

CS4450

Computer Networks: Architecture and Protocols

Lecture 17 Putting ALL the Pieces Together

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Recap: Last lecture

- **THE Internet Protocol**

- Functionality: delivering the data
- Three key ideas:
 - **Addressing** (IP addressing)
 - **Routing** (using a variety of protocols)
 - **Packet header as an interface** (Encapsulating data into packets)
- Why do packet headers look like the way they look?

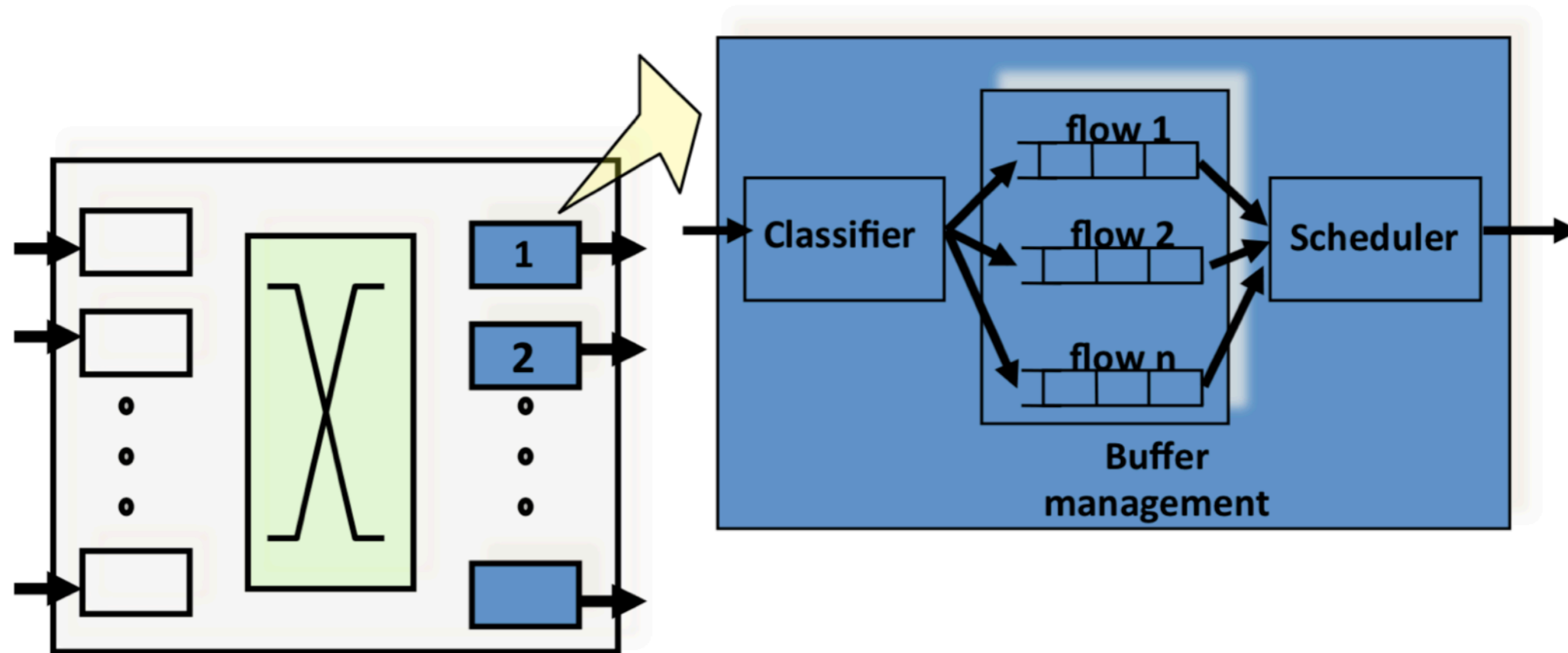
- **Switch and Router Architecture**

- Main challenge: processing packets **quickly**
- **Longest Prefix match**

Recap: Input Linecards

- Main challenge is processing speed
- Tasks involved
 - Update packet header (easy)
 - Longest prefix match lookup on destinations address (harder)
- Mostly implemented with specialized hardware

Output Linecard



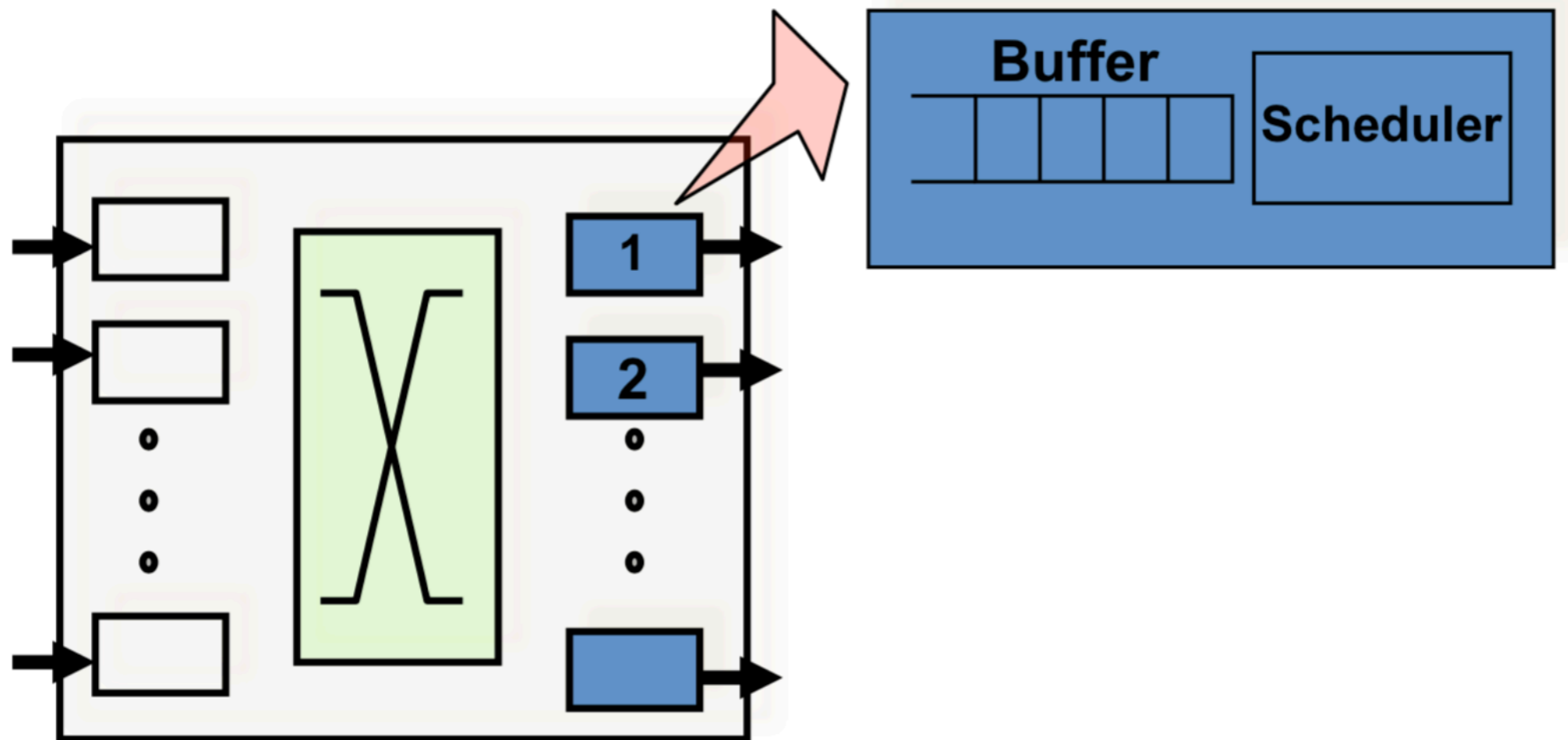
- **Packet Classification:** map each packet to a “flow”
 - Flow (for now): set of packets between two particular endpoints
- **Buffer Management:** decide when and which packet to drop
- **Scheduler:** decide when and which packet to transmit

Output Linecard

- **Packet Classification:** map each packet to a “flow”
 - Flow (for now): set of packets between two particular endpoints
- **Buffer Management:** decide when and which packet to drop
- **Scheduler:** decide when and which packet to transmit
- Used to implement various forms of policy
 - Deny all e-mail traffic from ISP X to Y (**access control**)
 - Route IP telephony traffic from X to Y via PHY_CIRCUIT (**policy**)
 - Ensure that no more than 50 Mbps are injected from ISP-X (**QoS**)

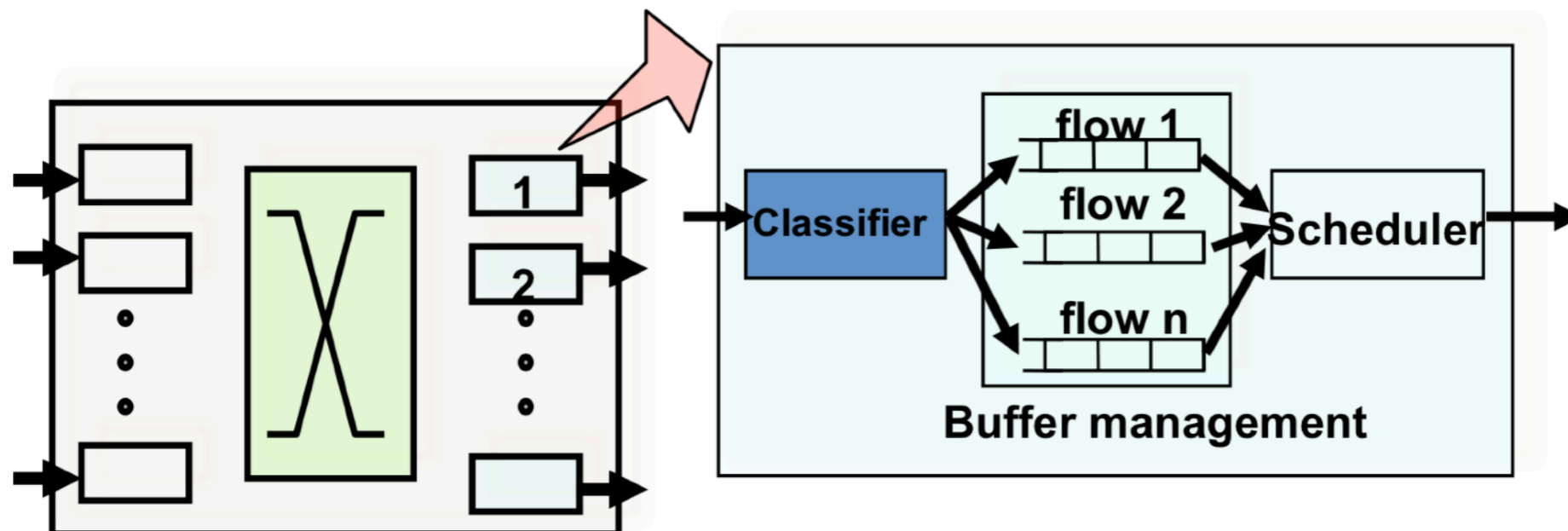
Simplest FIFO Router

- **No classification**
- **Drop tail buffer management:** when buffer is full drop incoming packet
- **First In First Out (FIFO) Scheduling:** schedule packets in order of arrival



