

19 Mar 2025

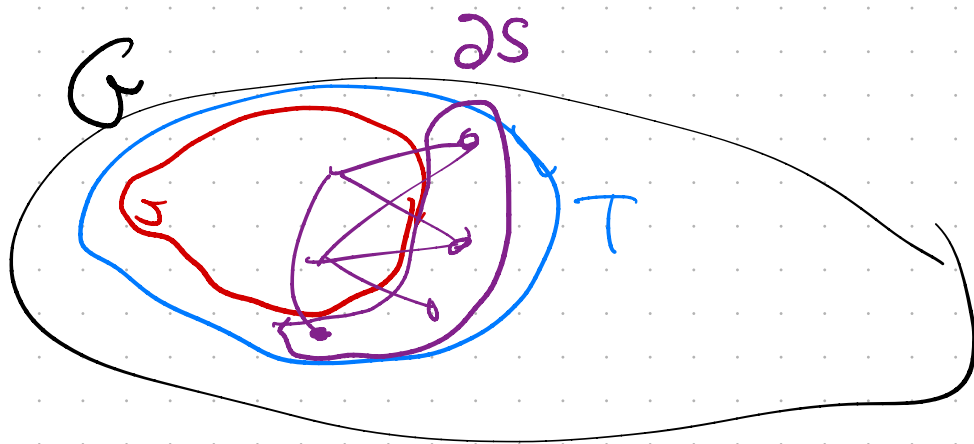
Finishing expansion

Introducing Ramsey Theory

Theorem If $p \geq \frac{7 \ln(n)}{n}$ then $G(n, p)$ is a $\frac{1}{2}$ -expander with probability $\rightarrow 1$ as $n \rightarrow \infty$.

Proof. Call a pair of vertex sets S, T an $\frac{1}{2}$ -expansion refuter ($\frac{1}{2}$ -ER) if

- $S \subset T$
- $|S| > \frac{2}{3}|T|$, $|S| \leq \frac{n}{2}$
- All elements of ∂S belong to T .



G is not a $\frac{1}{2}$ -expander



G has a $\frac{1}{2}$ -ER.

$$(\Downarrow) \exists S \quad |S| \leq \frac{n}{2}, \quad |\partial S| < \frac{1}{2}|S|$$

$$\Rightarrow T = S \cup \partial S, \quad |S| > \frac{2}{3}|T|$$

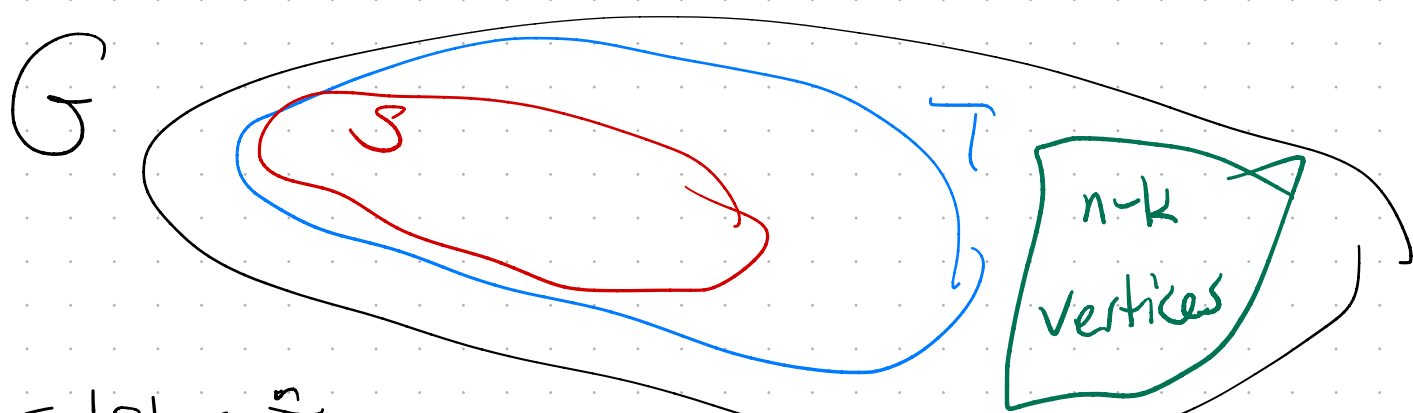
so S, T is a $\frac{1}{2}$ -ER.

