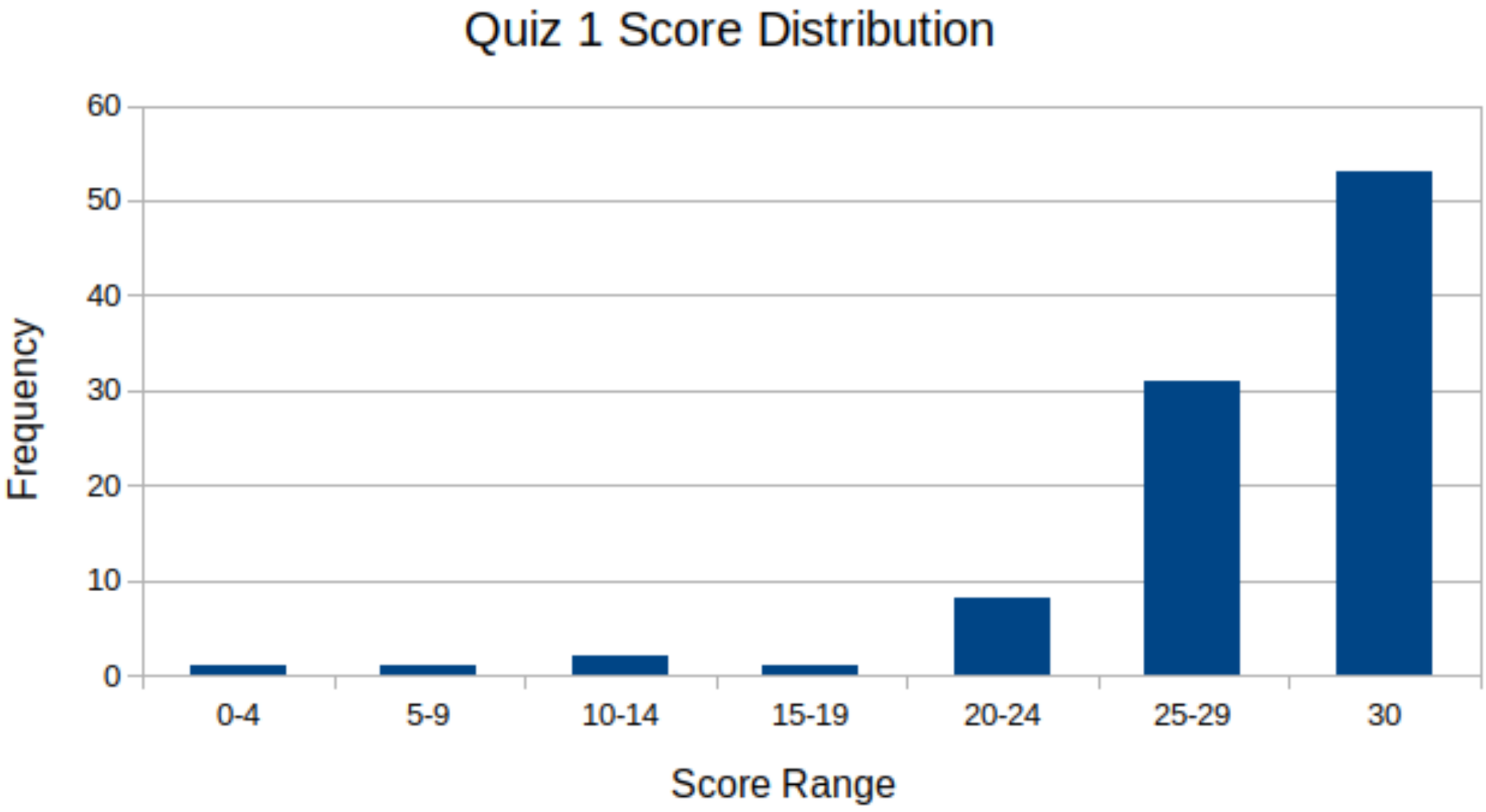
CS4450

Computer Networks: Architecture and Protocols

Lecture 10
Fundamentals of Routing

Rachit Agarwal


Quiz 1 Score Distribution
Frequency
60
50
40
30
20
10
0
0-4
5-9
10-14
15-19
20-24
25-29
30
Score Range

Goals for Today's Lecture

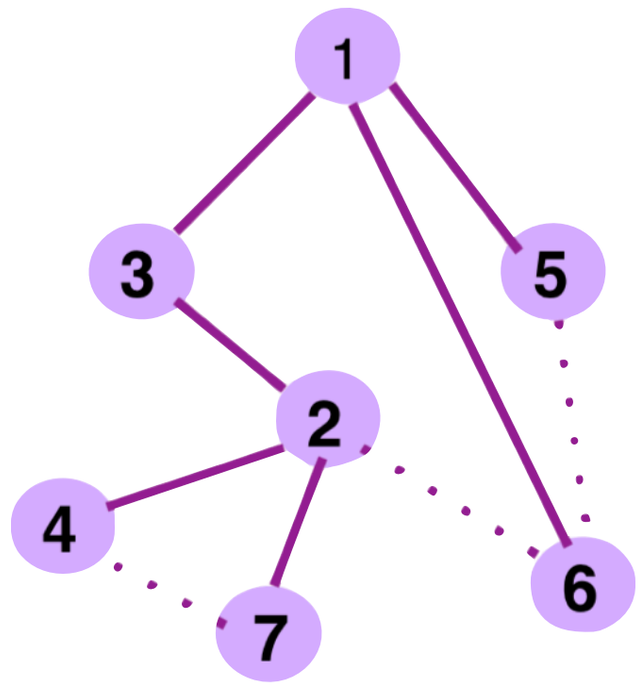
- **Recap: Why Network Layer?**
- **Recap: The right way to think about Routing Tables…**
- Correctness definition for routing tables
- Our first Network Layer Protocol: Link State

Recap from last lecture

Recap: Why do we need network layer?