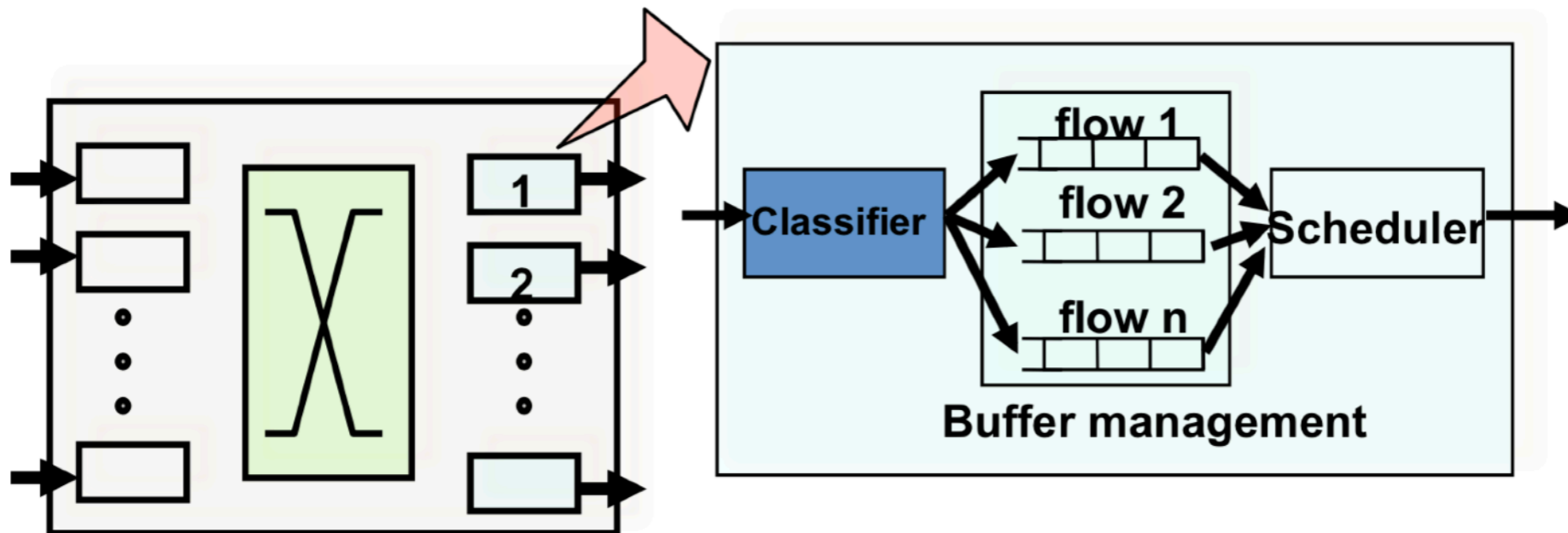
Packet Classification

- Classify an IP packet based on the number of fields in the packet header
 - Source/destination IP address (32 bits)
 - Source/destination TCP port number (16 bits)
 - Type of Service (TOS) byte (8 bits)
 - Type of Protocol (8 bits)
- In general fields are specified by range
 - Classification requires a multi-dimensional range search



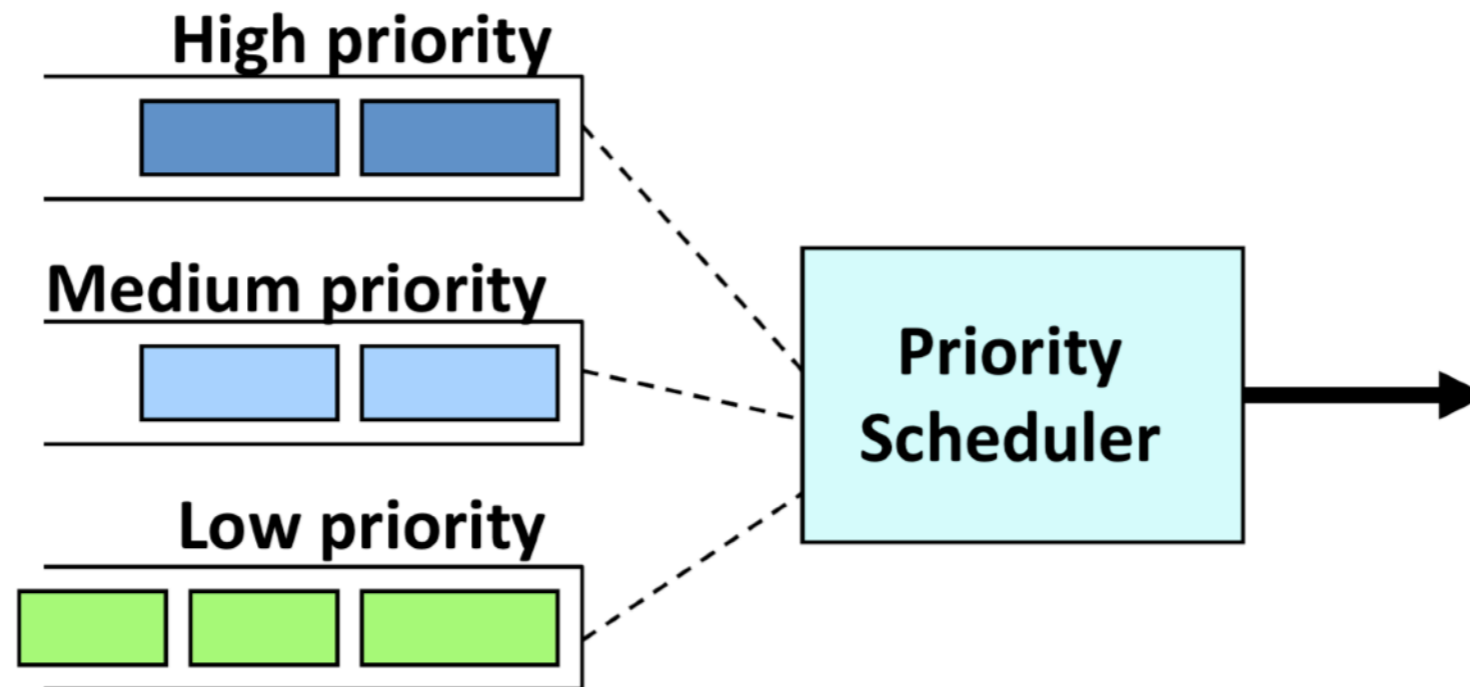
Scheduler

- One queue per flow
- Scheduler decides from which queue to send a packet
- Goals of scheduling algorithm
 - Fast!
 - Depends on the policy being implemented (fairness, priority, etc.)



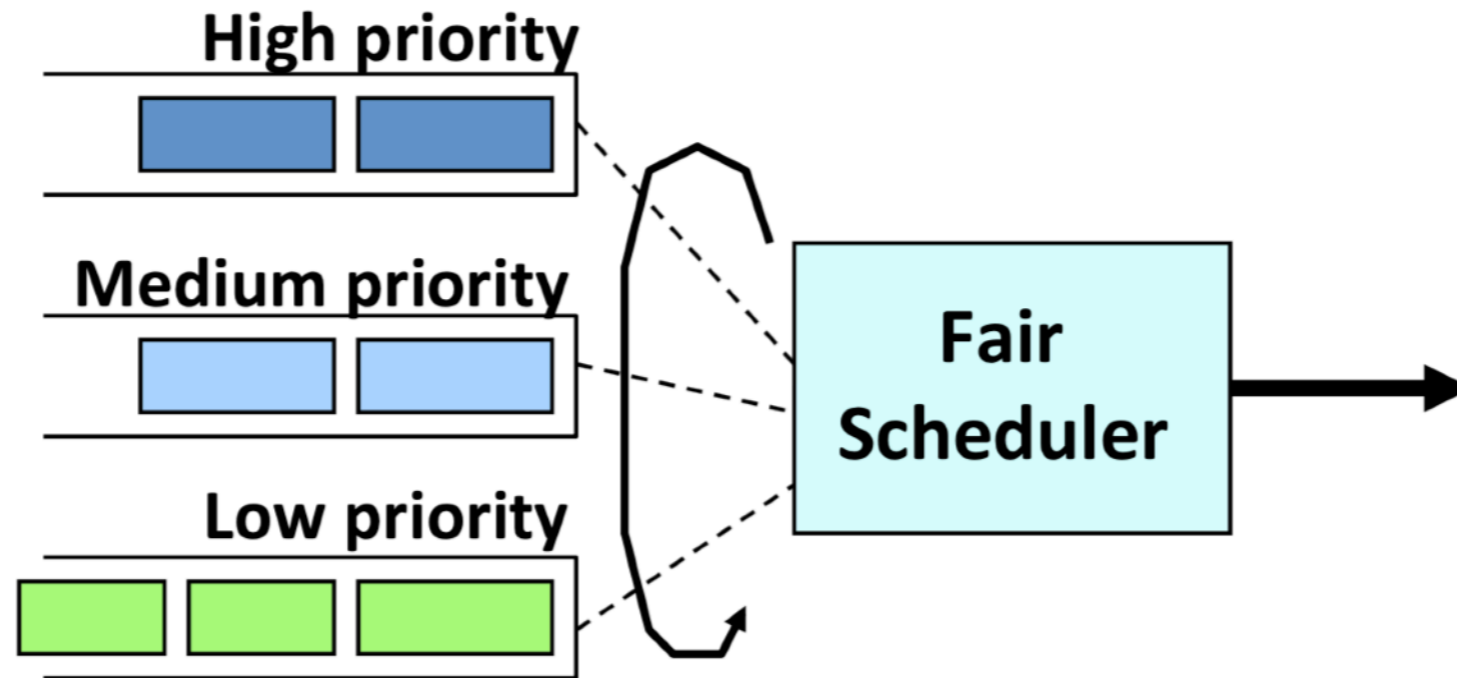
Example: Priority Scheduler

- Packets in the highest priority queue are always served before the packets in the lower priority queues



