

Game Playing

An AI Favorite

- structured task
- not initially thought to require large amounts of knowledge
- focus on games of perfect information

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Game Playing

Initial State is the initial board/position

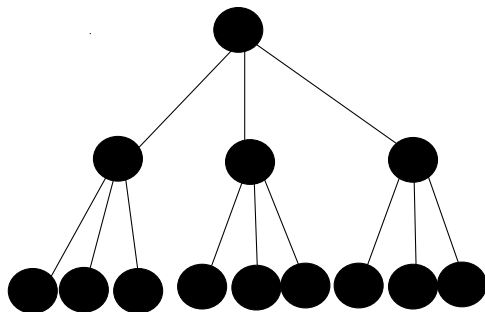
Successor Function defines the set of legal moves from any position

Terminal Test determines when the game is over

Utility Function gives a numeric outcome for the game

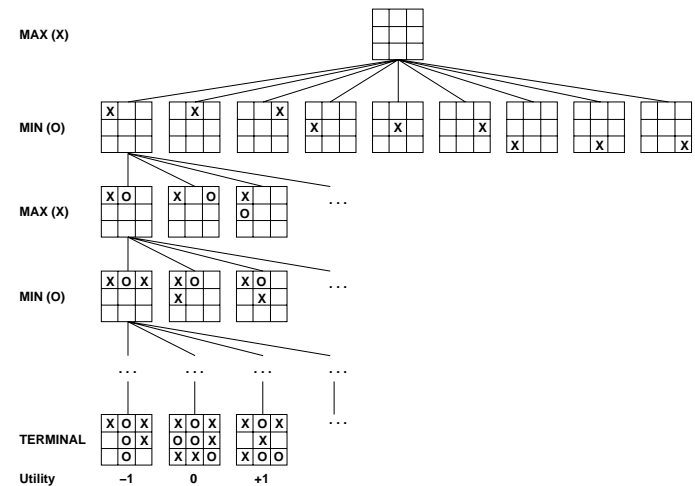
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Game Playing as Search



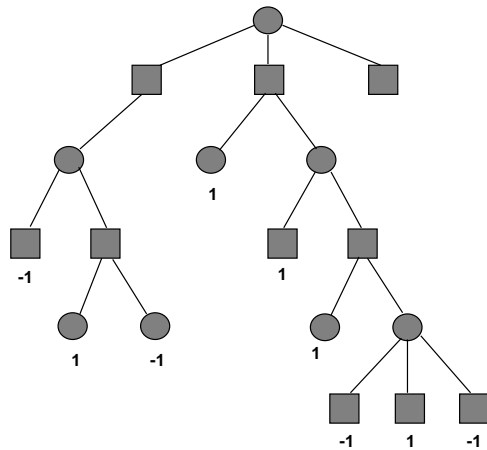
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Partial Search Tree for Tic-Tac-Toe



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Simple Minimax



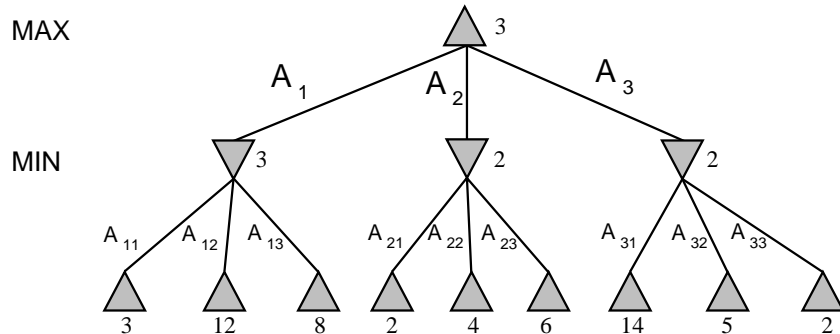
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Simplified Minimax Algorithm

1. Expand the entire tree below the root.
2. Evaluate the terminal nodes as wins for the minimizer or maximizer.
3. Select an unlabeled node, n , all of whose children have been assigned values. If there is no such node, we're done — return the value assigned to the root.
4. If n is a minimizer move, assign it a value that is the minimum of the values of its children. If n is a maximizer move, assign it a value that is the maximum of the values of its children. Return to Step 3.