(\uparrow) If (S, T) is a $\frac{1}{2}$ -ER then
 $|S| \leq \frac{n}{2}$ (2nd line of def.)

$\partial S \subseteq T \setminus S$ (3rd line of def.)

$|\partial S| < \frac{1}{2} |S|$ (2nd line of def.)

$T \setminus S$ has $< \frac{1}{2}$ as many elements as S .



$$\frac{2}{3}k < |S| \leq \frac{n}{2} \Rightarrow k < \frac{3}{4}n \Rightarrow n-k > \frac{1}{4}n.$$

$$\mathbb{E} \left[\# \text{ of } \frac{1}{2}\text{-ER in } G(n, p) \right]$$

$$= \sum_T \sum_{S \subseteq T} \mathbb{P} \left((S, T) \text{ is a } \frac{1}{2}\text{-ER} \right)$$

$$\leq \sum_{k=1}^n \binom{n}{k} \cdot 2^k \cdot (1-p)^{\binom{\frac{2}{3}k}{2} \binom{\frac{1}{4}n}{2}}$$

$$= \sum_{k=1}^n \binom{n}{k} \left[2 (1-p)^{n/6} \right]^k$$

$$= -1 + \left(1 + 2 (1-p)^{n/6} \right)^n$$

$$\leq 1 + [1 + 2e^{-pn/6}]^n$$

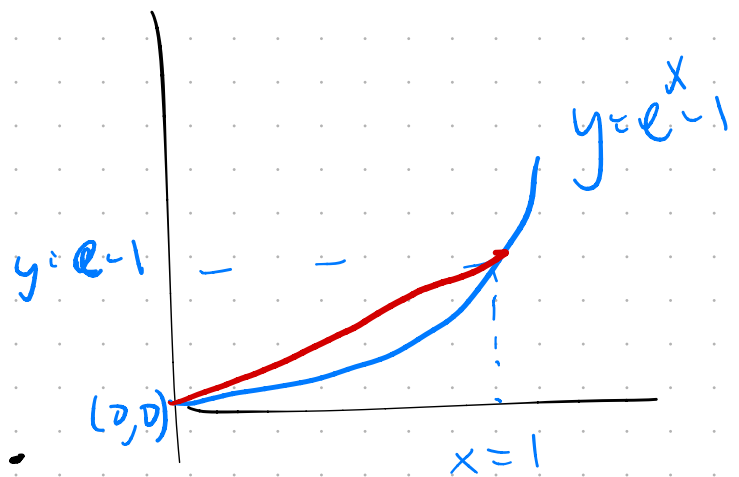
$$p \geq \frac{7 \ln n}{n}$$

$$\frac{p}{6} \geq \frac{7}{6} \ln(n)$$

$$\leq 1 + (1 + 2n^{-7/6})^n$$

$$\leq 1 + (e^{2n^{-7/6}})^n$$

$$\stackrel{[n \geq 64]}{\leq} 1 + e^{2n^{-1/6}} \leq (e-1) \cdot \frac{2}{n^{1/6}}$$