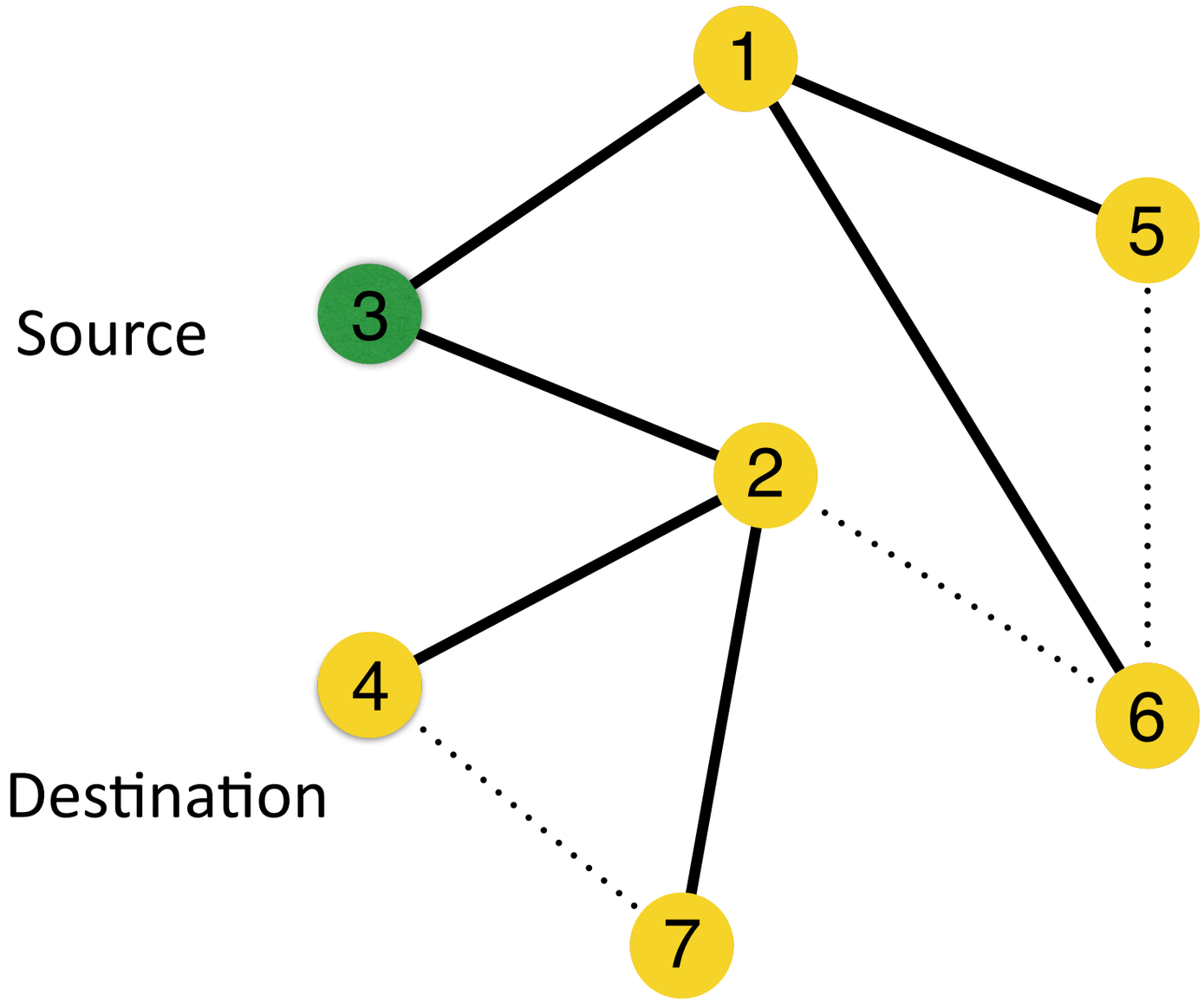
- Why not just use spanning trees across the entire network?

- Easy to design routing algorithms for (spanning) trees
    - **Step 1**: Source node "floods" its packet on its spanning tree links
    - **Step 2**: Whenever a node receives a packet:
        - Forwards incoming packet out to all links **other than the one that sent the packet**



Recap: Flooding Example

Eventually all nodes are covered



One copy of packet delivered to destination

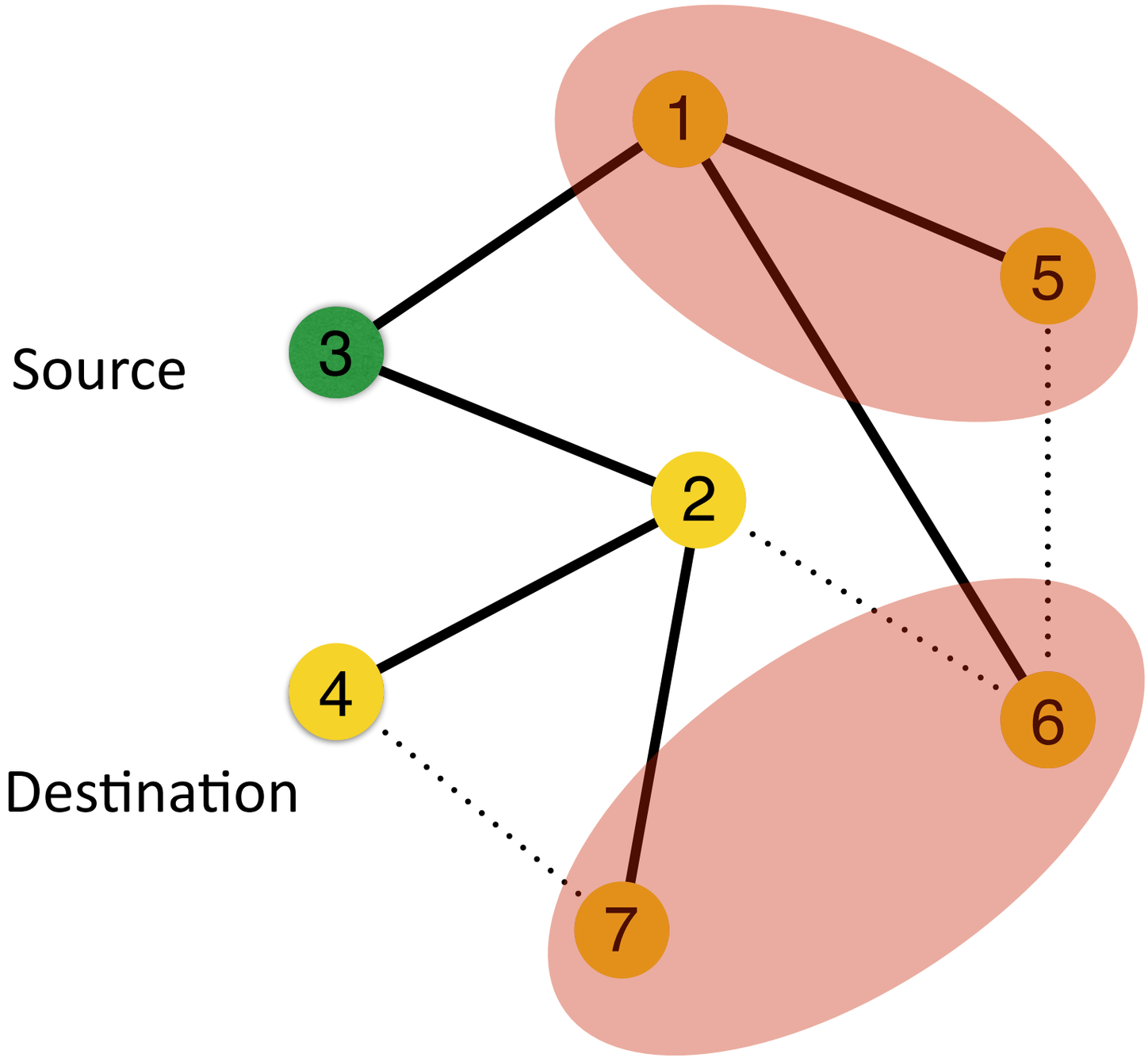
Recap: Why do we need network layer?

- Why not just use spanning trees across the entire network?