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Another Example



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Minimax

```

function MINIMAX-DECISION(game) returns an operator
  for each op in OPERATORS[game] do
    VALUE[op] ← MINIMAX-VALUE(APPLY(op, game), game)
  end
  return the op with the highest VALUE[op]
  
```

```

function MINIMAX-VALUE(state, game) returns a utility value
  if TERMINAL-TEST[game](state) then
    return UTILITY[game](state)
  else if MAX is to move in state then
    return the highest MINIMAX-VALUE of SUCCESSORS(state)
  else
    return the lowest MINIMAX-VALUE of SUCCESSORS(state)
  
```

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The Need for Imperfect Decisions

Problem: Minimax assumes the program has time to search to the terminal nodes.

Solution: Cut off search earlier and apply a heuristic evaluation function to the leaves.

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Static Evaluation Functions

Minimax depends on the translation of board quality into a single, summarizing number. Difficult. Expensive.

- Add up values of pieces each player has (weighted by importance of piece).
- Isolated pawns are bad.
- How well protected is your king?
- How much maneuverability to you have?
- Do you control the center of the board?
- Strategies change as the game proceeds.

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Design Issues for Heuristic Minimax

Evaluation Function: What features should we evaluate and how should we use them? An evaluation function should:

- 1.
- 2.
- 3.

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Linear Evaluation Functions

- $w_1f_1 + w_2f_2 + \dots + w_nf_n$
- This is what most game playing programs use
- Steps in designing an evaluation function:
 1. Pick informative features
 2. Find the weights that make the program play well

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Design Issues for Heuristic Minimax

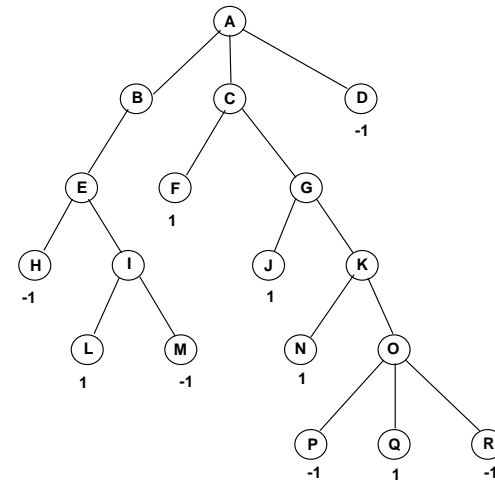
Search: search to a constant depth

Problems:

-
-

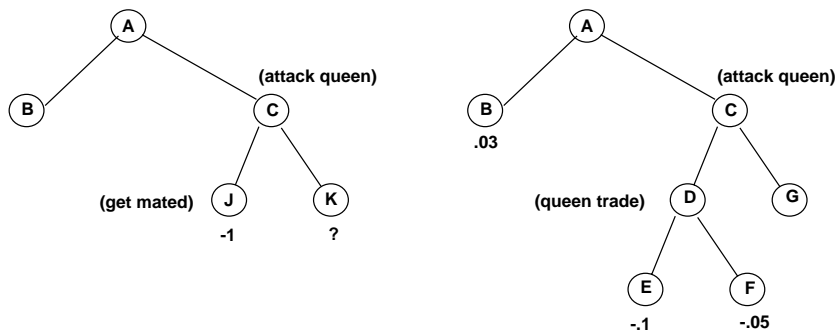
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Improving Minimax — $\alpha - \beta$ pruning



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Two More Examples



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Algebraic Solution

Let $g' = e(g)$. Then $c' = \min(-.05, g')$.

