

For each of questions A and B, match the grammars on the left with the sets they generate on the right. The correspondence is one-to-one. The start symbol in all cases is  $S$ .

A.

- |  |          |  |
|--|----------|--|
| 1. $S \rightarrow aSbb \mid \varepsilon$                               | <b>d</b> | a. $\{x \in \{a, b\}^* \mid x = \text{rev } x\}$ |
| 2. $S \rightarrow aaSb \mid \varepsilon$                               | <b>e</b> | b. $\{a, b\}^*$                                  |
| 3. $S \rightarrow aSb \mid bSa \mid SS \mid \varepsilon$               | <b>c</b> | c. $\{x \in \{a, b\}^* \mid \#a(x) = \#b(x)\}$   |
| 4. $S \rightarrow aSa \mid bSb \mid a \mid b \mid \varepsilon$         | <b>a</b> | d. $\{a^n b^m \mid m = 2n\}$                     |
| 5. $S \rightarrow aSb \mid bSa \mid a \mid b \mid SS \mid \varepsilon$ | <b>b</b> | e. $\{a^n b^m \mid n = 2m\}$                     |

B.

- |                               |                                      |                                      |   |   |
|-------------------------------|--------------------------------------|--------------------------------------|---|---|
| 6. $S \rightarrow aSb \mid T$ | $T \rightarrow bTa \mid \varepsilon$ | <b>g</b>                             | f. $\{a^n b^{n+m} a^m \mid n, m \geq 0\}$ |   |
| 7. $S \rightarrow TT$         | $T \rightarrow aTb \mid \varepsilon$ | <b>h</b>                             | g. $\{a^n b^m a^m b^n \mid n, m \geq 0\}$ |   |
| 8. $S \rightarrow TU$         | $T \rightarrow aTb \mid \varepsilon$ | $U \rightarrow bUa \mid \varepsilon$ | <b>f</b>                                  | h. $\{a^n b^n a^m b^m \mid n, m \geq 0\}$ |

C. The following is a grammar in Greibach normal form for the set of balanced parentheses. The start symbol is  $S$ .

$$S \rightarrow [B$$

$$B \rightarrow ] \mid ]S \mid [BB$$

Which of the following sentential forms would *not* occur in any derivation of the string  $[[[]][[]]$ ?

- $[[[BB$
- $[[[]B$
- $[[[] [BB$
- $[[[] []B$
- $[[[] []]B$  ✓
- $[[[] []] [B$

In this grammar, the number of  $B$ 's in any sentential form generated from  $S$  is always the same as the number of unmatched left parens in the terminal string generated so far.