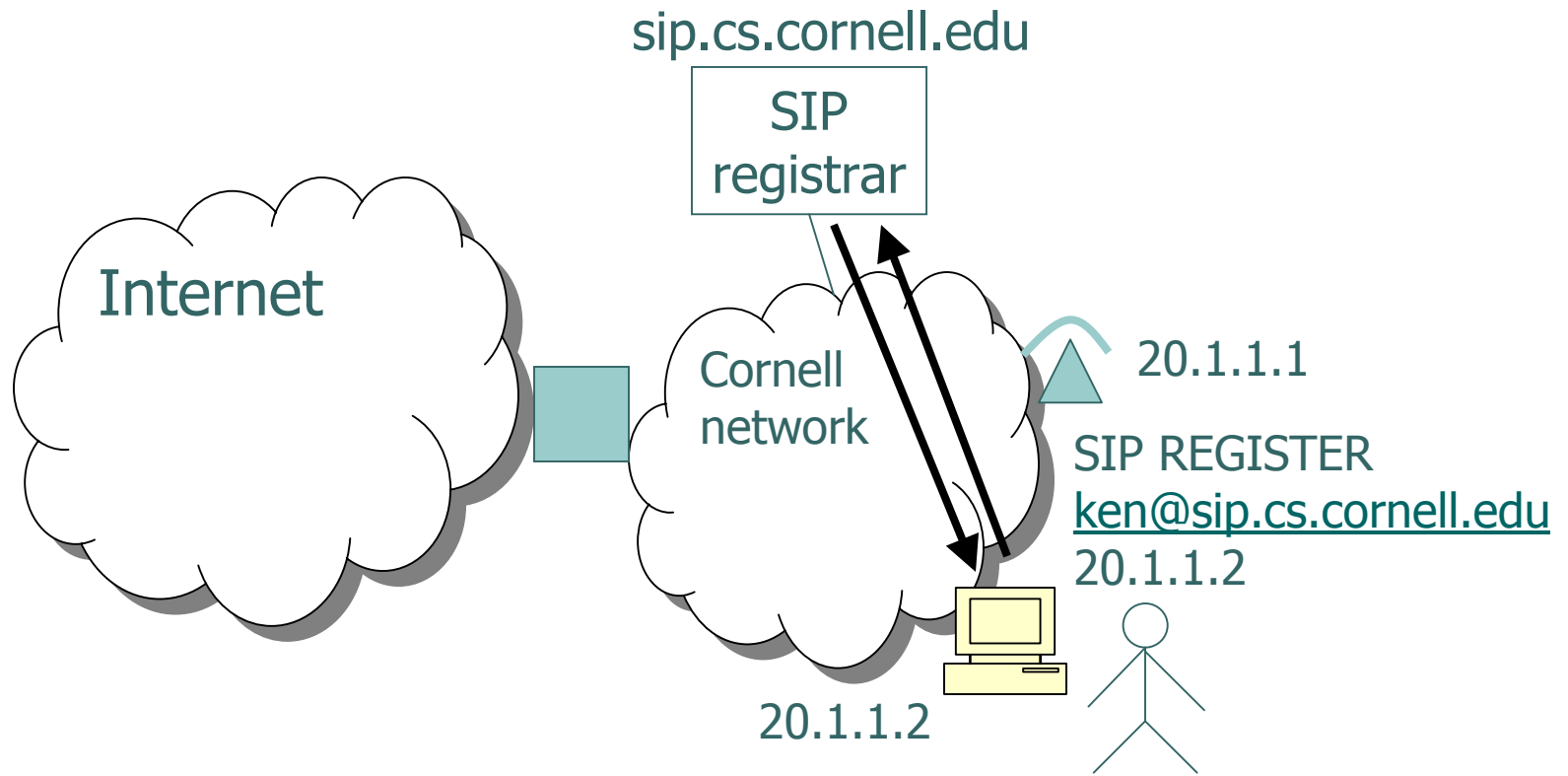
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Ken's VoIP desk phone periodically registers ken@sip.cs.cornell.edu