The value assigned to the root node a is

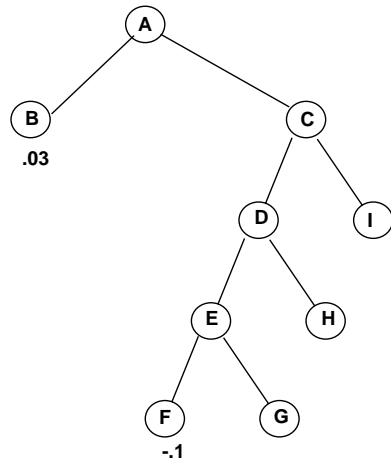
$$a' = \max(.03, \min(-.05, g')) = .03$$

because $\min(-.05, g') \leq -.05 < .03$.

The value assigned to a is independent of the value assigned to g .

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A deep $\alpha - \beta$ cutoff



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Player

Opponent

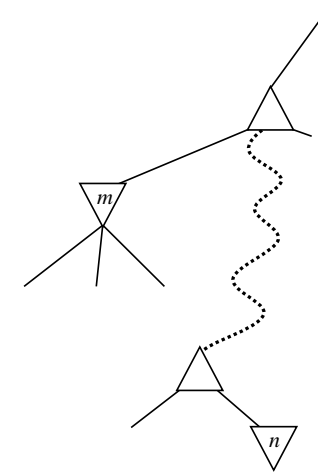
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Player

Opponent



If m is better than n for Player, never get to n in play.

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$\alpha - \beta$ Search

c = search cutoff

α = lower bound on Max's outcome; initially set to $-\infty$

β = upper bound on Min's outcome ; initially set to $+\infty$

We'll call $\alpha - \beta$ procedure recursively with a narrowing range between α and β .

Maximizing levels may reset α to a higher value; Minimizing levels may reset β to a lower value.

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$\alpha - \beta$ Search Algorithm

1. If the limit of search has been reached, compute $e(n)$ and report the result.
2. Otherwise, if the level is a **minimizing** level,
 - Until no more children or $\beta \leq \alpha$,
 - Use $\alpha - \beta$ search on child with current values of α and β ; note the value, v , returned.
 - If $v < \beta$, reset β to v .
 - Report β .

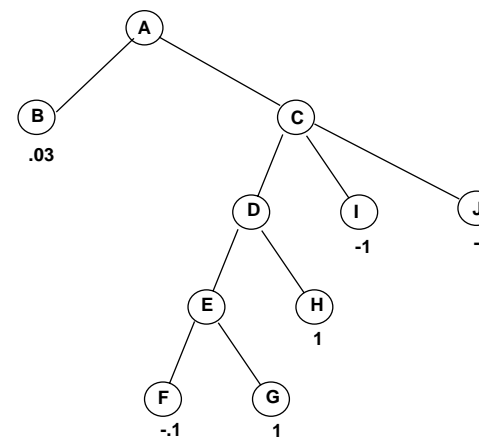
Slide CS472 – Adversarial Search 20

3. Otherwise, the level is a **maximizing** level:

- Until no more children or $\alpha \geq \beta$,
 - Use $\alpha - \beta$ search on child with current values of α and β ; note the value, v , returned.
 - If $v > \alpha$, reset α to v .
- Report α .

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Example

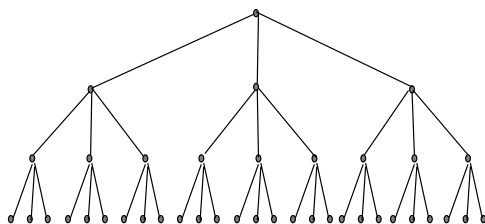


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Search Space Size Reductions

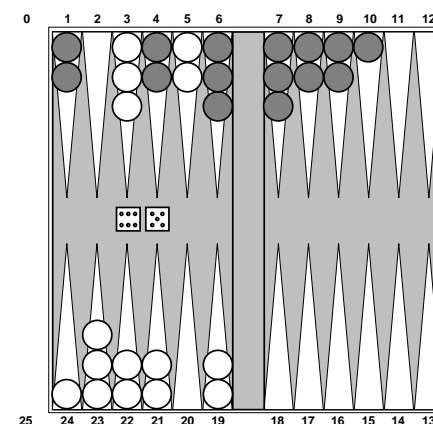
Worst Case: In an ordering where worst options evaluated first, all nodes must be examined.