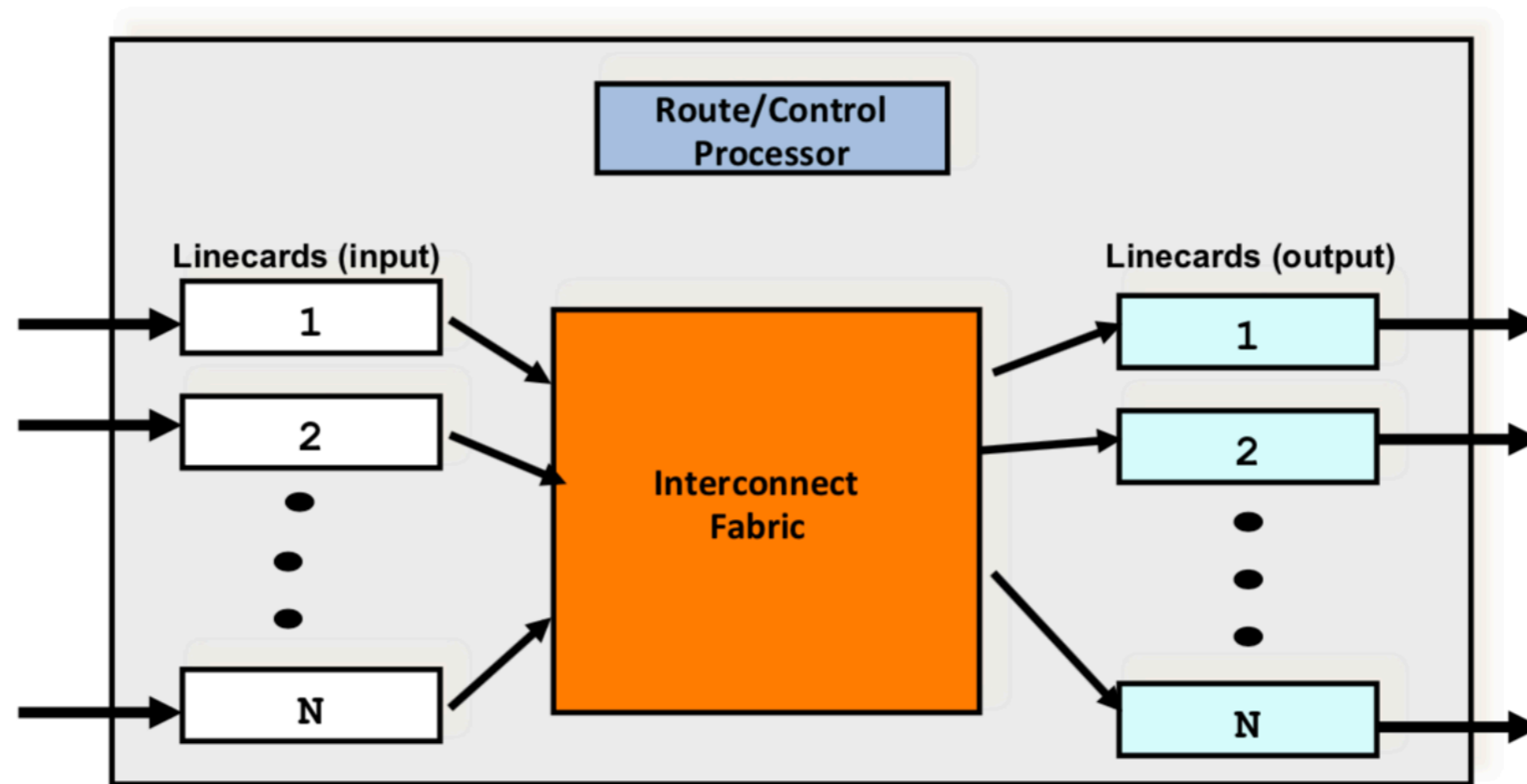
Example: Round Robin Scheduler

- Packets are served from each queue in turn

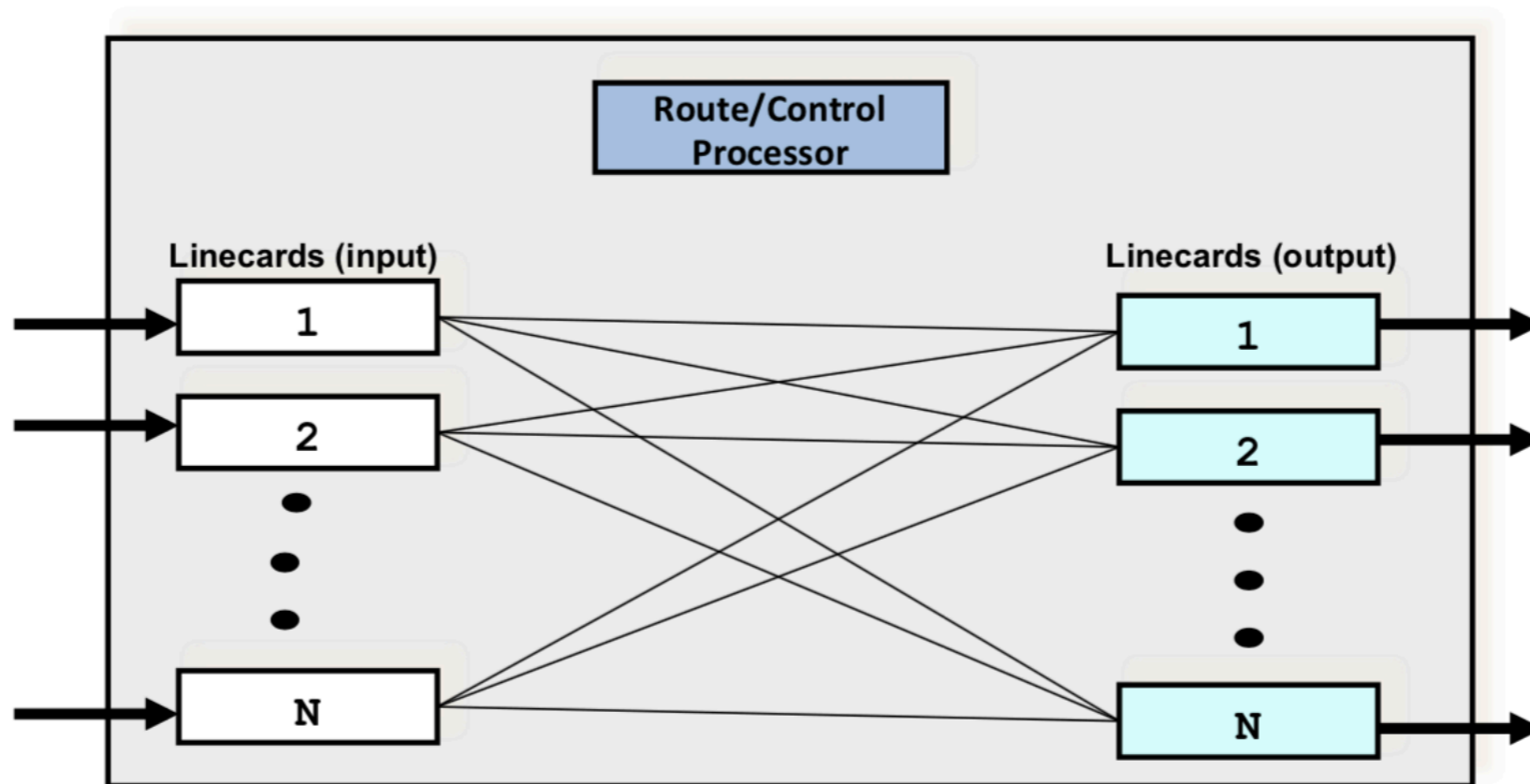


Connecting Input to Output: Switch Fabric

- Priority Scheduler: packets are served from each queue in turn



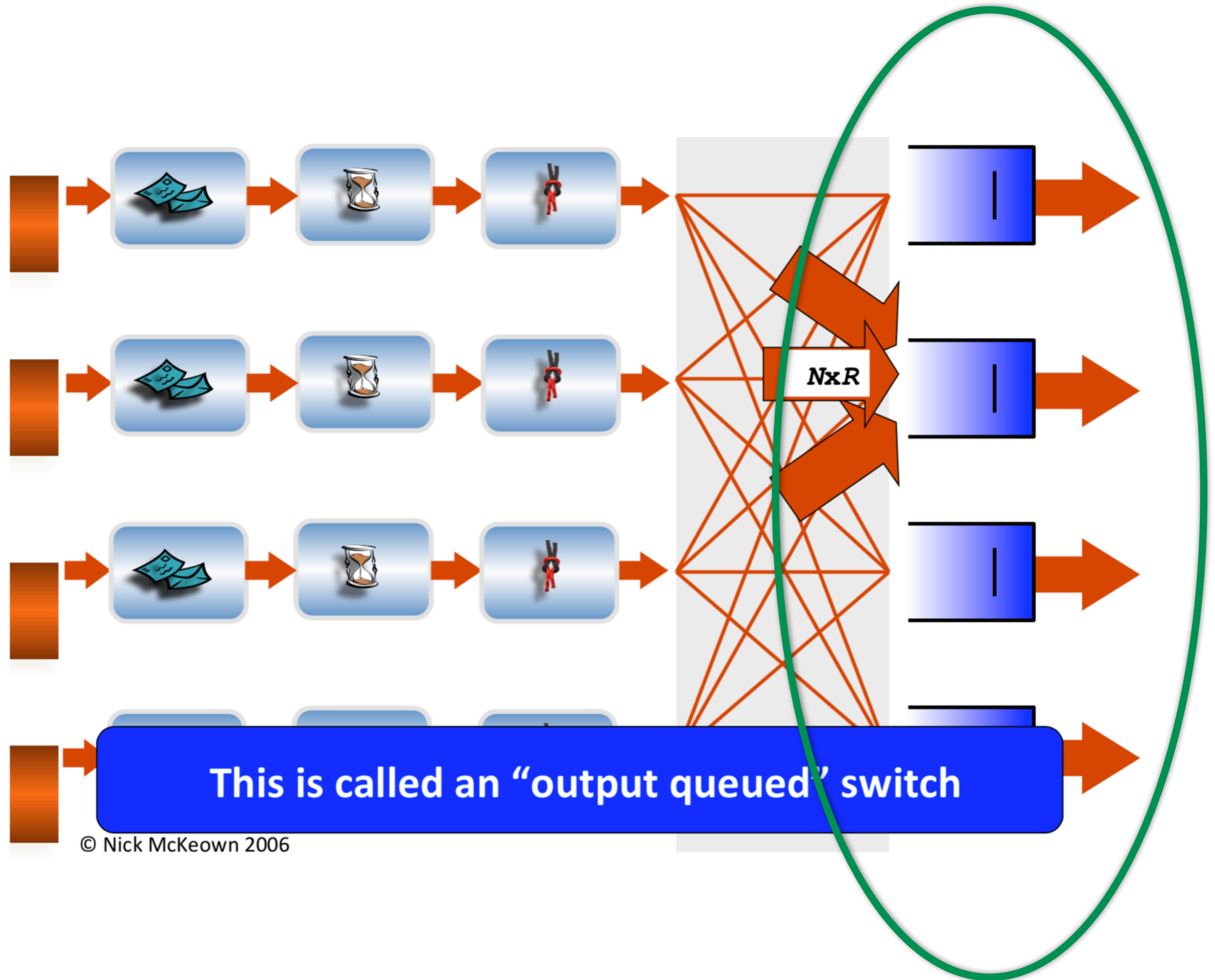
Today's Switch Fabrics: Mini Network!



