Moving to a different formalism...

SEND

+ MORE

MONEY

Consider search space for cryptarithmetic. DFS (depth-first search)

Is this (DFS) how humans tackle the problem?

And if not, what do humans do?

Slide CS472 – Constraint Satisfaction 1

Constraint Satisfaction Problems (CSP)

A powerful representation for (discrete) search problems.

A Constraint Satisfaction Problem (CSP) is defined by:

X is a set of n variables X_1, X_2, \ldots, X_n , each defined by a finite domain D_1, D_2, \ldots, D_n of possible values.

C is a set of constraints C_1, C_2, \ldots, C_m . Each C_i involves some subset of the variables; specifies the allowable combinations of values for that subset.

A **solution** is an assignment of values to the variables that satisfies all constraints.

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Human problem solving appears much more sophisticated!

For example, we derive new constraints on the fly. In a sense, we try to solve problems with **little** or **no** search!

In example, we can immediately derive that M=1. It then follows that S=8 or S=9. Etc. (derive more!)

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Cryptarithmetic as a CSP

TWO

+ TW

FOUR

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Variables:

$$T = \{0, \dots, 9\}; W = \{0, \dots, 9\}; O = \{0, \dots, 9\};$$

$$F = \{0, \dots, 9\}; U = \{0, \dots, 9\}; R = \{0, \dots, 9\};$$

Constraints:

$$O + O = R + 10 * X_1$$

$$X_1 + W + W = U + 10 * X_2$$

$$X_2 + T + T = O + 10 * X_3$$

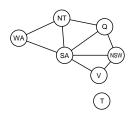
$$X_3 = F$$

$$TWO + TWO = FOUR;$$
each letter has a different digit (F \neq T, F \neq U, etc);

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Map-Coloring Problem



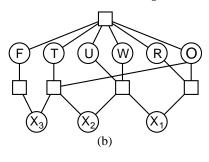


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Cryptarithmetic Constraint Graph



(a)



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Constraint Satisfaction Problems (CSP)

For a given CSP the problem is one of the following:

find all solutions

find one solution

just a feasible solution, or

a "reasonably good" feasible solution, or

the optimal solution given an objective function

determine if a solution exists

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How to View a CSP as a Search Problem?

Initial State – state in which all the variables are unassigned.

Successor function – assign a value to a variable from a set of possible values.

Goal test – check if all the variables are assigned and all the constraints are satisfied.

Path cost – assumes constant cost for each step

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CSP – Goal Decomposed into Constraints

How to exploit it?

Backtracking search: a DFS that chooses values for variables one at a time, checking for *consistency* with the constraints.

An uninformed search

Branching Factor

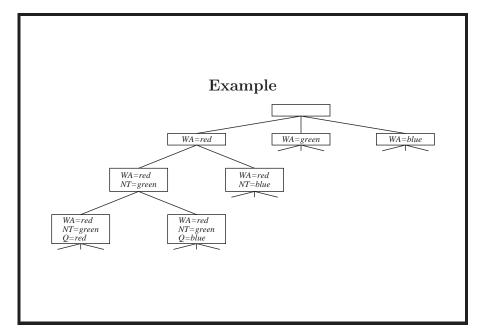
Hypothesis 1 – any unassigned variable at a given state can be assigned a value by an operator: branching factor as high as sum of size of all domains.

Better approach – since order of variable assignment not relevant, consider as the successors of a node just the different values of a *single* unassigned variable: max branching factor = max size of domain.

Maximum Depth of Search Tree

n the number of variables; all the solutions are at depth n. What are the implications in terms of using DFS vs. BFS

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Slide CS472 - Constraint Satisfaction 11

Constraint propagation: "looking ahead"

- Forward Checking each time variable is instantiated, delete from domains of the uninstantiated variables all of those values that conflict with current variable assignment.
- Arc Consistency state is arc-consistent, if every variable has some value that is consistent with each of its constraints (consider pairs of variables)
- **K-Consistency** generalizes arc-consistency. Consistency of groups of K variables.

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Constraint propagation: "looking ahead"

• Variable and value ordering:

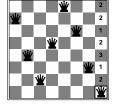
Minimum remaining values (MRV): choose the variable with the *fewest* possible values.

Degree heuristic: assign a value to the variable that is involved in the largest number of constraints on other unassigned variables.

Least-constraining value heuristic: choose a value that rules out the smallest number of values in variables connected to the current variable by constraints.

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Local Search for CSPs







Comparison of CSP Algorithms

Problem	BT	BT+MRV	BT+FC	BT+FC+	Min-Conf
USA	(>1,000K)	(>1,000K)	2K	60	64
n-queens	(>40,000K)	13,500 K	(>40,000K)	817K	4K

Dramatic recent progress in Constraint Satisfaction.

For example, methods can now handle problems with 10,000 to 100,000 variables, and up to 1,000,000 constraints.

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