Best Case: If nodes ordered so that the best options are evaluated first, then what?



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Backgammon – Board



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Backgammon – Rules

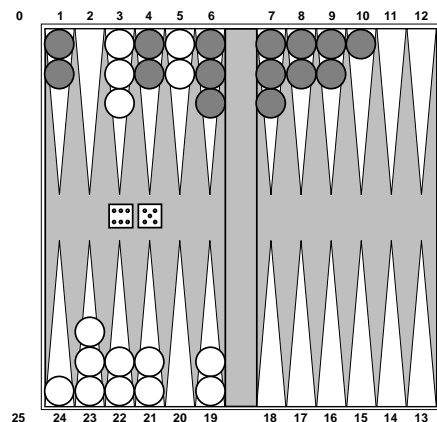
- Goal: move all of your pieces off the board before your opponent does.
- Black moves counterclockwise toward 0.
- White moves clockwise toward 25.
- A piece can move to any position except one where there are two or more of the opponent's pieces.
- If it moves to a position with one opponent piece, that piece is captured and has to start its journey from the beginning.

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Backgammon – Rules

- If you roll doubles you take 4 moves (example: roll 5,5, make moves 5,5,5,5).
- Moves can be made by one or two pieces (in the case of doubles by 1, 2, 3 or 4 pieces)
- And a few other rules that concern *bearing off* and *forced moves*.

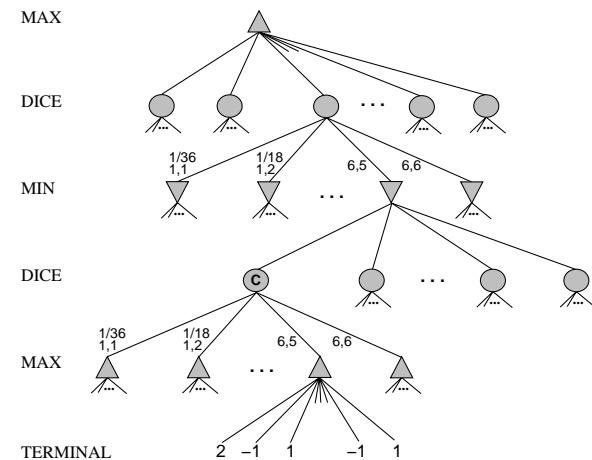
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White has rolled 6-5 and has 4 legal moves: (5-10,5-11), (5-11,19-24), (5-10,10-16) and (5-11,11-16).

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Game Tree for Backgammon



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Expectiminimax

Expectiminimax (n) =

utility(n) for n, a terminal state

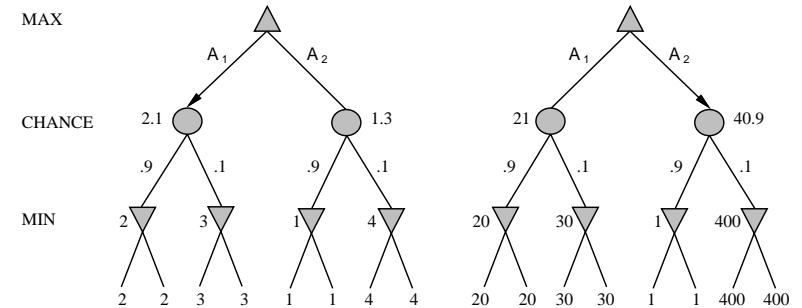
$\max_{s \in Succ(n)} \text{expectiminimax}(s)$ for n, a Max node

$\min_{s \in Succ(n)} \text{expectiminimax}(s)$ for n, a Min node

$\sum_{s \in Succ(n)} P(s) * \text{expectiminimax}(s)$ for n, a chance node

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Evaluation function



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State of the Art in Backgammon

- 1980: *BKG* using two-ply (depth 2) search and lots of luck defeated the human world champion.
- 1992: Tesauro combines Samuel's learning method with neural networks to develop a new evaluation function, resulting in a program ranked among the top 3 players in the world.