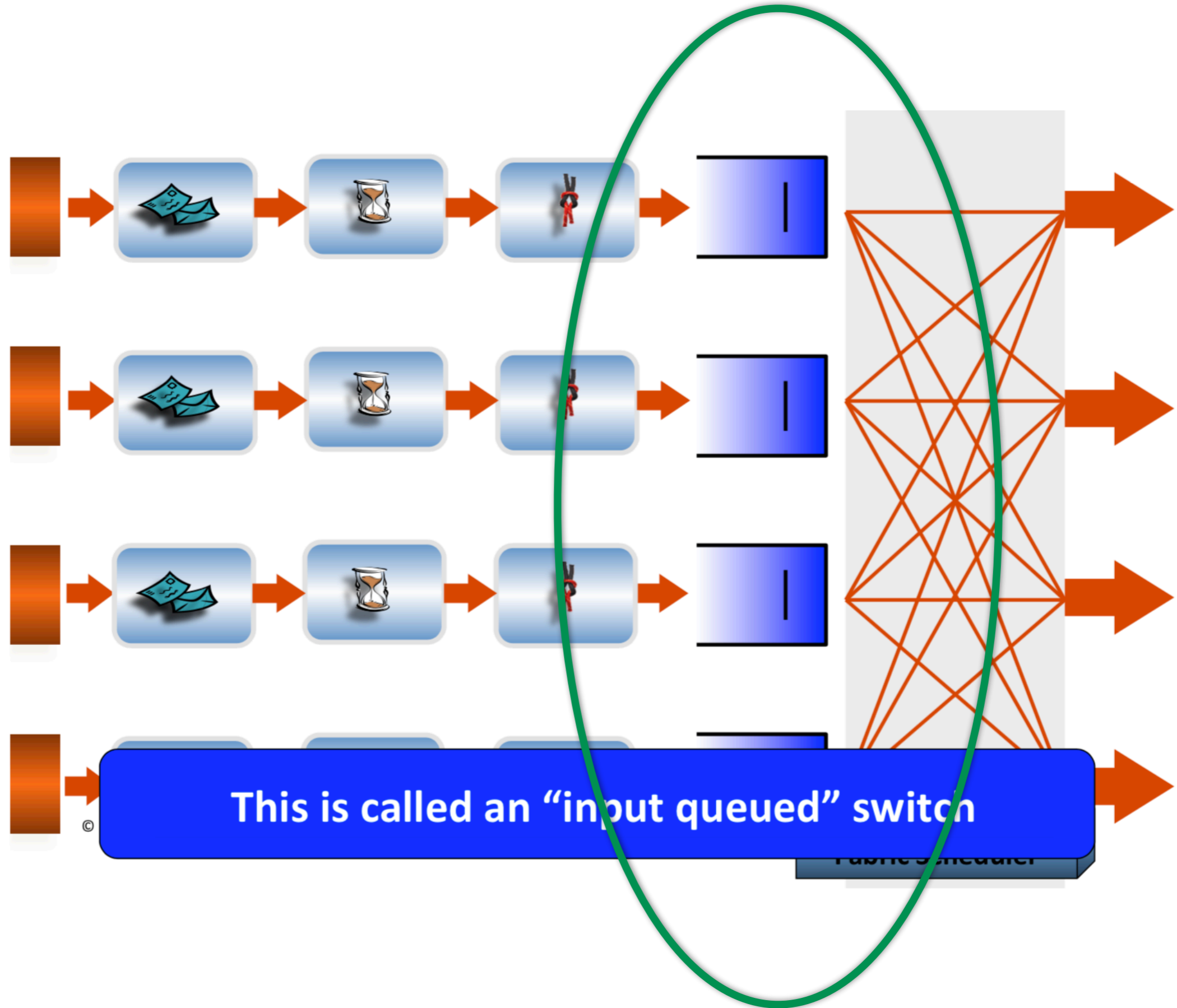
What's Hard About the Switch Fabric?

Queueing!

Third Generation Router: Switched Interconnects



Third Generation Router: Switched Interconnects



Reality is More Complicated

- Commercial high-speed routers use
 - Combination of input and output queueing
 - Complex multi-stage “topologies”
 - Distributed multi-stage schedulers (for scalability)

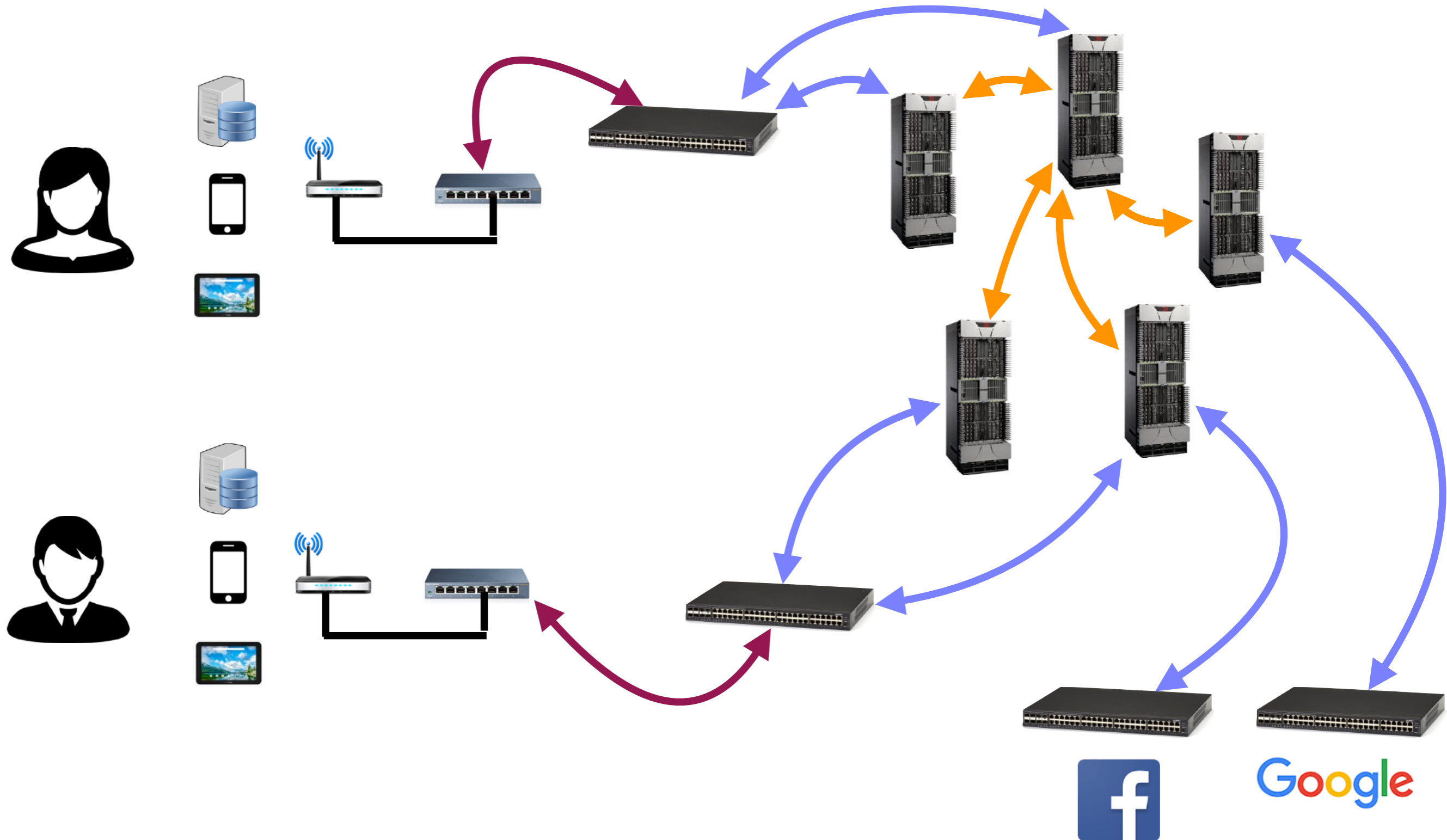
IP Routers Recap

- Core building block of Internet infrastructure
- Scalable Routing -> Longest Prefix Matching
- Need fast implementations for
 - Longest prefix matching
 - Switch fabric scheduling

Discovery Protocols

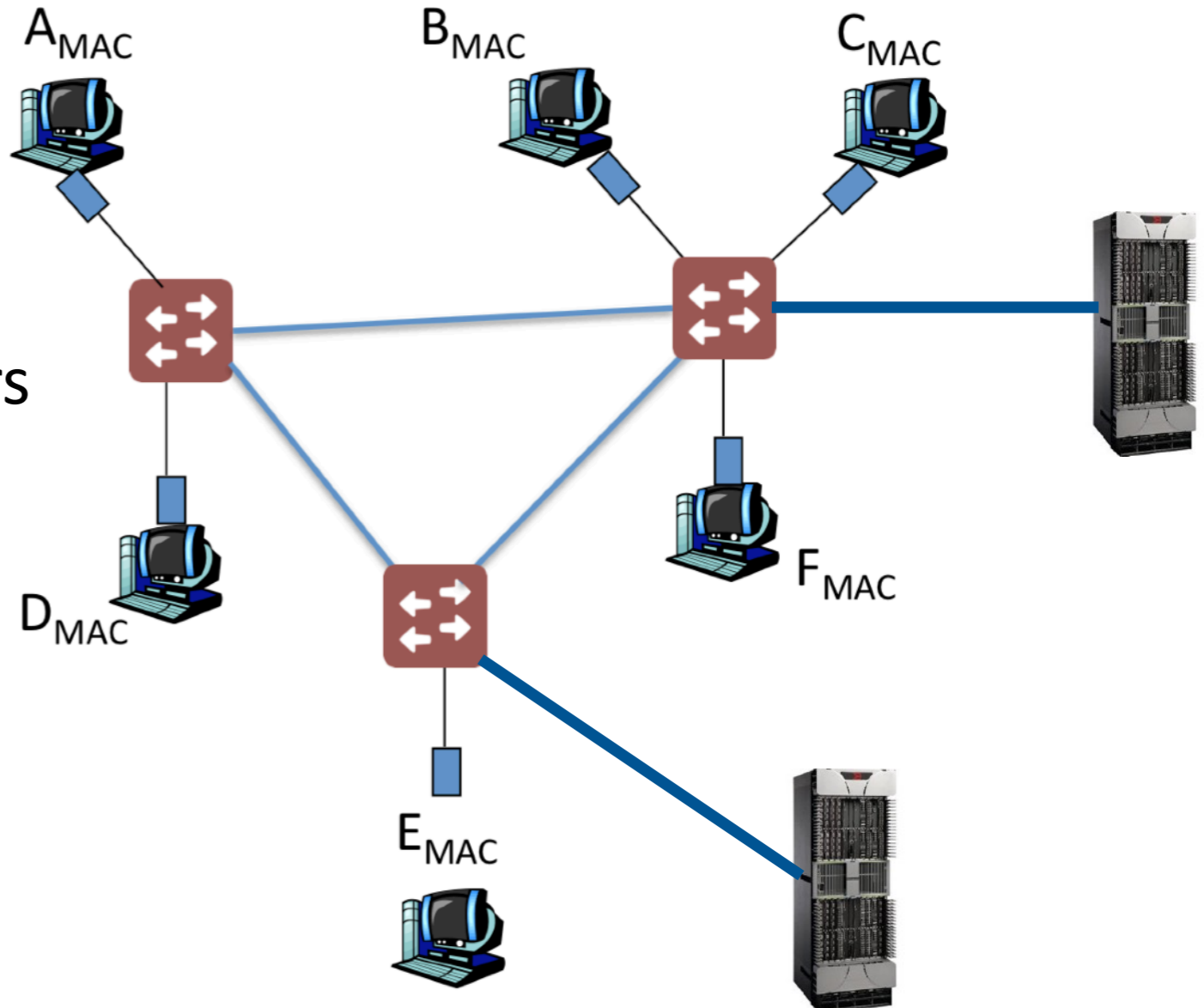
What is a computer network?