- Easy to design routing algorithms for (spanning) trees
    - **Step 1**: Source node "floods" its packet on its spanning tree links
    - **Step 2**: Whenever a node receives a packet:
        - Forwards incoming packet out to all links **other than the one that sent the packet**

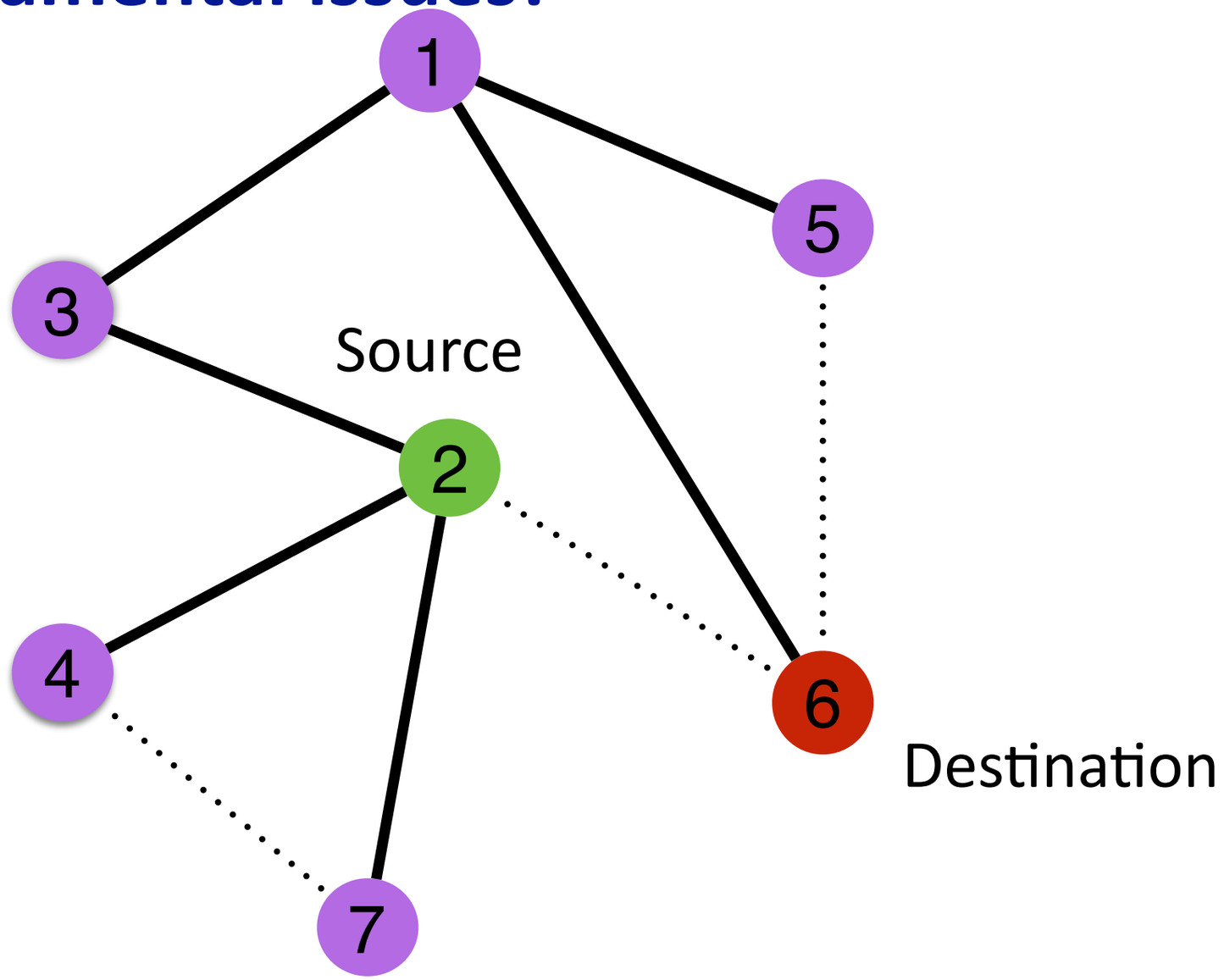
- **Amazing properties:**
    - No routing tables needed!
    - No packets will ever loop.
    - At least (and exactly) one packet must reach the destination
        - Assuming no failures

Recap: Three fundamental issues!



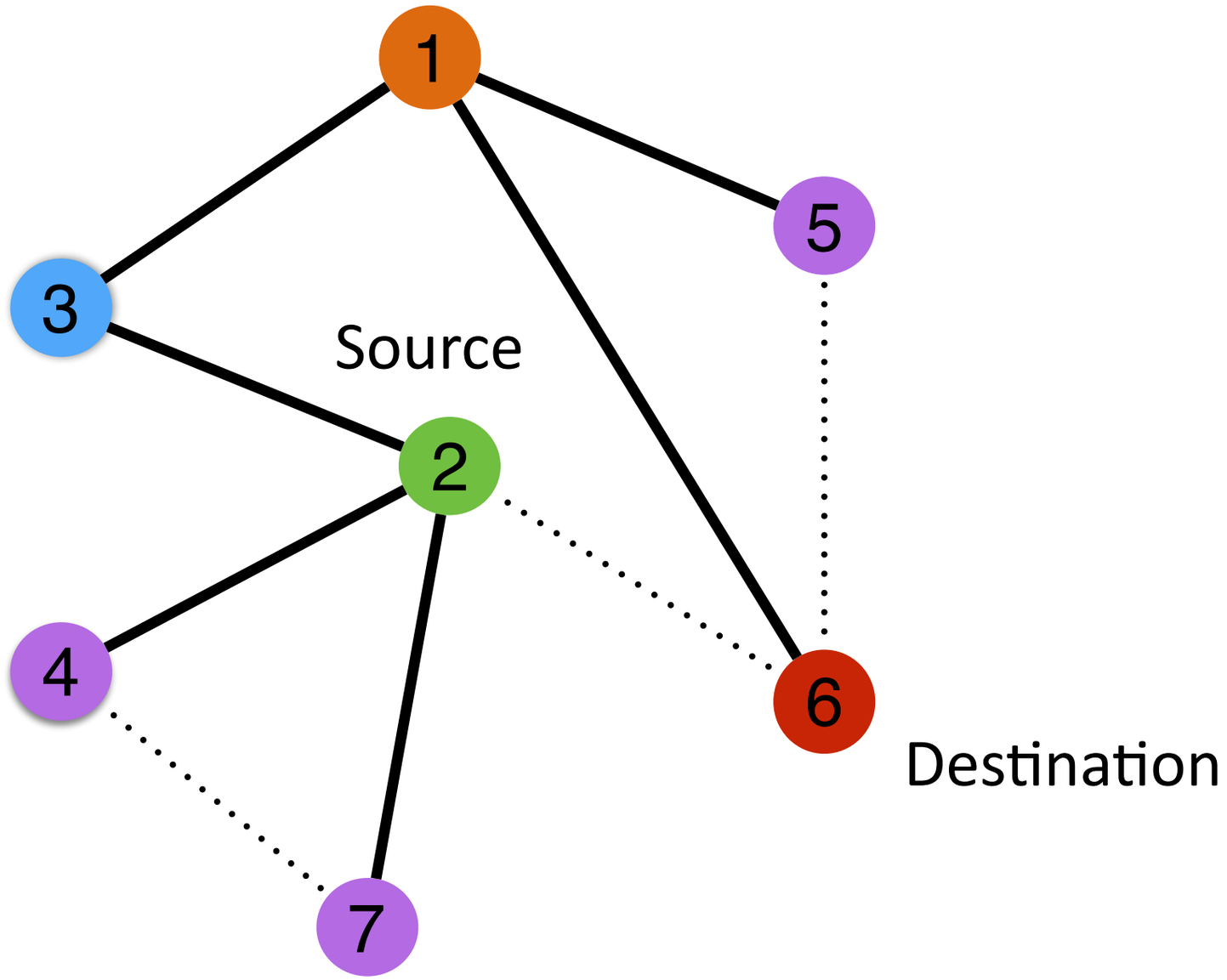
Issue 1: Each host has to do unnecessary packet processing!
(to decide whether the packet is destined to the host)

Recap: Three fundamental issues!