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State of the Art in Checkers

- 1952: Samuel developed a checkers program that learned its own evaluation function through self play.
- 1990: *Chinook* (J. Schaeffer) wins the U.S. Open. At the world championship, Marion Tinsley beat *Chinook*.

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State of the Art in Go

Large branching factor makes regular search methods inappropriate.

Best computer Go programs ranked only “weak amateur”.

Employ pattern recognition techniques and limited search.

\$2,000,000 prize available for first computer program to defeat a top level player.

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Othello

- Smaller search space than chess; usually 5 to 15 legal moves.
- Evaluation function expertise had to be developed from scratch.
- 1997: Logistello defeated the human world champion, 6-0.
- Generally acknowledged that humans are no match for computers at Othello.

Slide CS472 – Adversarial Search 34

History of Chess in AI

500	legal chess
1200	occasional player
2000	world-ranked
2900	Gary Kasparov

Early 1950's Shannon and Turing both had programs that (barely) played legal chess (500 rank).

1950's Alex Bernstein's system, (500+ ϵ).

1957 Herb Simon claims that a computer chess program would be world chess champion in 10 years...yeah, right.

Slide CS472 – Adversarial Search 35

1966 McCarthy arranges computer chess match, Stanford vs. Russia. Long, drawn-out match. Russia wins.

1967 Richard Greenblatt, MIT. First of the modern chess programs, *MacHack* (1100 rating).

1968 McCarthy, Michie, Papert bet Levy (rated 2325) that a computer program would beat him within 10 years.

1970 ACM started running chess tournaments. Chess 3.0-6 (rated 1400).

1973 By 1973...Slate: “It had become too painful even to look at Chess 3.6 any more, let alone work on it.”

1973 Chess 4.0: smart plausible-move generator rather than

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speeding up the search. Improved rapidly when put on faster machines.

1976 Chess 4.5: ranking of 2070.

1977 Chess 4.5 vs. Levy. Levy wins.

1980's Programs depend on search speed rather than knowledge (2300 range).

1993 DEEP THOUGHT: Sophisticated special-purpose computer; $\alpha - \beta$ search; searches 10-ply; singular extensions; rated about 2600.

1995 DEEP BLUE: searches 14-ply; considers 100–200 billion positions per move; regularly reaches depth 14;

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evaluation function has 8000+ features; singular extensions to 40-ply; opening book of 4000 positions; end-game database for 5-6 pieces.

1997 DEEP BLUE: first match won against world-champion (Kasparov).

2002 IBM declines re-match. FRITZ played world champion Vladimir Kramnik. 8 games. Ended in a draw.

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Concludes “Search”

- **Problem Solving as Search**

- **Uninformed search:** DFS / BFS / Uniform cost search
time / space complexity
size search space: up to approx. 10^{11} nodes
special case: **Constraint Satisfaction / CSPs**
generic framework: variables & constraints
backtrack search (DFS); propagation (forward-checking / arc-consistency, variable / value ordering

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- **Informed Search:** use heuristic function guide to goal

Greedy best-first search

A search / provably optimal*

Search space up to approximately 10^{25}

Local search

Greedy / Hillclimbing

Simulated annealing

Tabu search

Genetic Algorithms / Genetic Programming

search space 10^{100} to 10^{1000}

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- **Adversarial Search / Game Playing**

- minimax**

- Up to around 10^{10} nodes, 6 — 7 ply in chess.

- alpha-beta pruning**

- Up to around 10^{20} nodes, 14 ply in chess.

- provably optimal*

Slide CS472 – Adversarial Search 41

Search and AI

Why such a central role?

Basically, because lots of tasks in AI are **intractable**.

Search is “only” way to handle them.

Many applications of search, in e.g.,

Learning / Reasoning / Planning / NLU / Vision

Good thing: much recent progress (10^{30} quite feasible; sometimes up to 10^{1000}). **Qualitative difference**

from only a few years ago!

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