A set of network elements connected together, that implement a set of protocols for the purpose of sharing resources at the end hosts

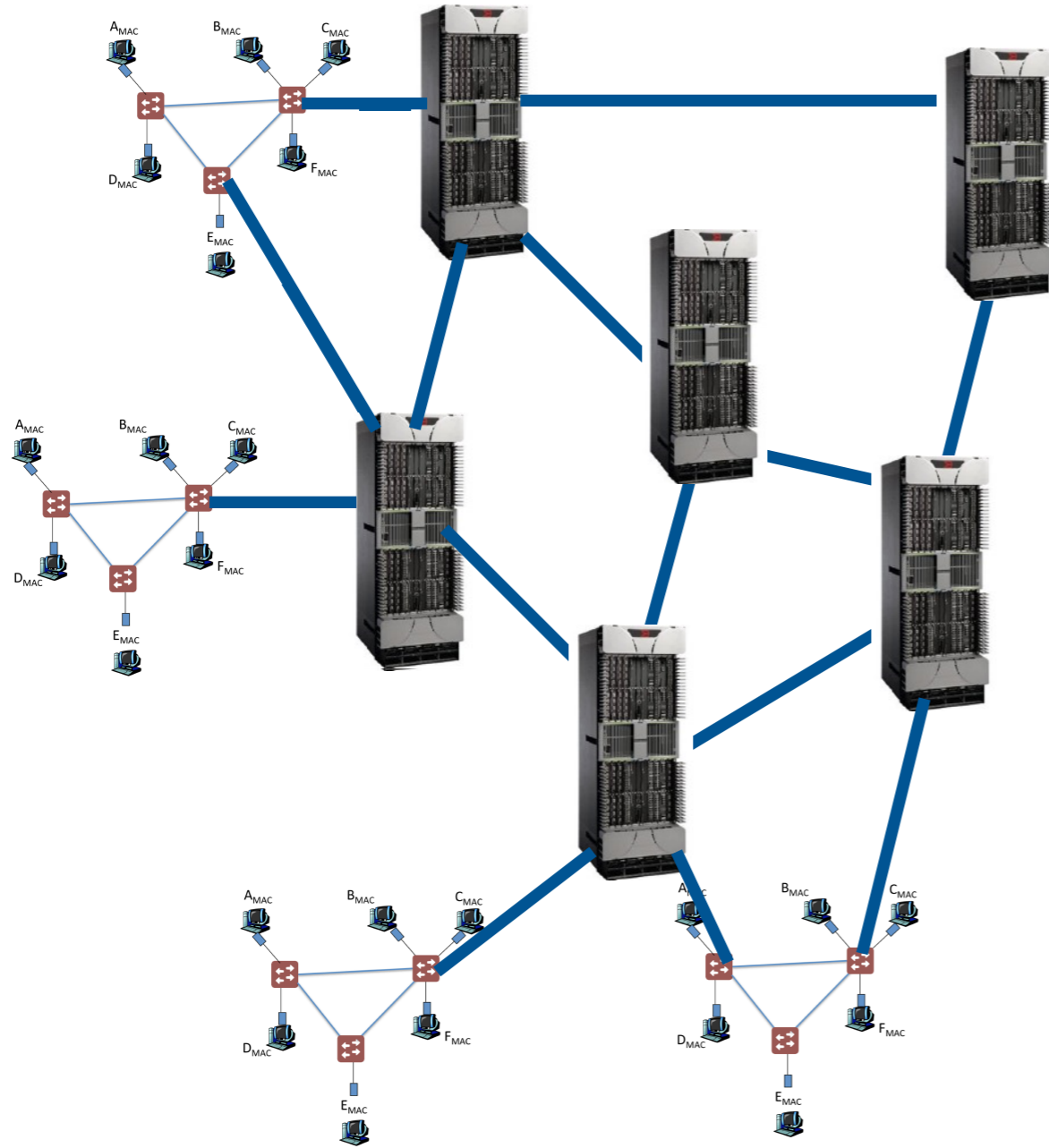


What does Internet actually look like?

- **The smallest component:**
 - A Network Interface Card (NIC), or a machine, or a server
 - Has a **Link Layer MAC name/address**
- **Multiple NICs connected in a Local Area Network (LAN) via**
 - Broadcast Ethernet,
 - Or, Switched Ethernet
- **Switches in LAN**
 - Connected to larger routers

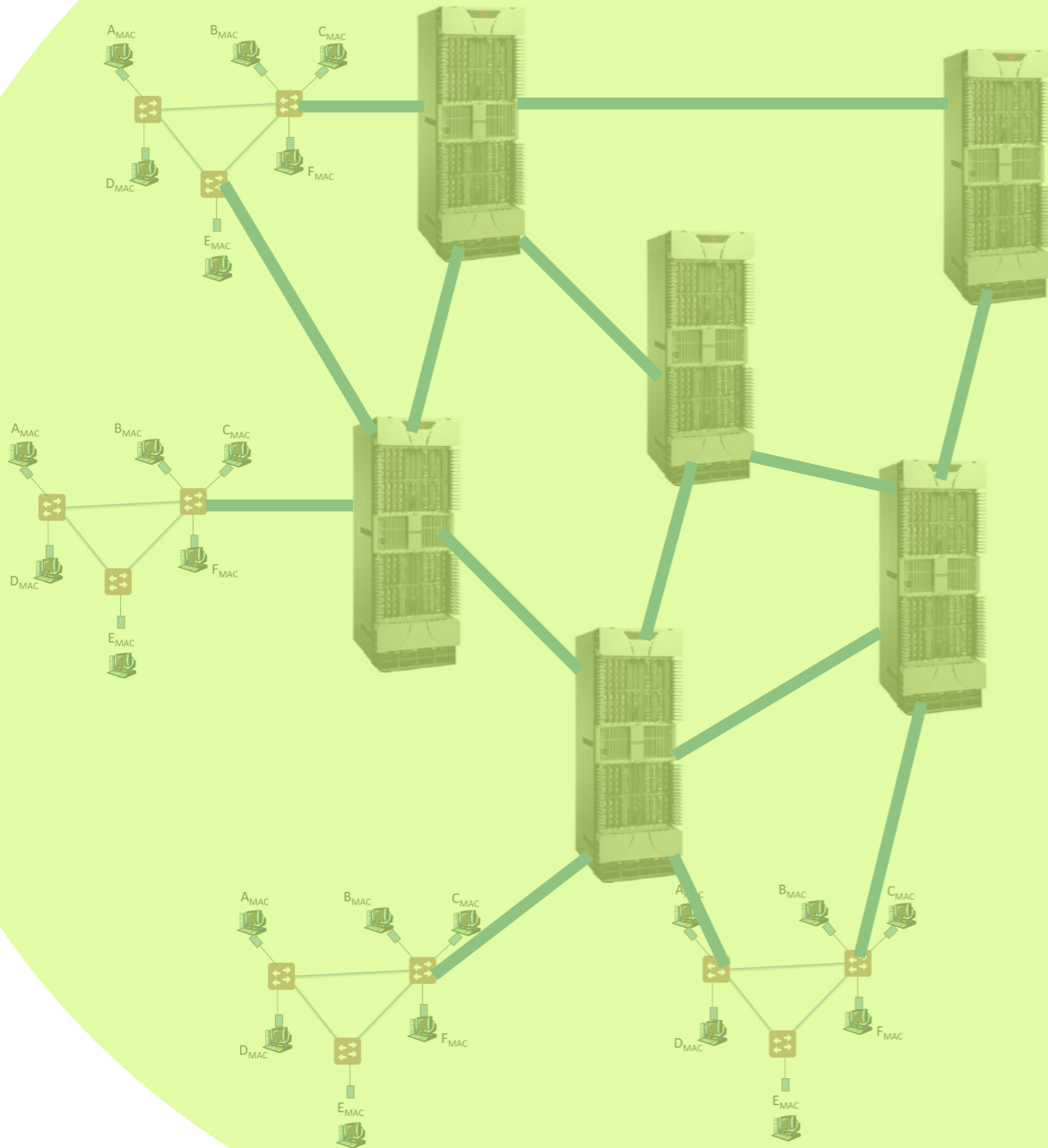


What does Internet actually look like?