Issue 2: Higher latency!
(The packets unnecessarily traverse much longer paths)

Recap: Three fundamental issues!



Issue 3: Wasted bandwidth!

(2-6 and 3-1 packets unnecessarily have to share bandwidth)

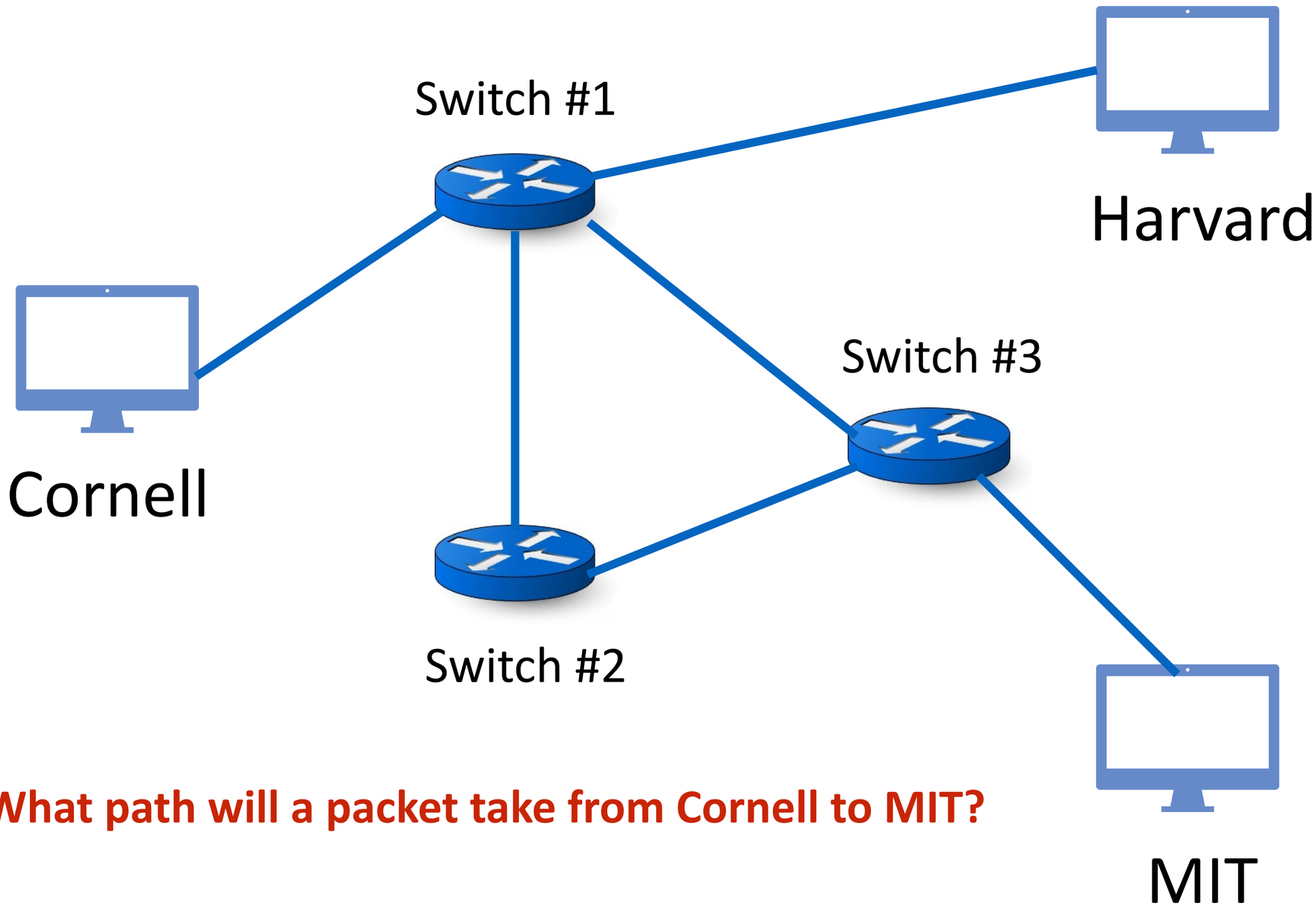
Questions?

Why do we need a network layer?

- Network layer performs “routing” of packets to alleviate these issues
- Uses routing tables
- Lets understand routing tables first

