

TOPICS

1. Smullyan's completeness proof for 1st Order tableaux rules. Please read pages 57-61.

2. Extending completeness to Refinement Logic.
The rules without realizers and a sequent rule.

$\&$ $\frac{F(X \& Y)}{FX \mid FY}$	$H \vdash X \& Y$ $H \vdash X$ $H \vdash Y$	$\frac{T(X \& Y)}{TX \mid TY}$	$H, X \& Y, H' \vdash G$ $H, x:X, y:Y, H' \vdash G$
\vee $\frac{F(X \vee Y)}{FX \mid FY}$	$H \vdash X \vee Y$ $H \vdash X$ \dots $H \vdash X \vee Y$ $H \vdash Y$	$\frac{T(X \vee Y)}{TX \mid TY}$	$H, d:X \vee Y, H' \vdash G$ $H, x:X, H' \vdash G$ $H, y:Y, H' \vdash G$
\Rightarrow $\frac{F(X \Rightarrow Y)}{TX \mid FY}$	$H \vdash X \Rightarrow Y$ $H, x:X \vdash Y$	$\frac{T(X \Rightarrow Y)}{FX \mid TY}$	$H, f:X \Rightarrow Y, H' \vdash G$ $" " " " \vdash X$ $" " " " , g:Y H' \vdash G$
\sim $\frac{F \sim X}{TX}$	$H \vdash X \Rightarrow \perp$ $H, x:X \vdash \perp$	$\frac{T(\sim X)}{FX}$	$H, f:X \Rightarrow \perp, H' \vdash G$ $H, " " " , H' \vdash X$ $H, " " " , g:\perp H' \vdash G \text{ aux(3)}$

Example

$$\underline{F(X \Rightarrow X)}$$

$$\begin{array}{l} TX \\ FX \\ \times \end{array}$$

$$H \vdash X \Rightarrow X$$

$$H, x: X \vdash X$$

Tableaux proof of $X \vee \neg X$

$$\frac{F(X \vee (X \Rightarrow \perp))}{\begin{array}{c} FX, \underline{F(X \Rightarrow \perp)} \\ \begin{array}{c} TX \\ F\perp \\ \times \\ FX, TX \end{array} \end{array}}$$

For refinement logic we
need a new rule

$$H \vdash A \vee (A \Rightarrow \perp)$$

by magic(A)

Classical logic assumes a non-empty domain D , so
in refinement logic we need

$$H \vdash D \text{ by } d_0 \quad \text{or} \quad H \vdash D \text{ by } \text{obj}(d_0)$$

In refinement logic we can add a sequentias rule
(also called a cut rule).

$$\frac{\begin{array}{c} H \vdash G \text{ by } \text{seq } T \quad \text{ap}(\lambda(x. \overline{p}) ; \overline{q}) \\ H \vdash T \text{ by } t \end{array}}{H, x:T \vdash G \text{ by } g(x)}$$

Exercise Give a tableaux proof of

$$\exists y:D. (P(y) \Rightarrow \forall x:D. P(x))$$