Basic SIP operation

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Ken moves to a computer down the hall, start VoIP app

Basic SIP operation

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INVITE
ken@sip.cs.cornell.edu,
ken-dog@sip.verizon.com,
30.1.1.1:4567, codec...

sip.cs.cornell.edu

SIP
registrar

Internet

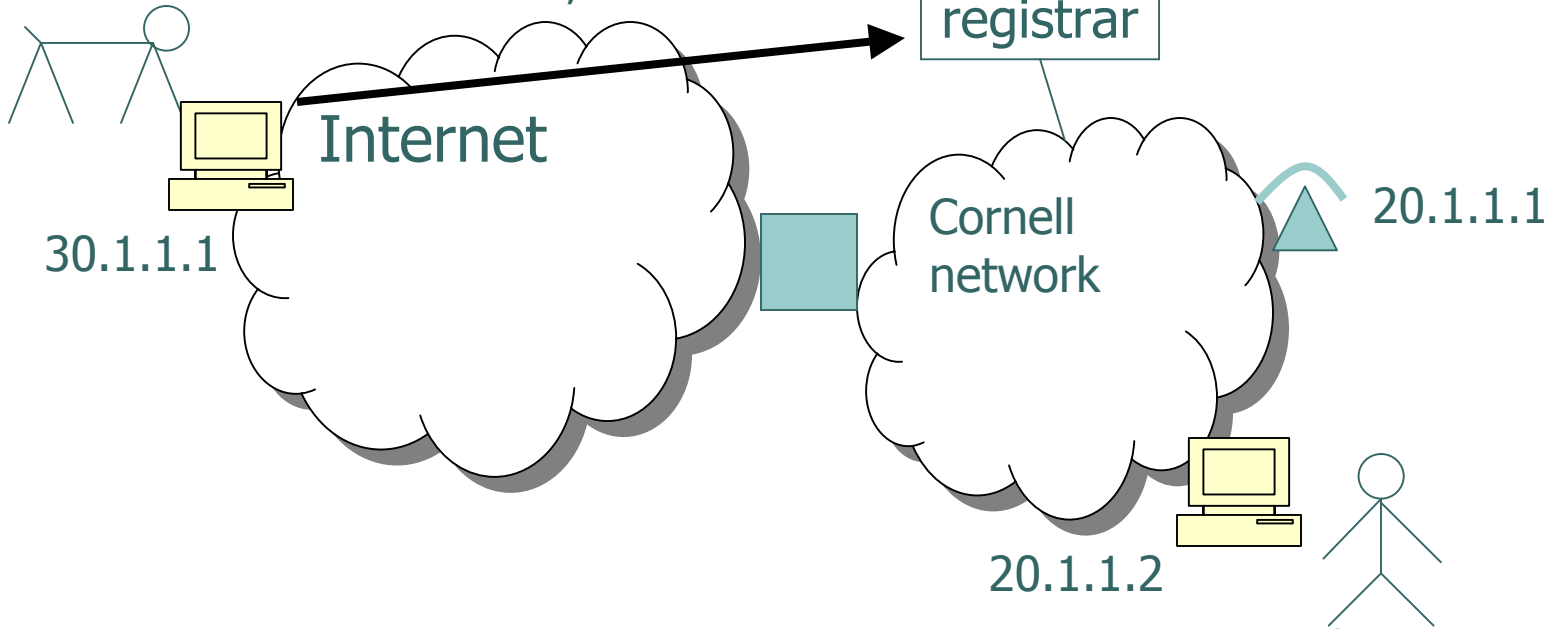
30.1.1.1

Cornell
network

20.1.1.1

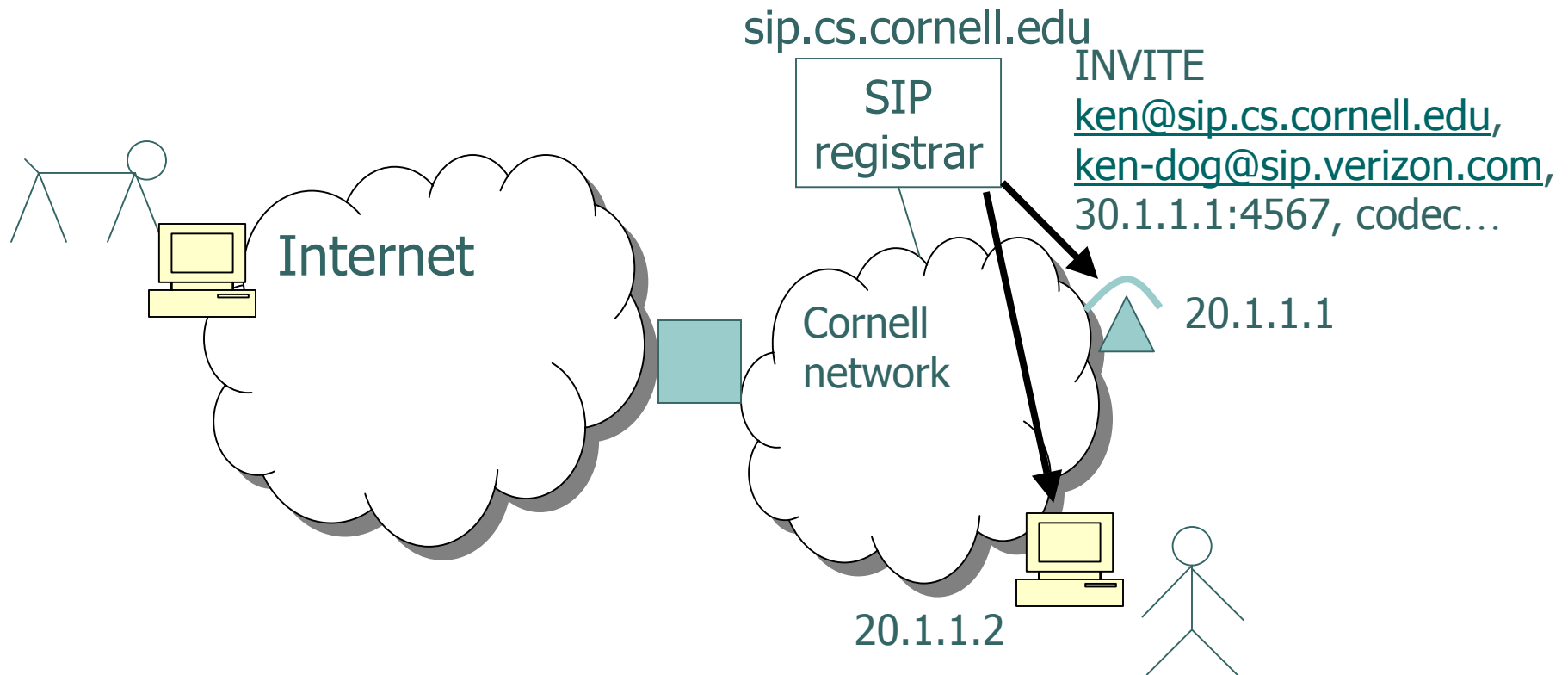
20.1.1.2

