*If a course is easy, some students are happy.*
Put in first-order form:

$$\forall c, easy(c) \Rightarrow \exists s, happy(s)$$

Skolemize:

$$\forall c, easy(c) \Rightarrow happy(f(c))$$

Drop universal quantifier:

$$easy(c) \Rightarrow happy(f(c))$$

Eliminate implications:

$$\neg easy(c) \vee happy(f(c))$$

*If a course has a final, no students are happy.*
Put in first-order form:

$$\forall c, final(c) \Rightarrow \neg\exists s, happy(s)$$

Move negation inward:

$$\forall c, final(c) \Rightarrow \forall s, \neg happy(s)$$

Move quantifiers outward:

$$\forall c, \forall s, final(c) \Rightarrow \neg happy(s)$$

Drop universal quantifiers:

$$final(c) \Rightarrow \neg happy(s)$$

Eliminate implications:

$$\neg final(c) \vee \neg happy(s)$$

So, we have two clauses (with variables renamed to get rid of duplicates):

$$\neg easy(c) \vee happy(f(c))$$

$$\neg final(k) \vee \neg happy(s)$$

Now, we want to show *If a course has a final, the course isn't easy.*:
Put in first-order form:

$$\forall c, final(c) \Rightarrow \neg easy(c)$$

Drop universal quantifier:

$$final(c) \Rightarrow \neg easy(c)$$

Eliminate implications:

$$\neg final(c) \vee \neg easy(c)$$

Rename variables to avoid conflicts with other clauses:

$$\neg final(j) \vee \neg easy(j)$$

For resolution, we want to negate the clause we're trying to prove:

$$final(j) \wedge easy(j)$$

Now, our complete clause set at the start of resolution looks like:
A.

$$\neg easy(c) \lor happy(f(c))$$

B.

$$\neg final(k) \lor \neg happy(s)$$

C.

$$final(j)$$

D.

$$easy(j)$$

Resolution:
Combine A and D with the resolution rule to get:
E.

$$happy(f(c))$$

Combine B and C with the resolution rule to get:
F.

$$\neg happy(s)$$

Combine E and F with the resolution rule to get:

$$\emptyset$$

Intuitively, E and F contradict eachother (E states that a particular student, $f(c)$, is happy, while F states that no student is happy), so we've derived a contradiction by assuming the negation of $\neg final(j) \lor \neg easy(j)$.