By Markov Ineq,

$$\Pr(G(n,p) \text{ has a } \frac{1}{2}\text{-ER}) \leq \frac{2(e-1)}{n^{1/6}}$$

$\Pr(G(n,p) \text{ is a } \frac{1}{2}\text{-expander})$

$$\geq 1 - \frac{2(e-1)}{n^{1/6}},$$

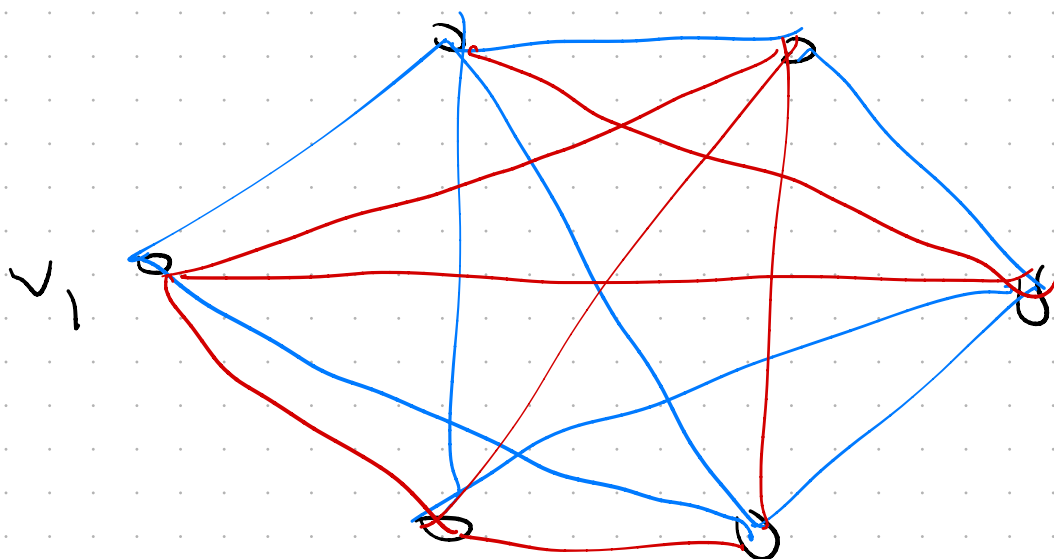
RAMSEY THEORY

Def. A clique in an undirected graph is a vertex set S such that every two elements of S form an edge.

An indep set in an undirected graph is a vertex set S such that every two elements of S don't form an edge.

Theorem. Every graph with ≥ 6 vertices contains either a clique or an independent set of size 3.

Proof. Take any 6 vertices $v_1, v_2, v_3, v_4, v_5, v_6$.



Connect every pair with a
blue edge if $(v_i, v_j) \in E(G)$
red edge otherwise.

v_i either has 3 blue neighbors
or 3 red neighbors.

Assume red. If any 2 of
those are connected by a
red edge, we have a
red triangle.

If they are all connected by
blue edges, we have a
blue triangle.

$$\underline{R(3,3) \leq 6.}$$

$R(k, l)$ is the
minimum n such that every n -vertex
graph has either a clique of size k
or indep set of size l .

Theorem [Ramsey, 1930]

$$R(k, l) \leq 2^{k+l-3},$$

Proof. Suppose $n = 2^{k+l-3}$.

Color every pair of vertices
blue if (v_i, v_j) is edge of G
red if not.

We seek a clique with all edges
same color, either blue with
 k vertices or red with
 l vertices. $t = k+l-3$.

Construct sequence of vertices

v_0, v_1, \dots, v_t

and vertex sets $S_0 \supset S_1 \supset \dots \supset S_t$

Starting with $v_0 =$ any vertex
 $S_0 =$ all vertices.

Always $v_i \in S_i$.

Look at colors of edges

from v_i to $S_i \setminus \{v_i\}$.

Majority red $\Rightarrow S_{i+1}$ = red neighbors

Majority blue $\Rightarrow S_{i+1}$ = blue neighbors

v_{i+1} = arbitrary element of S_{i+1} .

$$|S_{i+1}| \geq \left\lceil \frac{|S_i| - 1}{2} \right\rceil.$$

$$|S_0| \geq 2^{k+l-3}$$

$\Rightarrow S_1, \dots, S_t$ all non-empty.

v_0 , v_1 , v_2 ... , v_{t-1} , v_t

If $k-1$ of the vertices in v_0, \dots, v_{t-1} are blue \Rightarrow blue k -clique.

If $l-1$ of them are red
 \Rightarrow red l -clique.

If $\leq k-2$ blue, $\leq l-2$ red
 $\Rightarrow \leq k+l-4$ colored vertices in
 v_0, \dots, v_{t-1} \searrow

Thm (Erdős, 1947)

$$R(k, k) \geq 2^{\frac{k+1}{2}} \quad \forall k \geq 3.$$

Proof (next time).