Ken's dog wants to go for a walk,
activates its BoIP phone



Basic SIP operation

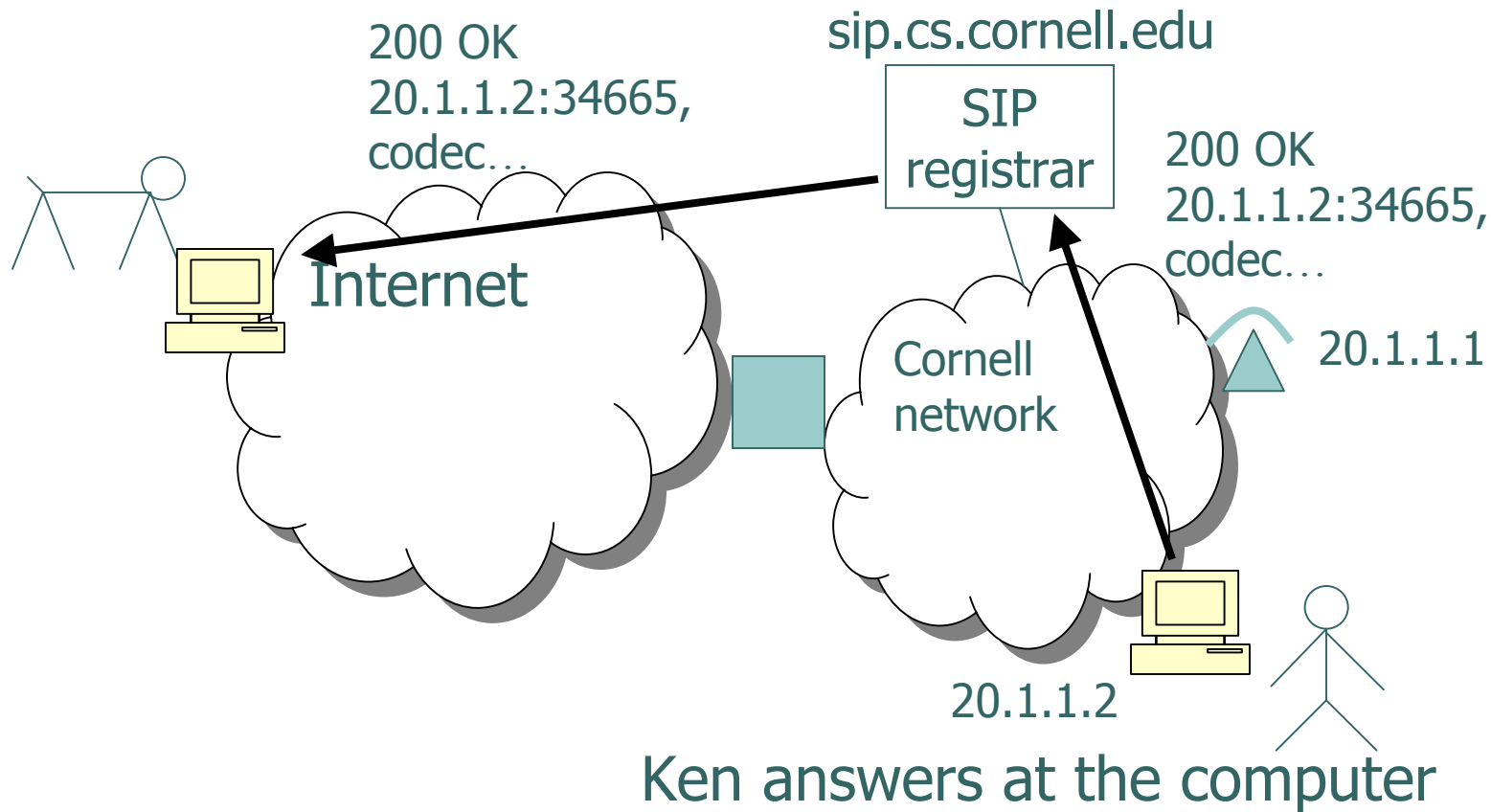
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The SIP registrar forks the INVITE, sends it to both devices