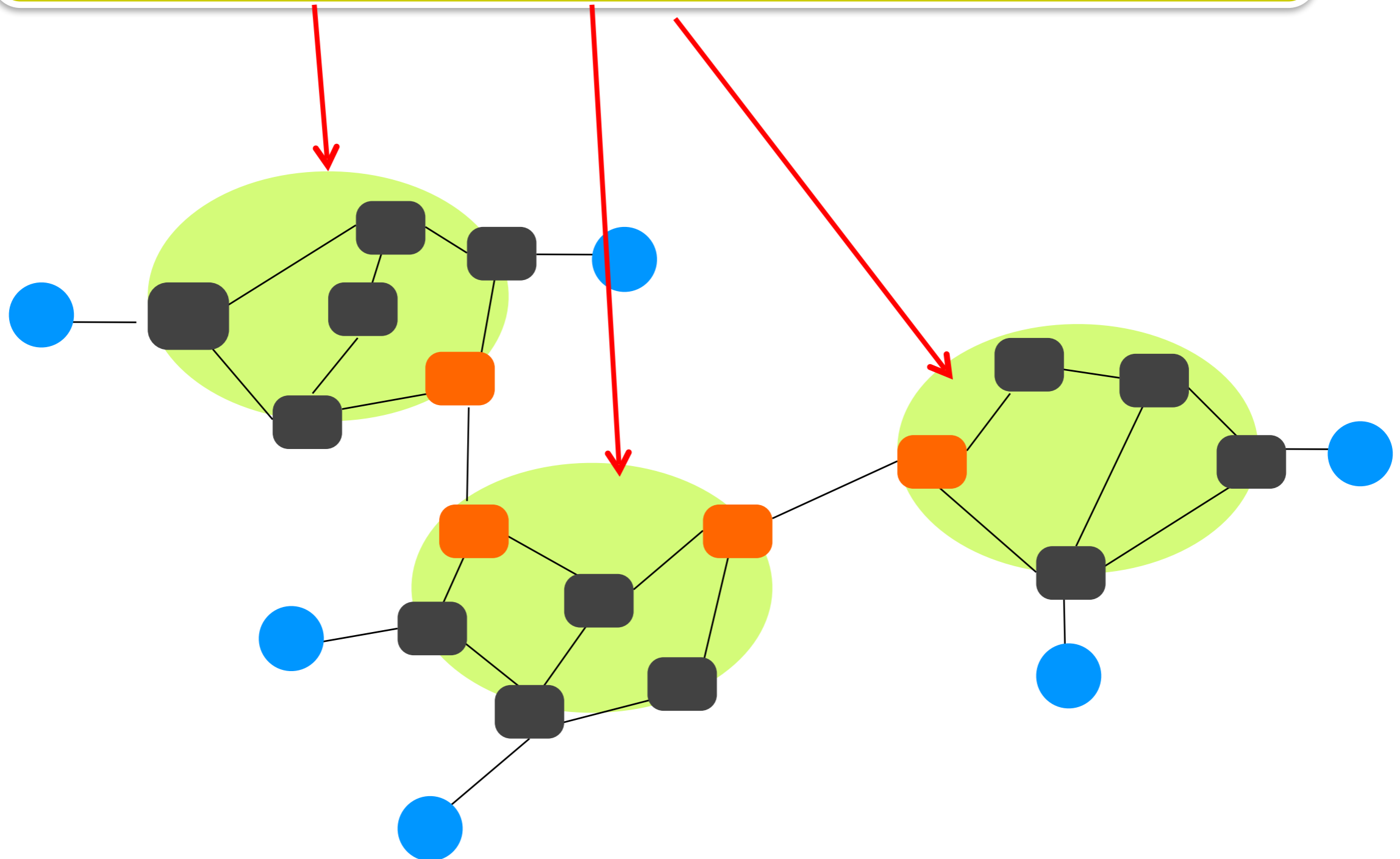
What does Internet actually look like?

“Autonomous System (AS)” or “Domain”
Region of a network under a single administrative entity

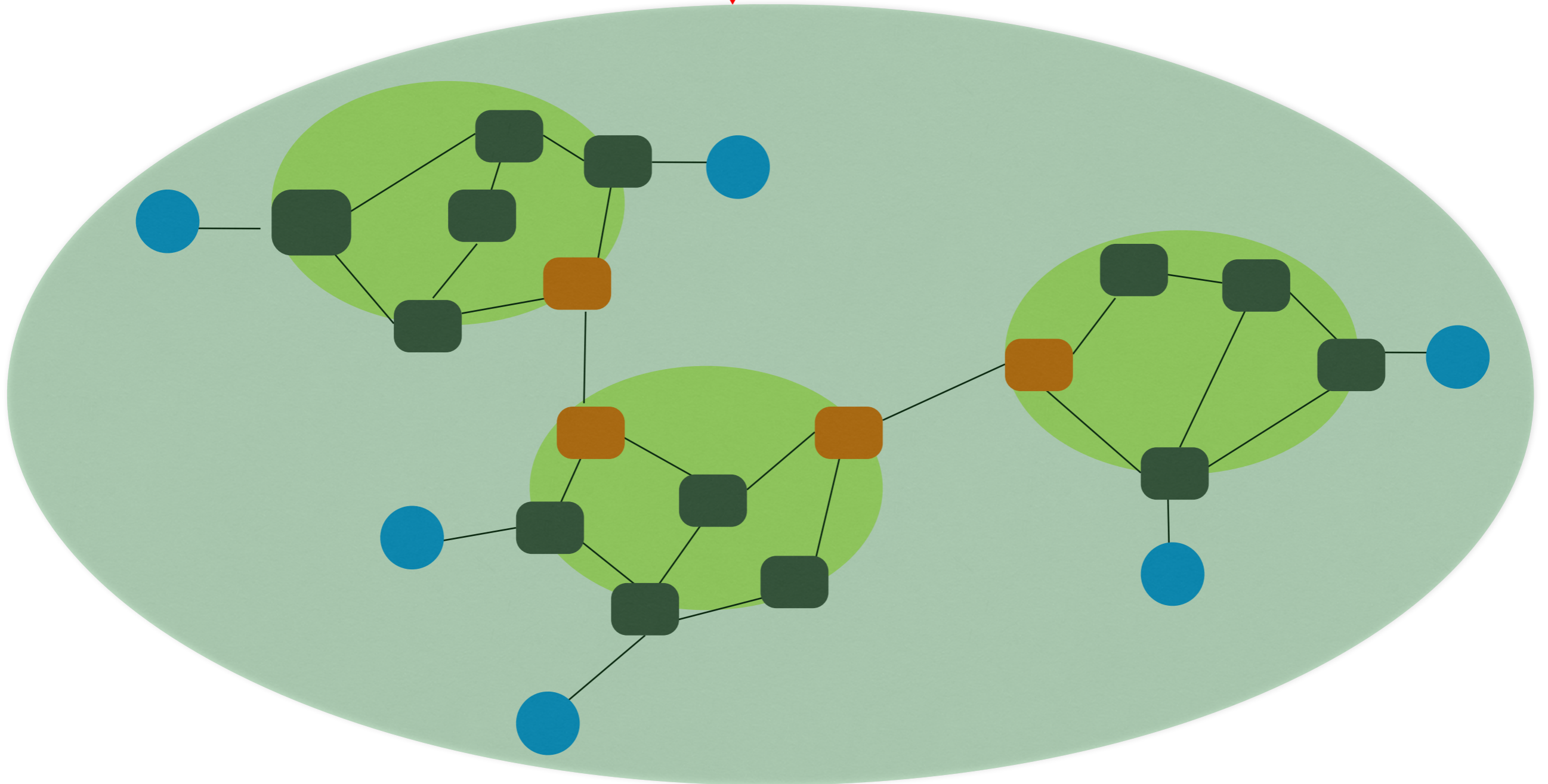


What does Internet actually look like?

Multiple “Autonomous Systems (AS)” or “Domains” connect together using Border Routers



This entire infrastructure is a part of the INTERNET :-)

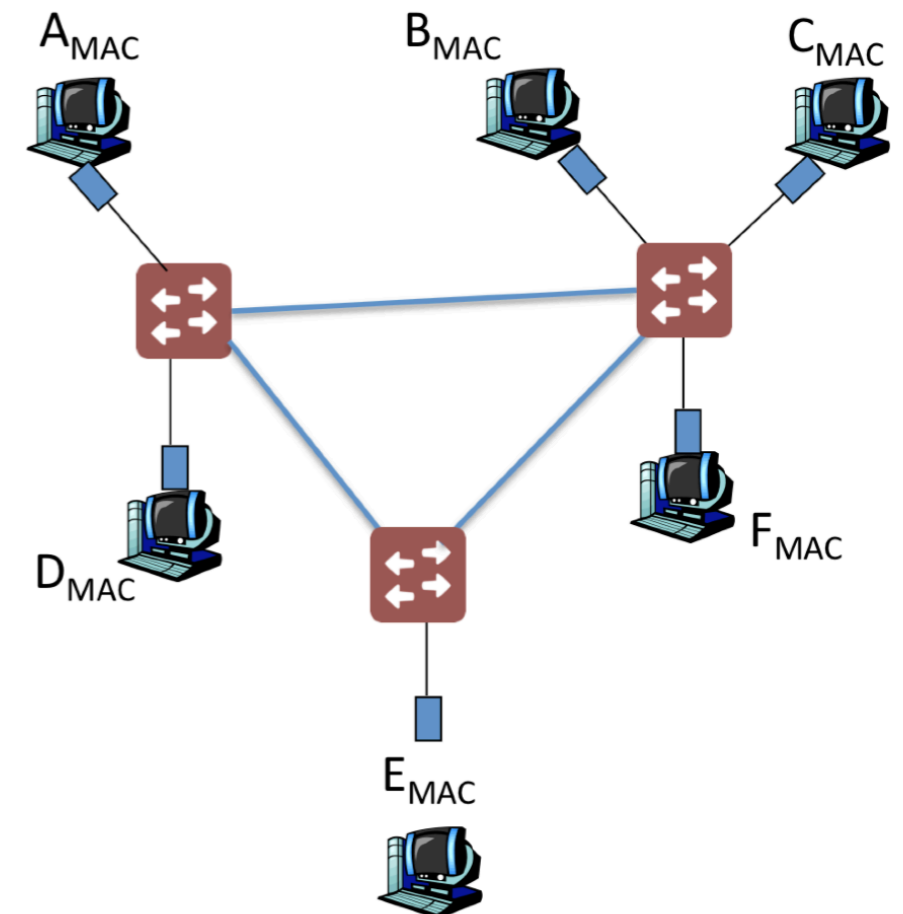


What is the other part of the Internet?

Protocols!

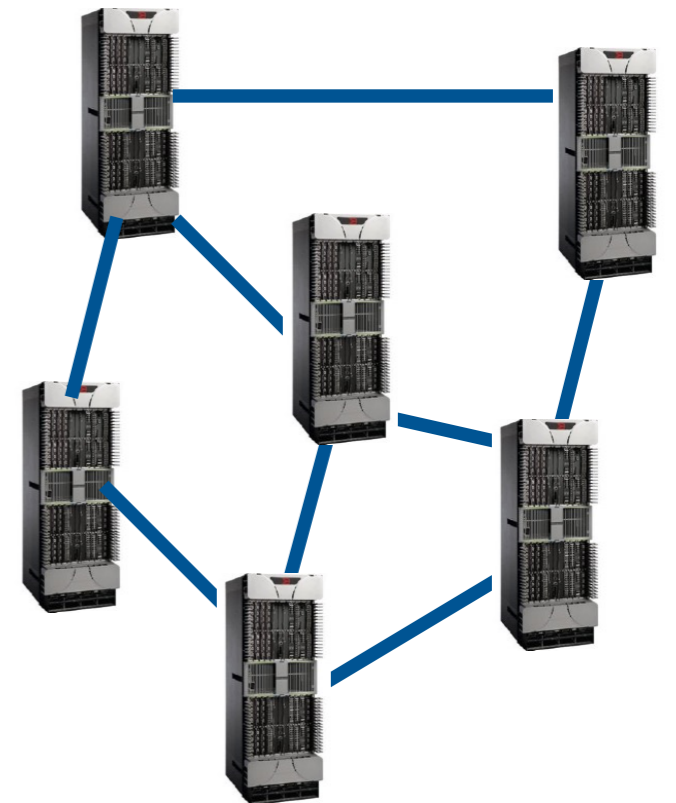
What protocols have we learnt on LAN?

