

Constraint Satisfaction

CS472/CS473 — Fall 2005

Slide CS472 – Constraint Satisfaction 1

Moving to a different formalism...

```

SEND
+  MORE
-----
MONEY
    
```

Consider state space for cryptarithmic (e.g. DFS).

Is this (DFS) how humans tackle the problem?

Human problem solving appears more **sophisticated!**

For example, we derive new constraints on the fly.

→ **little** or **no** search!

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Constraint Satisfaction Problems (CSP)

A powerful representation for (discrete) search problems.

A **Constraint Satisfaction Problem (CSP)** is defined by:

X is a set of n variables X_1, X_2, \dots, X_n ,
each defined by a finite domain D_1, D_2, \dots, D_n
of possible values.

C is a set of constraints C_1, C_2, \dots, C_m .

Each C_i involves a subset of the variables; specifies
the allowable combinations of values for that subset.

A **solution** is an assignment of values to the variables
that satisfies all constraints.

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Cryptarithmic as a CSP

Problem: TWO + TWO = FOUR

Variables:

$T \in \{0, \dots, 9\}; W \in \{0, \dots, 9\}; O \in \{0, \dots, 9\};$

$F \in \{0, \dots, 9\}; U \in \{0, \dots, 9\}; R \in \{0, \dots, 9\};$

$X_1 \in \{0, \dots, 1\}; X_2 \in \{0, \dots, 1\}; X_3 \in \{0, \dots, 1\};$

Constraints:

$O + O = R + 10 * X_1$

$X_1 + W + W = U + 10 * X_2$

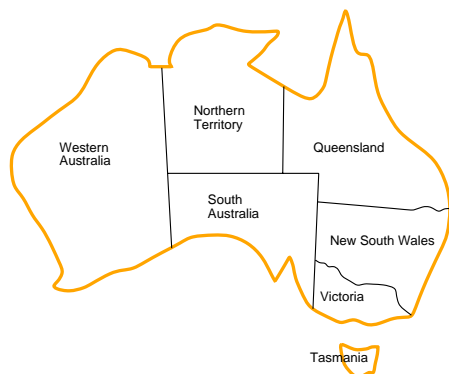
$X_2 + T + T = O + 10 * X_3$

$X_3 = F$

each letter has a different digit ($F \neq T, F \neq U$, etc);

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Map-Coloring Problem



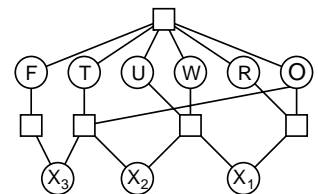
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Cryptarithmic Constraint Graph

```

  T W O
+ T W O
-----
F O U R
    
```

(a)



(b)

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Types of Constraints

Unary Constraints: Restriction on single variable

Binary Constraints: Restriction on pairs of variables

Higher-Order Constraints: Restriction on more than two variables

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Constraint Satisfaction Problems (CSP)

For a given CSP the problem is one of the following:

1. find all solutions
2. find one solution
 - just a feasible solution, or
 - a “reasonably good” feasible solution, or
 - the optimal solution given an objective function
3. determine if a solution exists

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How to View a CSP as a Search Problem?

Initial State – state in which all the variables are unassigned.

Successor function – assign a value to a variable from a set of possible values.

Goal test – check if all the variables are assigned and all the constraints are satisfied.

Path cost – assumes constant cost for each step

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Branching Factor

Approach 1 – any unassigned variable at a given state can be assigned a value by an operator: branching factor as high as sum of size of all domains.

Approach 2 – since order of variable assignment not relevant, consider as the successors of a node just the different values of a *single* unassigned variable: max branching factor = max size of domain.

Maximum Depth of Search Tree

n the number of variables \rightarrow all solutions at depth n .
Prefer DFS or BFS?

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CSP – Goal Decomposed into Constraints

Backtracking Search : a DFS that

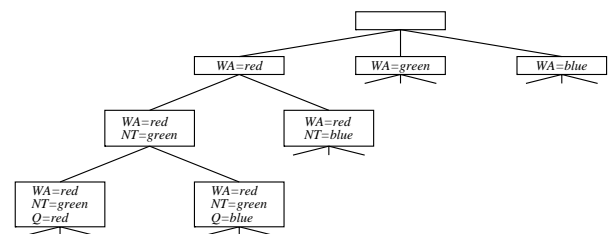
- chooses values for variables one at a time
- checks for *consistency* with the constraints.

Decisions during search :

- Which variable to choose next for assignment.
- Which value to choose next for the variable.

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Example



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General Purpose Heuristics

- **Variable and value ordering:**

Minimum remaining values (MRV): choose the variable with the *fewest* possible values.

Degree heuristic: assign a value to the variable that is involved in the largest number of constraints on other unassigned variables.

Least-constraining value heuristic: choose a value that rules out the smallest number of values in variables connected to the current variable by constraints.