Routing Packets via Routing Tables

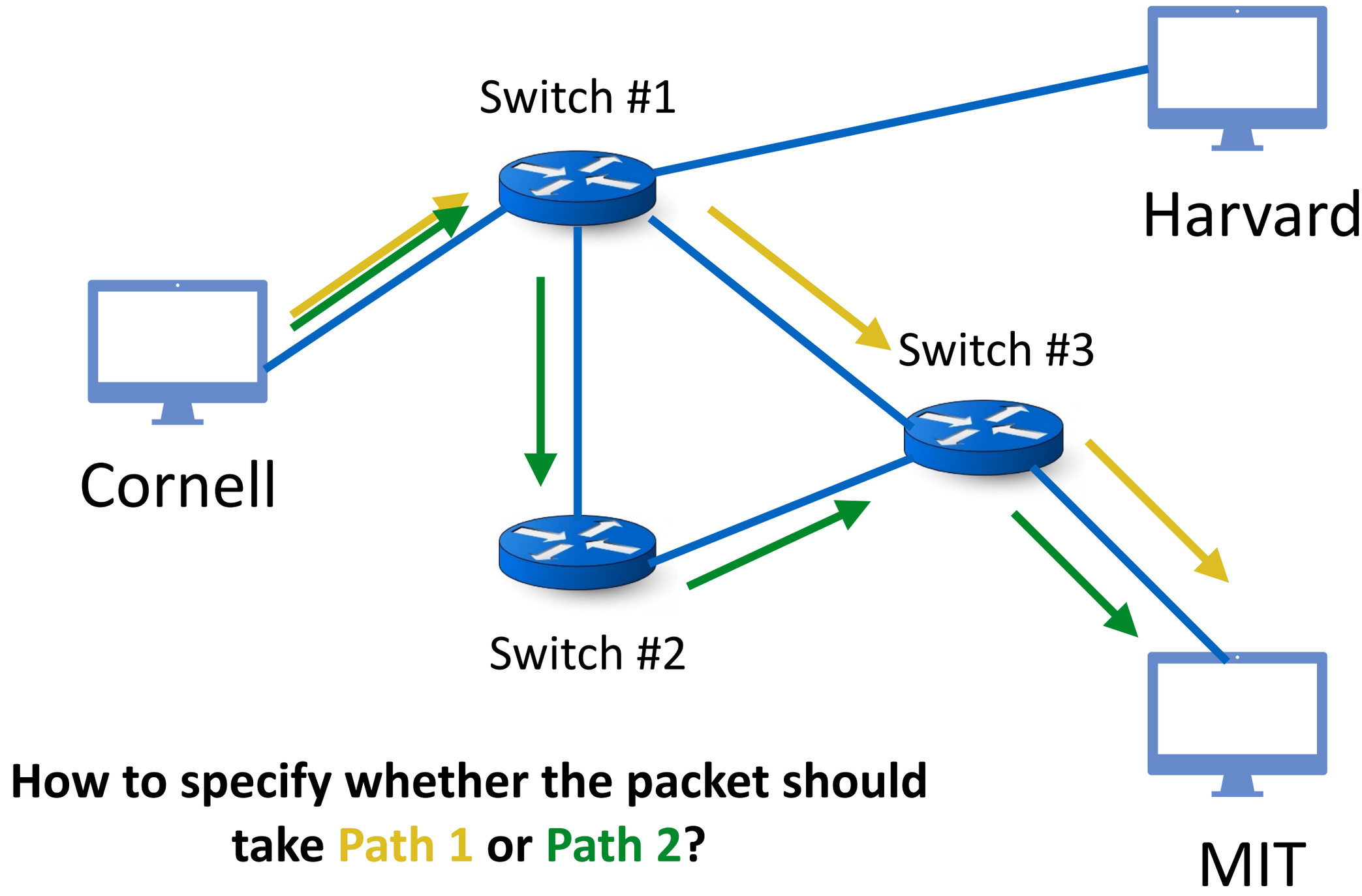
- **Routing tables allow finding path from source to destination**



What path will a packet take from Cornell to MIT?

Routing Packets via Routing Tables

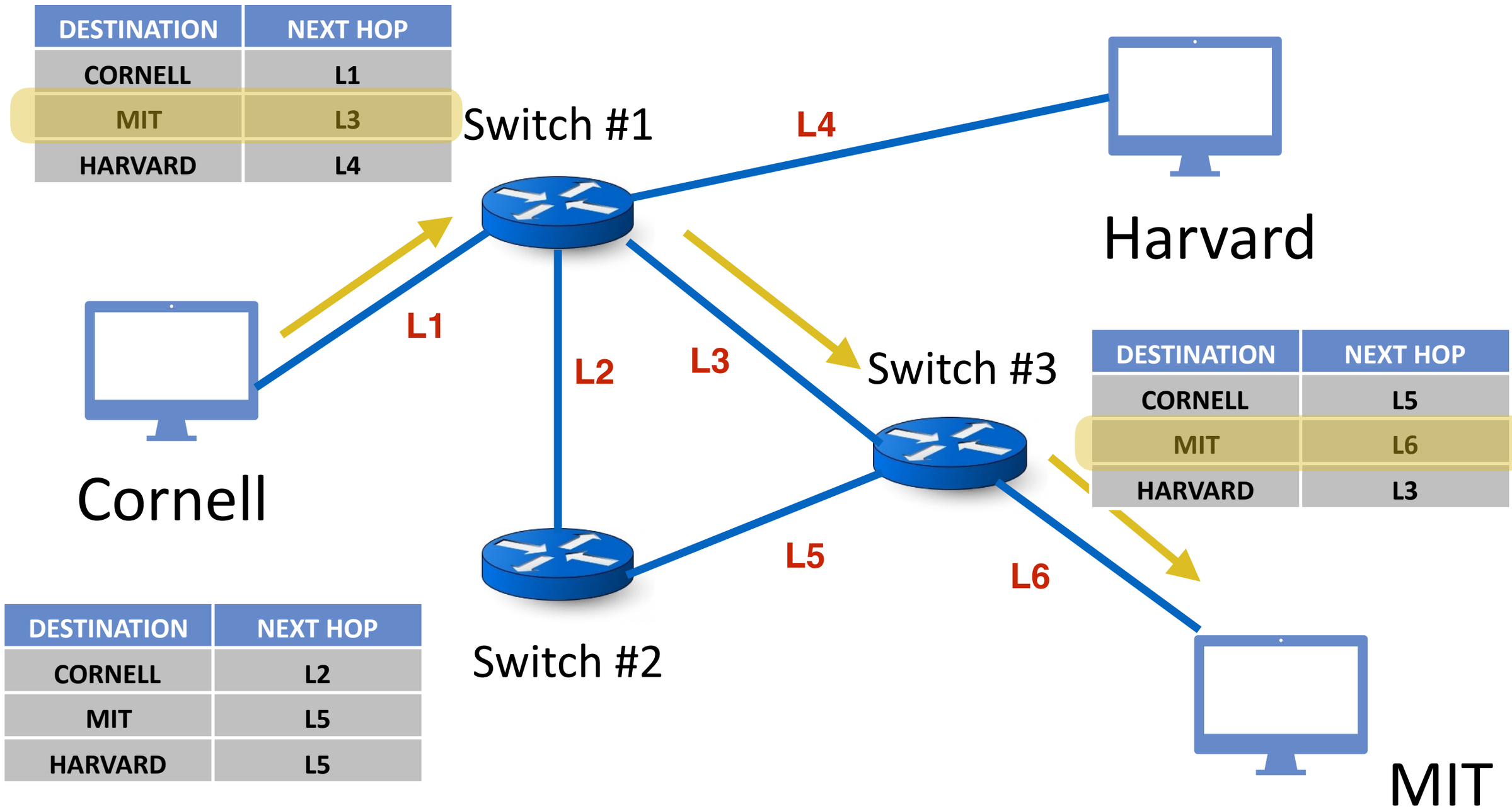
- **Finding path for a packet from source to destination**



How to specify whether the packet should take Path 1 or Path 2?