- **Addresses**
 - Link Layer MAC names/addresses: come with the hardware
- **CSMA/CD Protocol:**
 - For transmitting frames on broadcast Ethernet
- **Spanning Tree Protocol:**
 - For transmitting frames on switched Ethernet



What have we learnt beyond LAN?

- **Link-state and Distance-vector Protocols:**
 - For finding routes (and a next-hop) to an IP address within an ISP
- **Border Gateway Protocol:**
 - For finding routes to an IP address range
- **Forwarding at routers**
 - Store **routing tables** (map **destination prefixes** to outgoing port)
 - Longest prefix match for destination address lookup



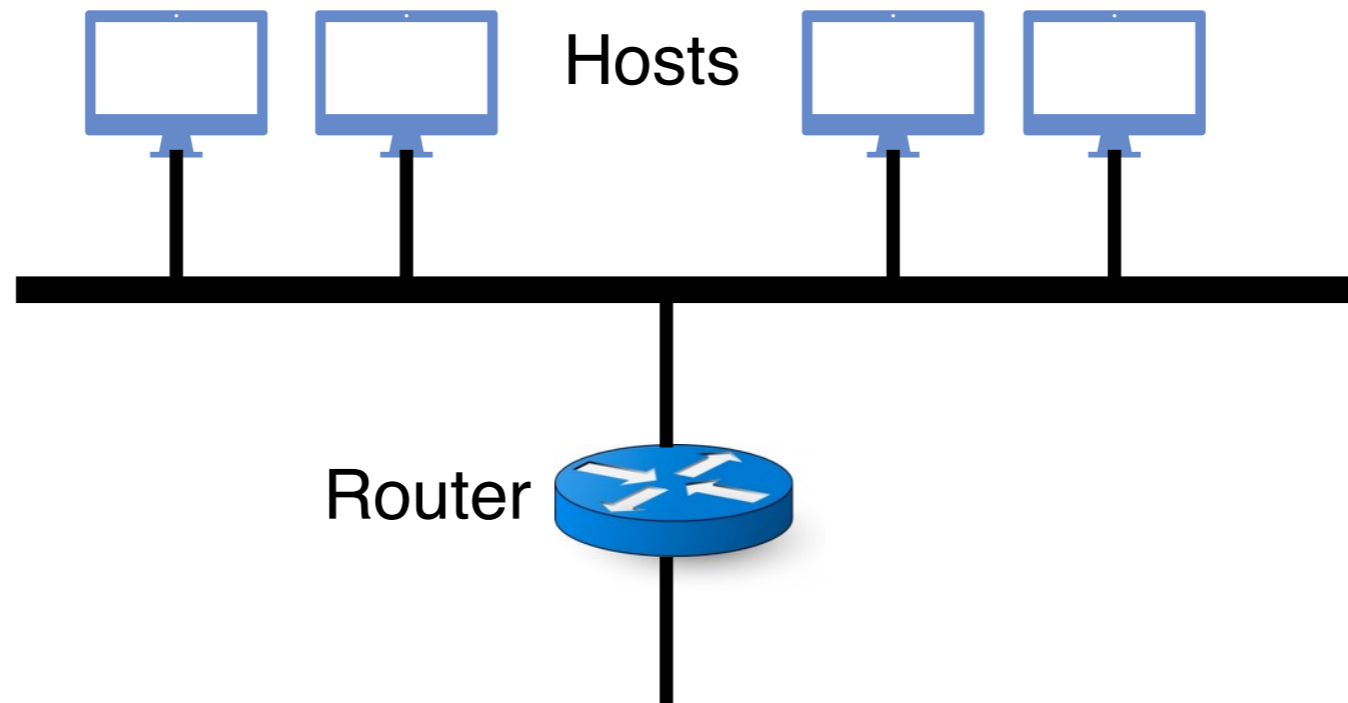
Suppose Host A wants to communication with Host B

Discovery

- Suppose I am host A
- I want to communicate with B (say, www.google.com)
- I was “born” knowing **only** my name — my MAC address :-)
- Must discover some information before I can communicate with B
 - What is my IP address?
 - What is B’s IP address?
 - Using DNS
 - Is B within my LAN?
 - If yes, what is B’s MAC address?
 - If not, what is the address of my first-hop router to B?
 - ...

DHCP and ARP

- Link layer discovery protocols
 - DHCP — Dynamic Host Configuration Protocol
 - ARP — Address Resolution Protocol
 - Configured to a single LAN
 - Rely on broadcast capability



DHCP and ARP

- Link layer discovery protocols
- Serve two functions
 1. Discovery of local end-hosts
 - For communication between hosts on the same LAN
 2. Bootstrap communication with remote hosts
 - What's my IP address?
 - Who/where is my local DNS server?
 - Who/where is my first hop router?

DHCP

- Dynamic Host Configuration Protocol
 - Defined in RFC 2131
- A host uses DHCP to discover
 - Its own IP address
 - Subnet masks — allows to test whether an IP address is local or not
 - IP address(es) for its local DNS name server(s)
 - IP address(es) for its first-hop “default” router(s)

DHCP: operation

1. One or more local DHCP servers maintain required information
 - IP address pool, netmask, DNS servers, etc.
 - Application that listens on UDP port 67

DHCP: operation

1. One or more local DHCP servers maintain required information
2. Client broadcasts a DHCP discovery message
 - L2 broadcast, to MAC address FF:FF:FF:FF:FF:FF

DHCP: operation

1. One or more local DHCP servers maintain required information
2. Client broadcasts a DHCP discovery message
3. One or more DHCP servers respond with a DHCP “offer” message
 - Proposed IP address for client, lease time
 - Other parameters

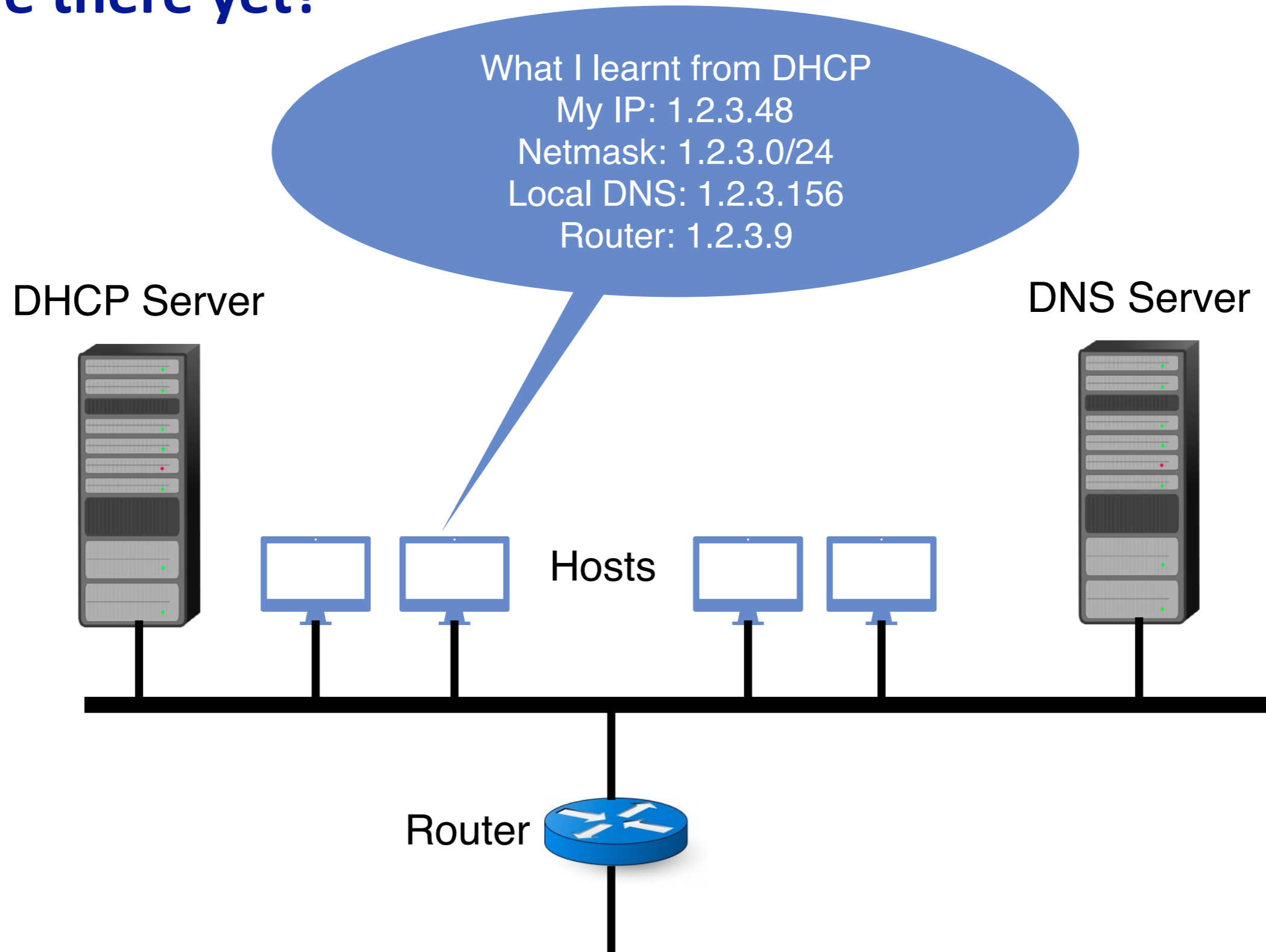
DHCP: operation

1. One or more local DHCP servers maintain required information
2. Client broadcasts a DHCP discovery message
3. One or more DHCP servers respond with a DHCP “offer” message
4. Client broadcasts a DHCP request message
 - Specifies which offer it wants
 - Echoes accepted parameters
 - Other DHCP servers learn they were not chosen

DHCP: operation

1. One or more local DHCP servers maintain required information
2. Client broadcasts a DHCP discovery message
3. One or more DHCP servers respond with a DHCP “offer” message
4. Client broadcasts a DHCP request message
5. Selected DHCP server responds with an ACK

Are we there yet?