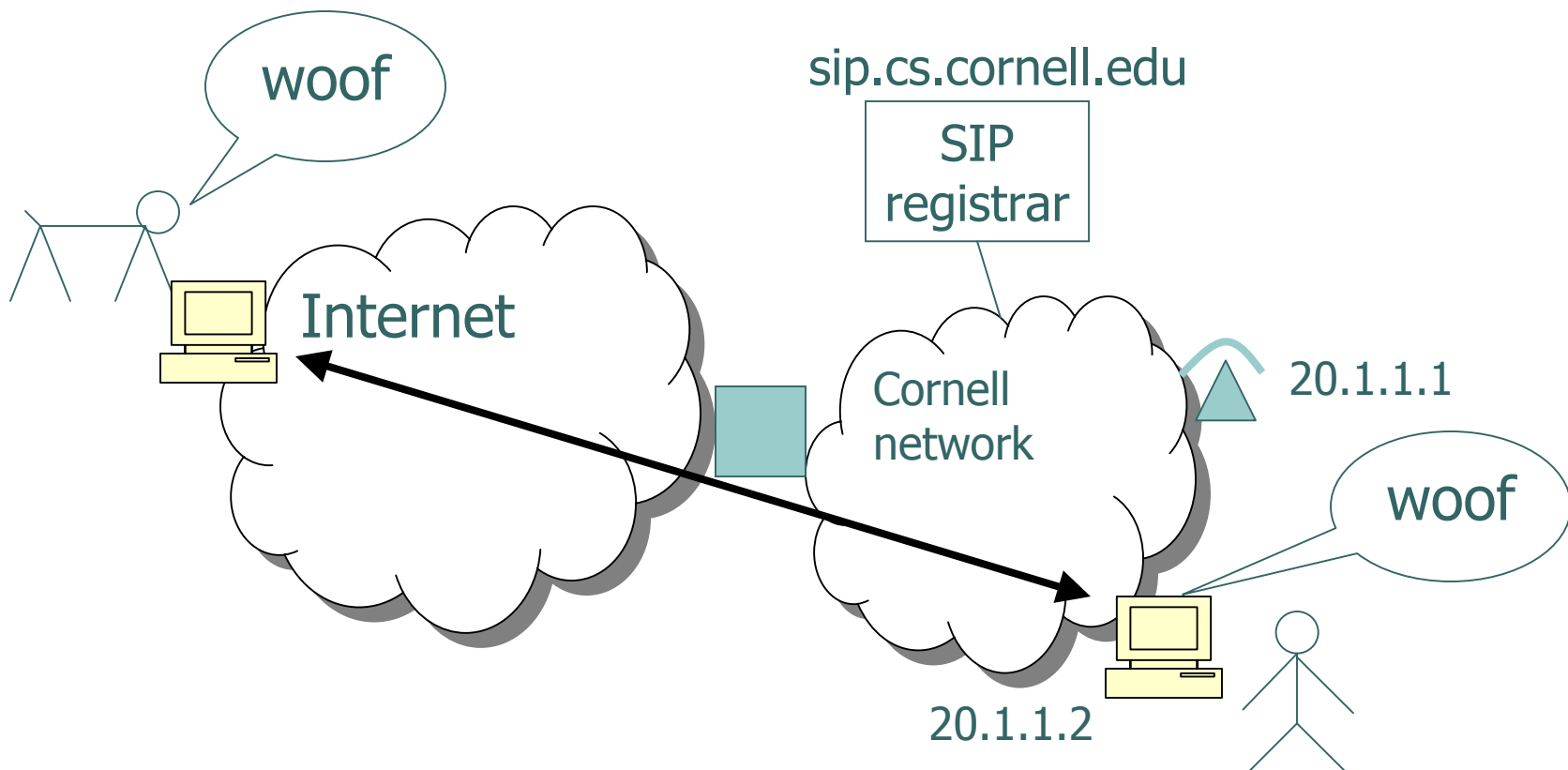
Basic SIP operation

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Basic SIP operation

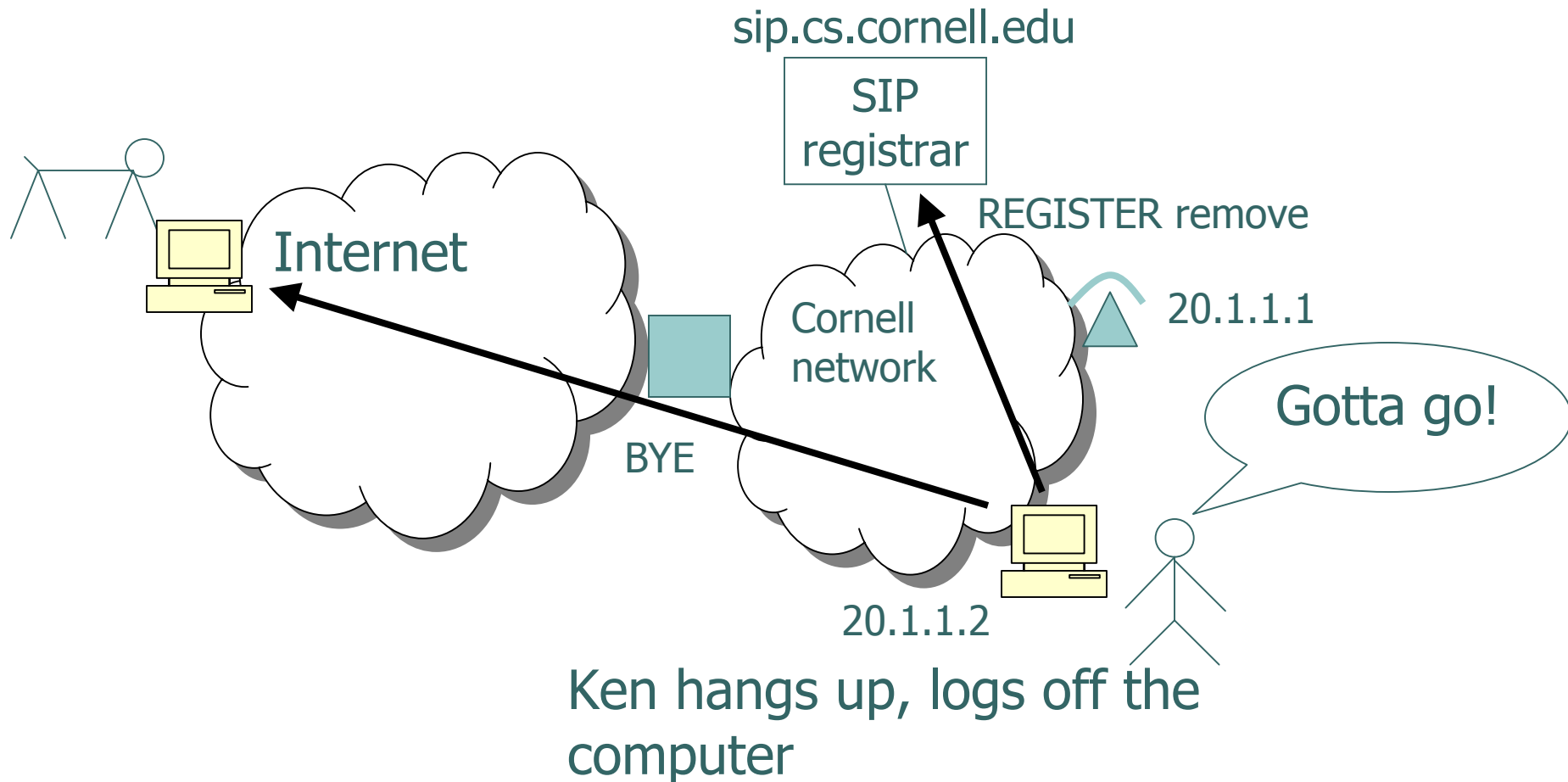
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The two devices establish a media stream over RTP

Basic SIP operation

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SIP methods

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- SIP base methods
 - REGISTER, INVITE, ACK, CANCEL, BYE, OPTIONS
- SIMPLE presence methods
 - SUBSCRIBE, NOTIFY
- SIMPLE message method
 - MESSAGE



SIP status

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- Hasn't reached “critical mass” yet
 - Though used in growing number of enterprises for voice (PBX replacement)
- Microsoft moving to SIP
 - Messenger based on SIMPLE
 - VoIP based on SIP
- Unlike IPv6, SIP doesn't have the vicious circle
 - No ISP involvement needed
 - Microsoft can bootstrap SIP all by itself



SIP future

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- Once SIP takes off, every P2P application will be built over it
 - Games, voice, video, chat, voice chat, presence, messaging, file sharing, etc.
 - Because it scales, has security, and allows easier integration of multiple communications channels
 - Example: A web-based help desk will be able to determine what applications you have (through presence, once you approve), and send you web pages, videos, etc., as part of the help service