Routing Table

- **Suppose packet follows Path 1: Cornell - S#1 - S#3 - MIT**

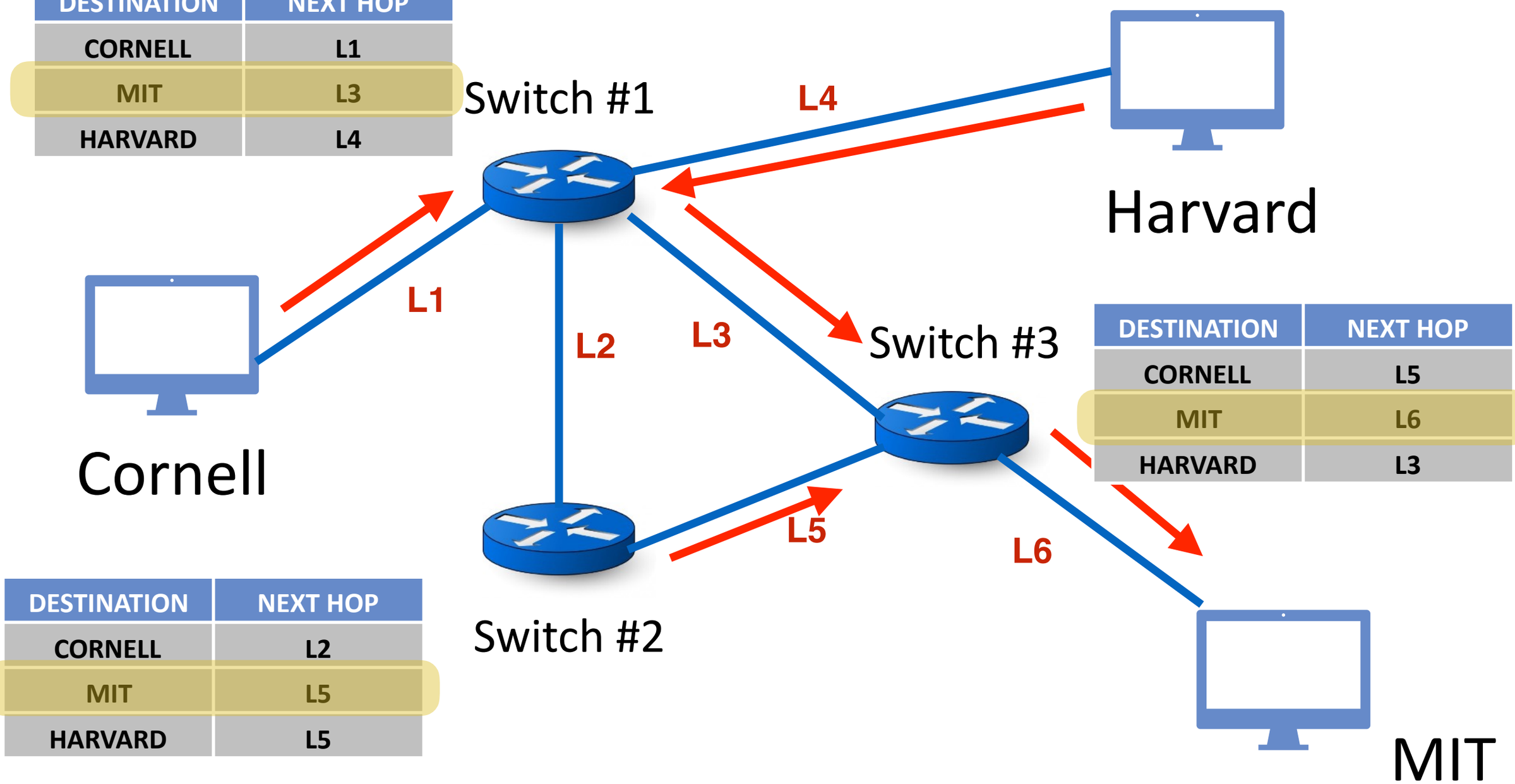


Each Switch stores a table indicating the next hop for corresponding destination of a packet (called a routing table)

Routing Table: The right way to think about them

- Lets focus on one destination - MIT

DESTINATION	NEXT HOP
CORNELL	L1
MIT	L3
HARVARD	L4



DESTINATION	NEXT HOP
CORNELL	L5
MIT	L6
HARVARD	L3

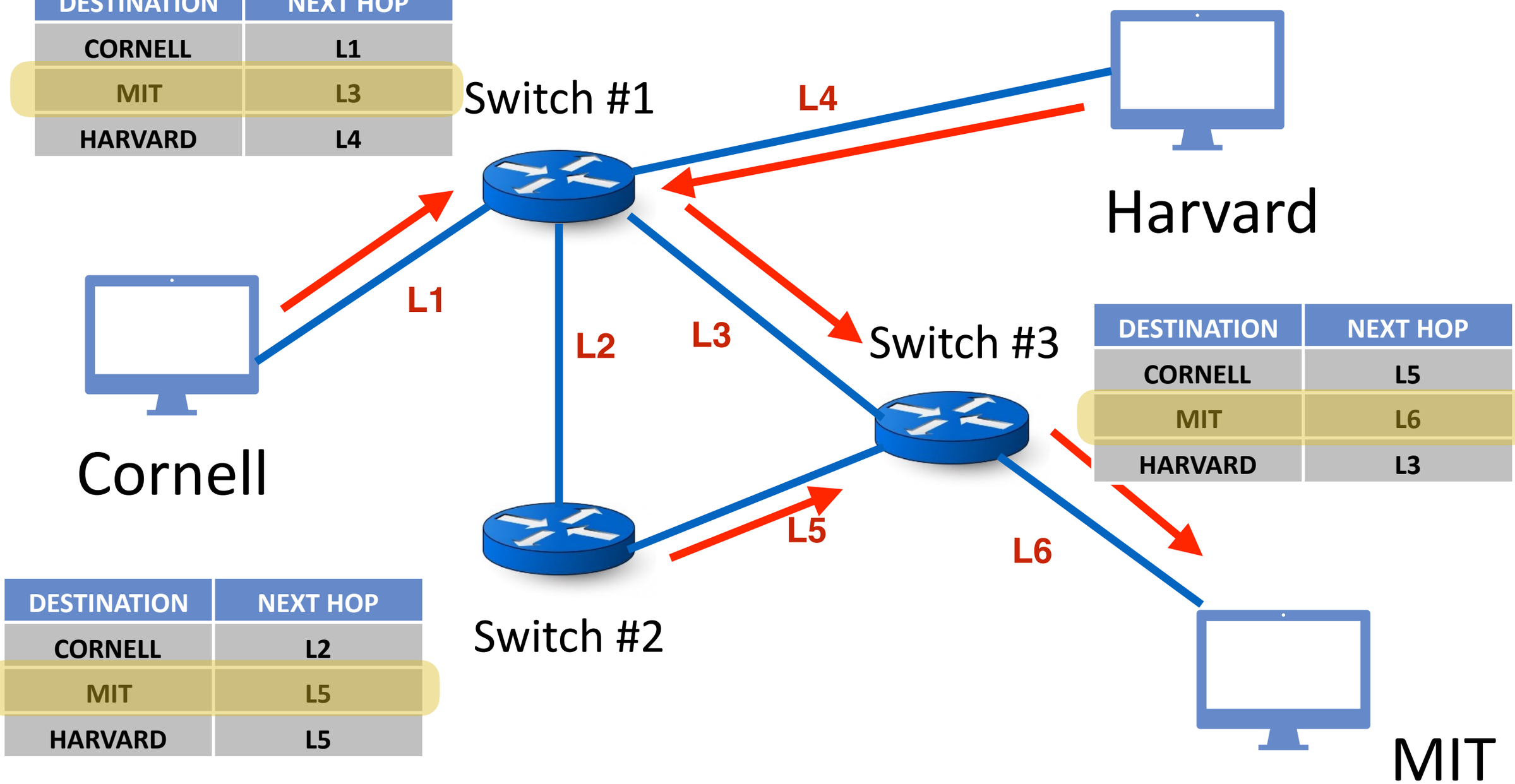
DESTINATION	NEXT HOP
CORNELL	L2
MIT	L5
HARVARD	L5

See something interesting?