ARP: Address Resolution Protocol

- Every host maintains an ARP table
 - List of (IP address — MAC address) pairs
 - For IP addresses within the same LAN
- Consult the table when sending a packet
 - Map destination IP address to destination MAC address
- But: what if IP address not in the table?
 - Either its not local (detected using DHCP)
 - If its local:
 - Sender broadcasts: “Who has IP address 1.2.3.156?”
 - Caches the answer in ARP table

Key Ideas in Both ARP and DHCP

- Broadcasting: can use broadcast to make contact
 - Scalable because of limited size
- Caching: remember the past for a while
 - Store the information you learn to reduce overhead

Taking Stock: Discovery

Layer	Examples	Structure	Configuration	Resolution Service
App Layer	<u>www.cs.cornell.edu</u>	Organizational hierarchy	~ manual	↕ DNS
Network Layer	123.45.6.78	Topological hierarchy	DHCP	
Link Layer	45-CC-4E-12-F0-97	Vendor(flat)	Hard-coded	↕ ARP

How does the Internet work?

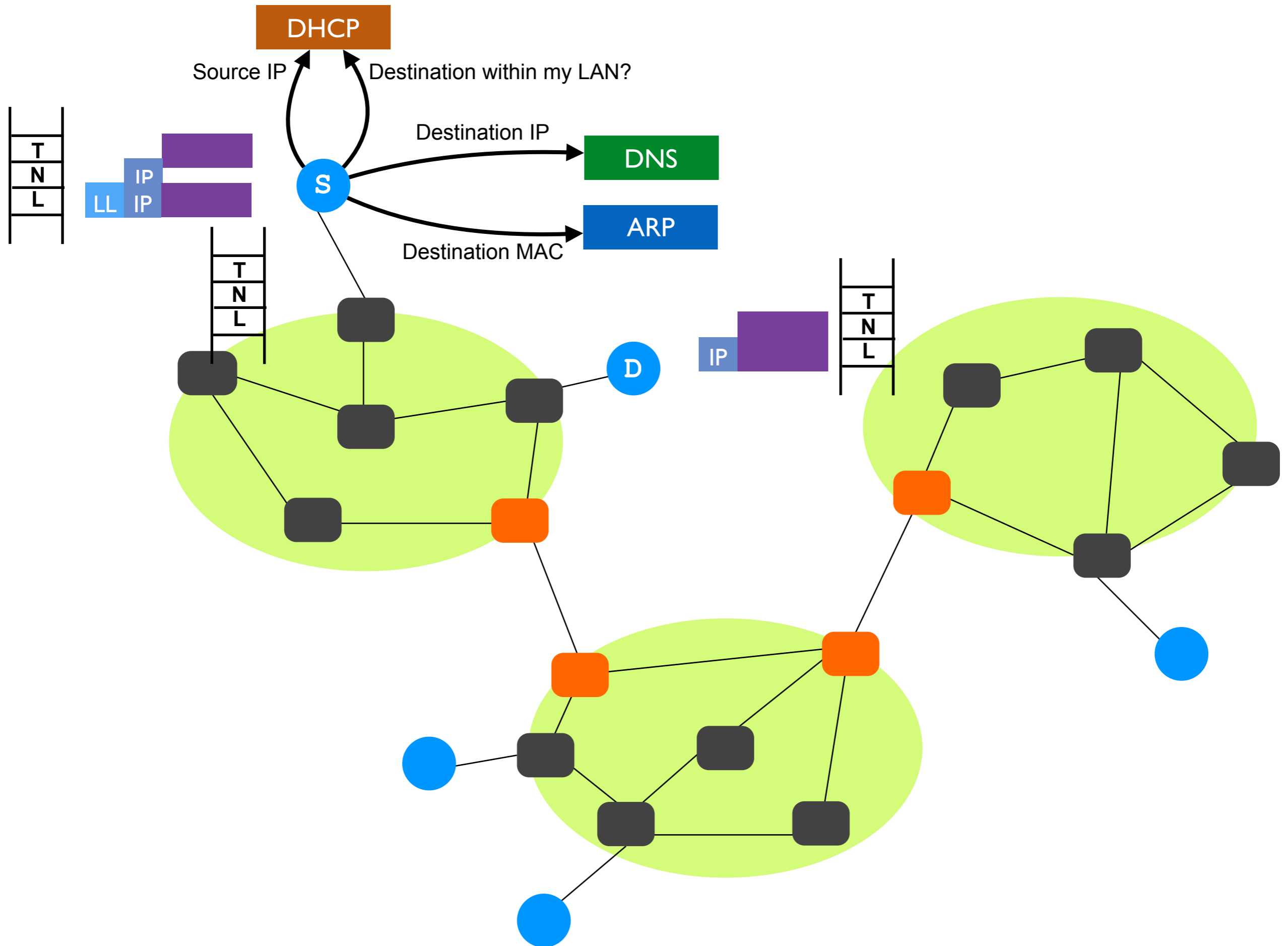
Are you ready?

(Count the number of protocols used for each packet)

How does Internet work — end-to-end?

- Network stack receives the packet from the application (roughly speaking)
- What is my IP address? (using DHCP)
- What is the destination IP address? (using DNS)
- Is destination IP address within my LAN? (using DHCP)
- **If destination IP address local:**
 - What is destination MAC address (using ARP)?
 - Convert packet into frames with correct source/destination address
 - Convert frames into bits
 - Forward the bits to the wire ...
- **Each switch:**
 - Forwards to destination (using STP/CSMA/CD)

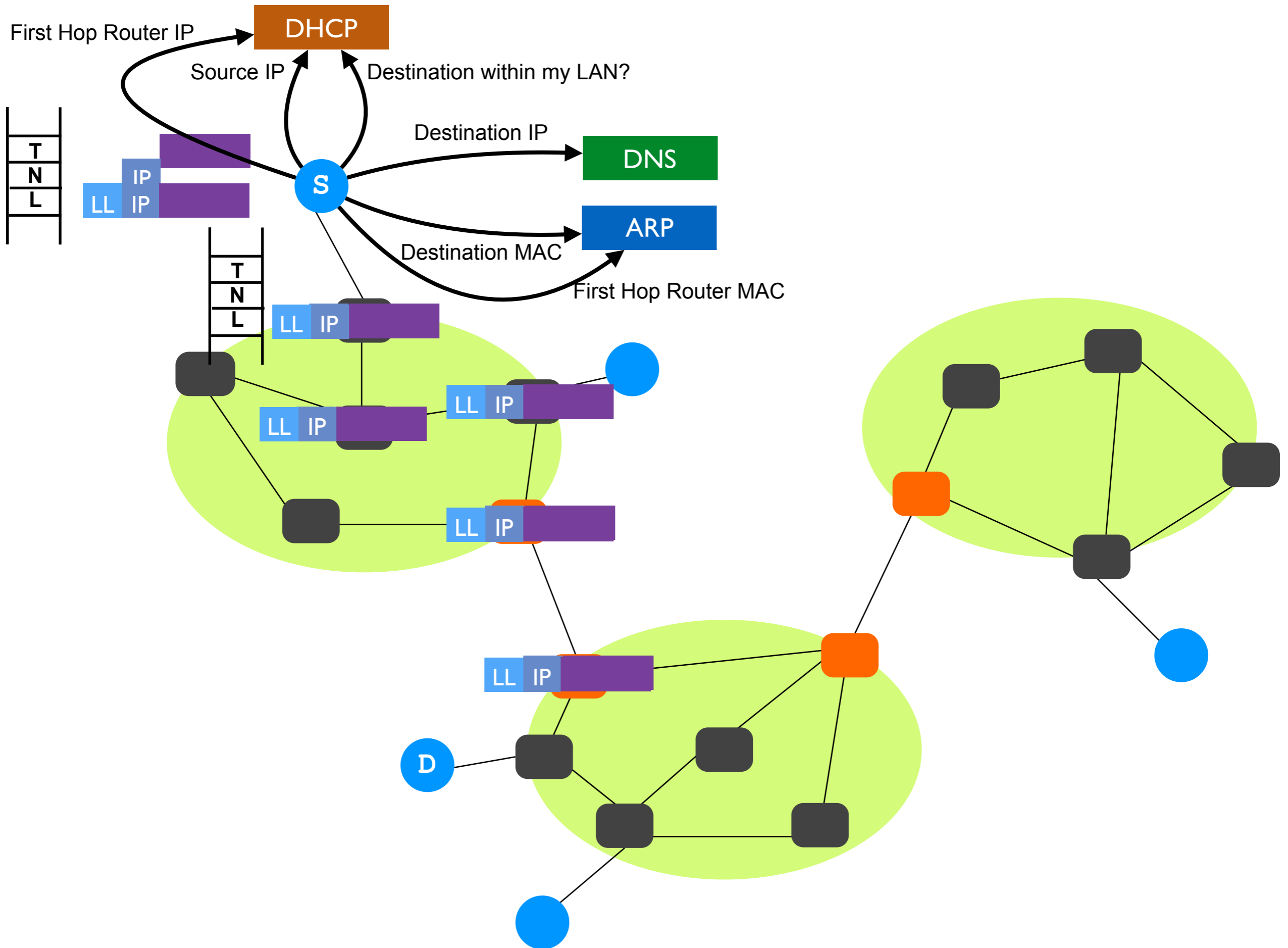
End-to-End I



How does Internet work — end-to-end?

- Network stack receives the packet from the application (roughly speaking)
- What is my IP address? (using DHCP)
- What is the destination IP address? (using DNS)
- Is destination IP address within my LAN? (using DHCP)
- **If destination IP address remote:**
 - **What is my next-hop router IP address? (using DHCP)**
 - **What is my next-hop router MAC address? (using ARP)**
 - Convert packet into frames with correct source/destination address
 - Convert frames into bits
 - Forward the bits to the wire ...
- **Each router**

End-to-End II



How does Internet work — end-to-end?

A router upon receiving a packet (implicit questions)

- **Is the destination in a LAN connected to me?**
 - Forward the packet to the destination
 - Using STP/CSMA/CD
- **Is the destination not in my LAN but in my ISP?**
 - Forward the packet to the next-hop router towards the destination
 - Using distance-vector routing algorithm
- **Is the destination in a different ISP?**
 - Forward the packet to the next-hop router towards the destination
 - Using BGP routing algorithm

Are We There Yet?

- Yes!
- How can we be sure?
- Lets go back to where we started

Recall the end-to-end story from our fifth lecture :-)

- Application opens a **socket** that allows it to connect to the **network stack**
- Maps **name** of the web site to its **address** using **DNS**
- The network stack at the source embeds the address and **port** for both the source and the destination in **packet header**
- Each **router** constructs a **routing table** using a distributed algorithm
- Each router uses destination address in the packet header to look up the **outgoing link** in the routing table
 - And when the link is free, forwards the packet
- When a packet arrives the destination:
 - The network stack at the destination uses the port to forward the packet to the right application

You now know how the Internet works!!!!

All that is remaining:

Reliability.