Routing Table: The right way to think about them

- Lets focus on one destination - MIT

DESTINATION	NEXT HOP
CORNELL	L1
MIT	L3
HARVARD	L4



DESTINATION	NEXT HOP
CORNELL	L5
MIT	L6
HARVARD	L3

DESTINATION	NEXT HOP
CORNELL	L2
MIT	L5
HARVARD	L5

Routing table entries for a particular destination form a (directed) spanning tree with that destination as the root!!!!

Routing Table: The right way to think about them

- Routing tables are nothing but ….
    - A collection of (directed) spanning tree
    - One for each destination

- **Routing Protocols**
    - “n” spanning tree protocols running in parallel

“Valid Routing Tables” (routing state)

- Global routing state is valid if:
    - it **always** results in deliver packets to their destinations

- **Goal of Routing Protocols**
    - Compute a valid state
    - But how to tell if a routing state is valid?…
    - Think about it, what could make routing incorrect?

Validity of a Routing State

- Global routing state valid **if and only if**:
    - There are no **dead ends** (other than destination)
    - There are no **loops**

- A **dead end** is when there is **no outgoing link**
    - A packet arrives, but ..
        - the routing table does not have an outgoing link
        - And that node is not the destination

- A **loop** is when a **packet cycles around** the same set of nodes forever

Example: Routing with Dead Ends

- Suppose packet wants to go from Cornell to MIT using given state:

DESTINATION	NEXT HOP
CORNELL	L1
MIT	L2
HARVARD	L4

Switch #1

L4

Harvard

L1

L2

L3

Switch #3

DESTINATION	NEXT HOP
CORNELL	L5
MIT	L6
HARVARD	L3

Cornell

L5

L6

Switch #2

DESTINATION	NEXT HOP
CORNELL	L2
HARVARD	L5

Dead End!

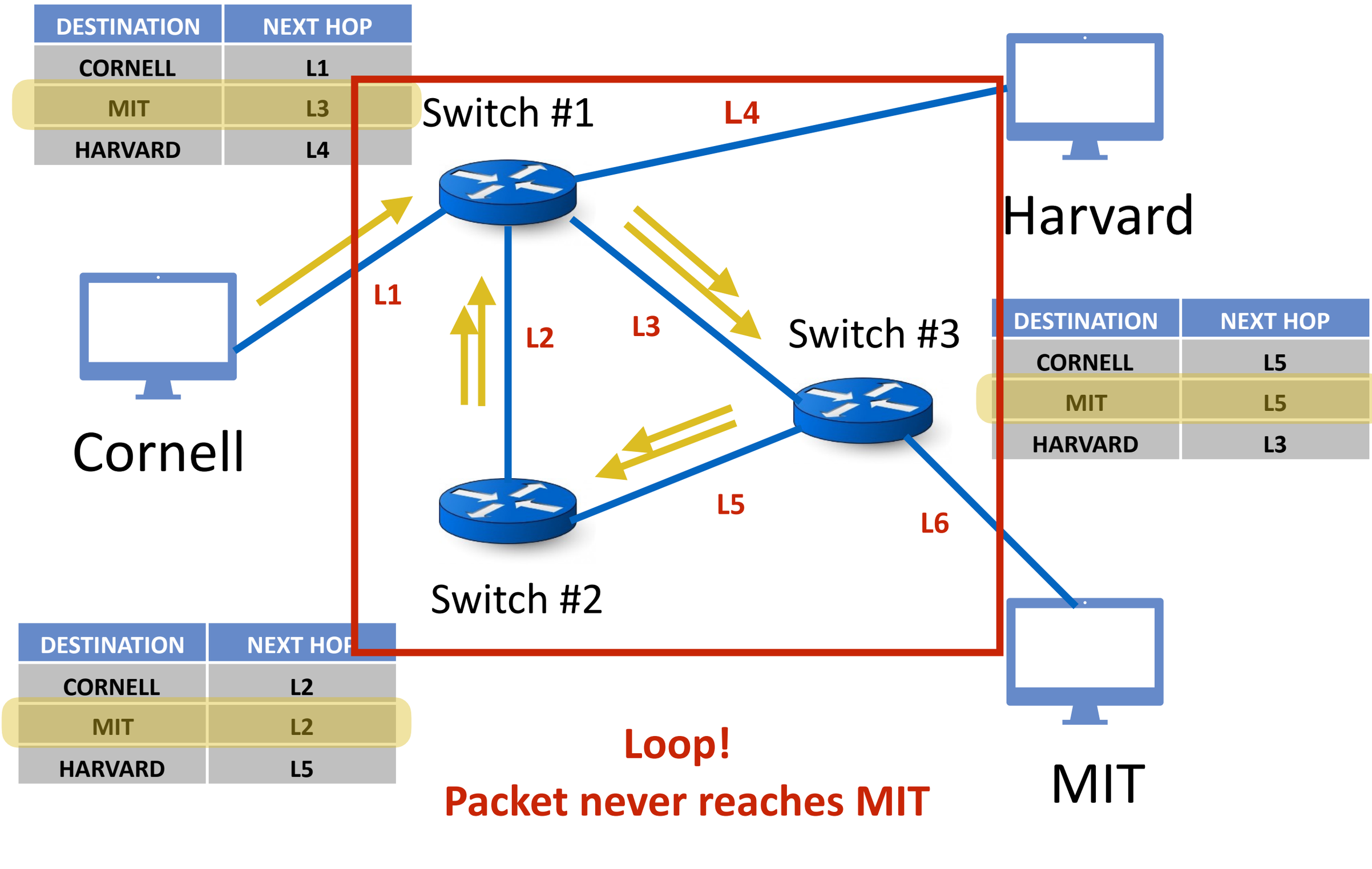
Packet never reaches MIT

MIT

No forwarding decision for MIT!

Example: Routing with Loops

- Suppose packet wants to go from Cornell to MIT using given state:



Loop!

Packet never reaches MIT

Two Questions

- How can we **verify** given routing state is valid?

- How can we **produce** valid routing state?

Checking Validity of a Routing State

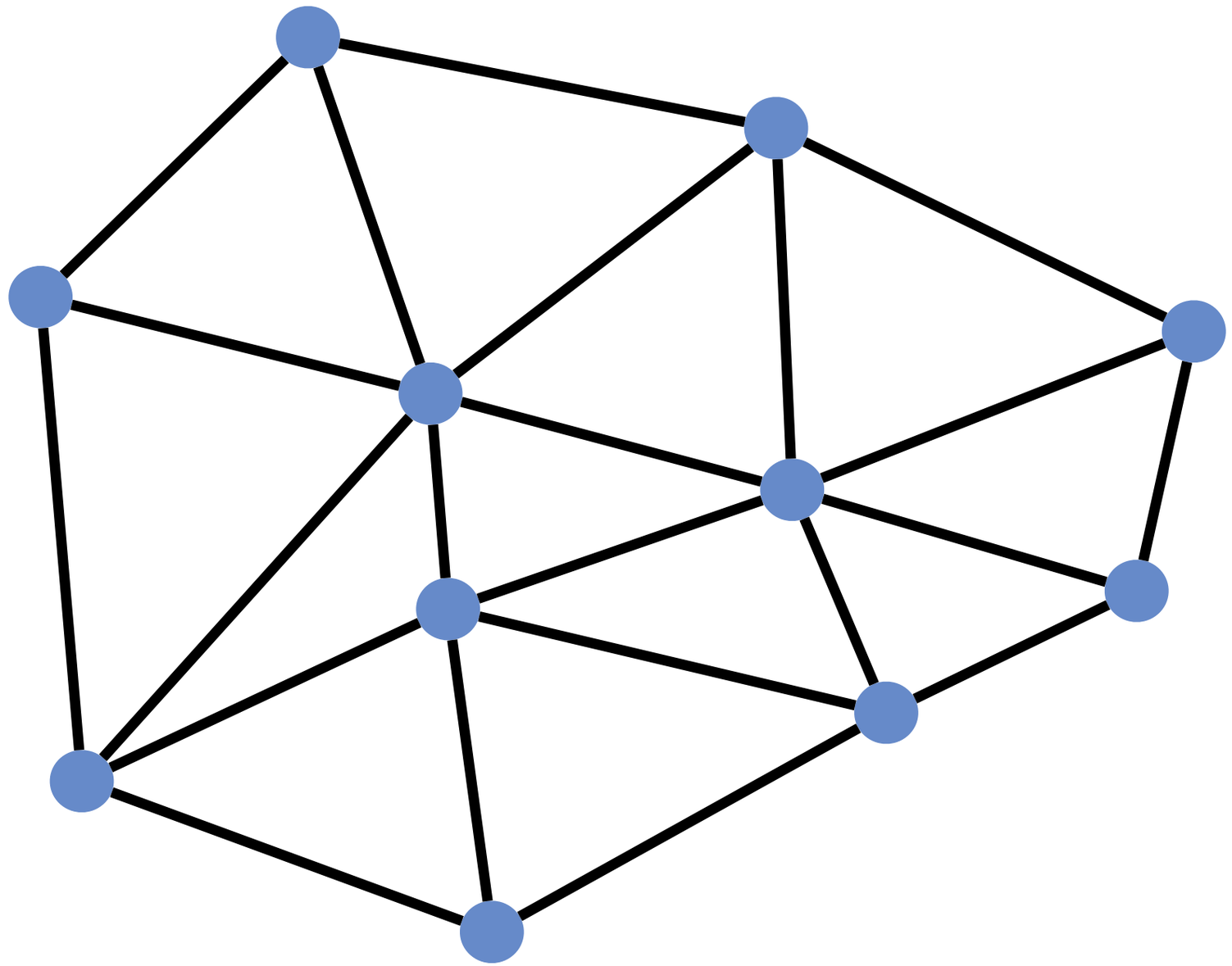
- Check validity of routing state for one destination at a time…

- For each node:
    - Mark the outgoing link with arrow for the required destination
    - There can only be one at each node

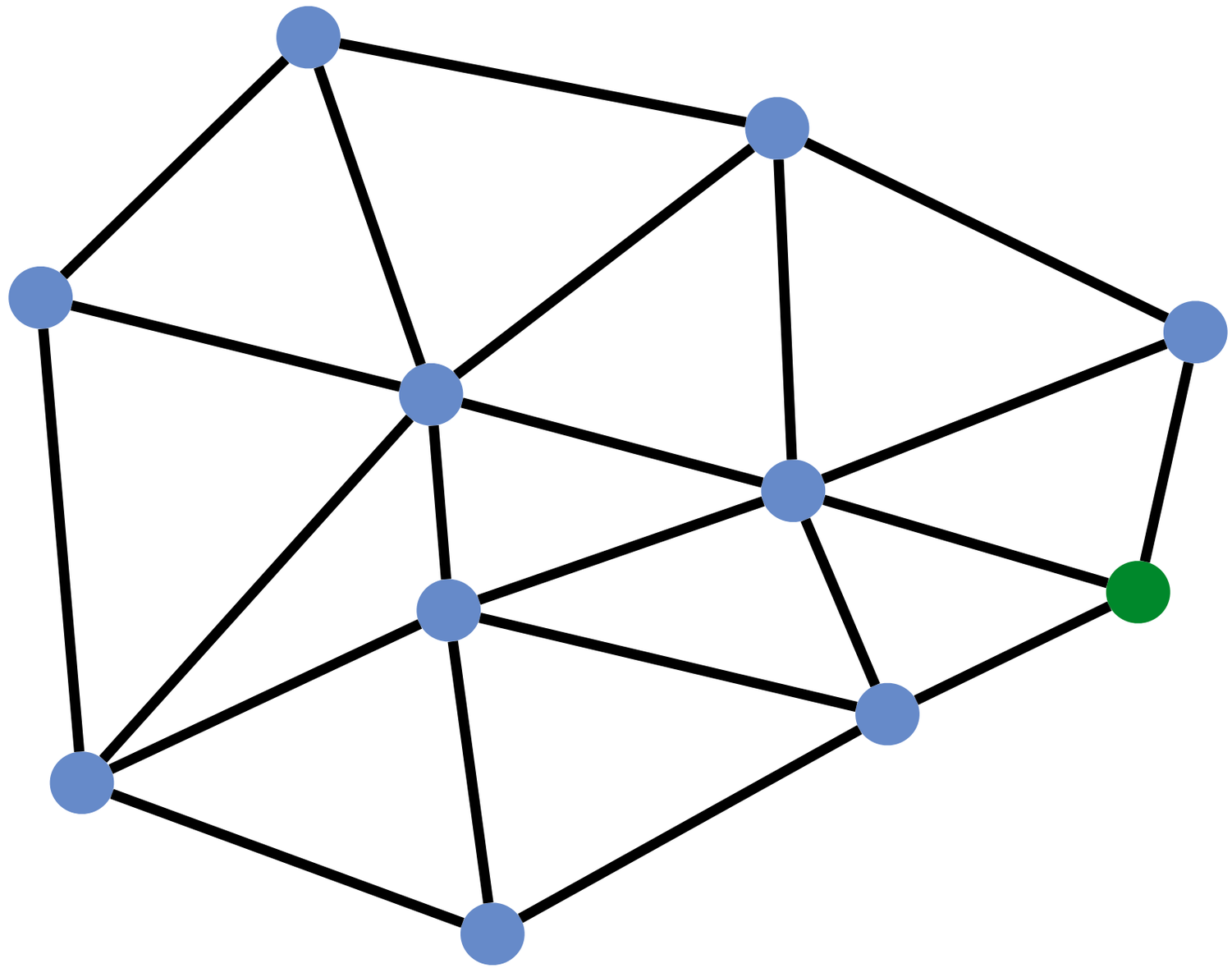
- Eliminate all links with no arrows

- Look what's left. **State is valid if and only if**
    - Remaining graph is a spanning tree with destination as sink
    - Why is this true?
        - Tree -> No loops
        - Spanning (tree) -> No dead ends

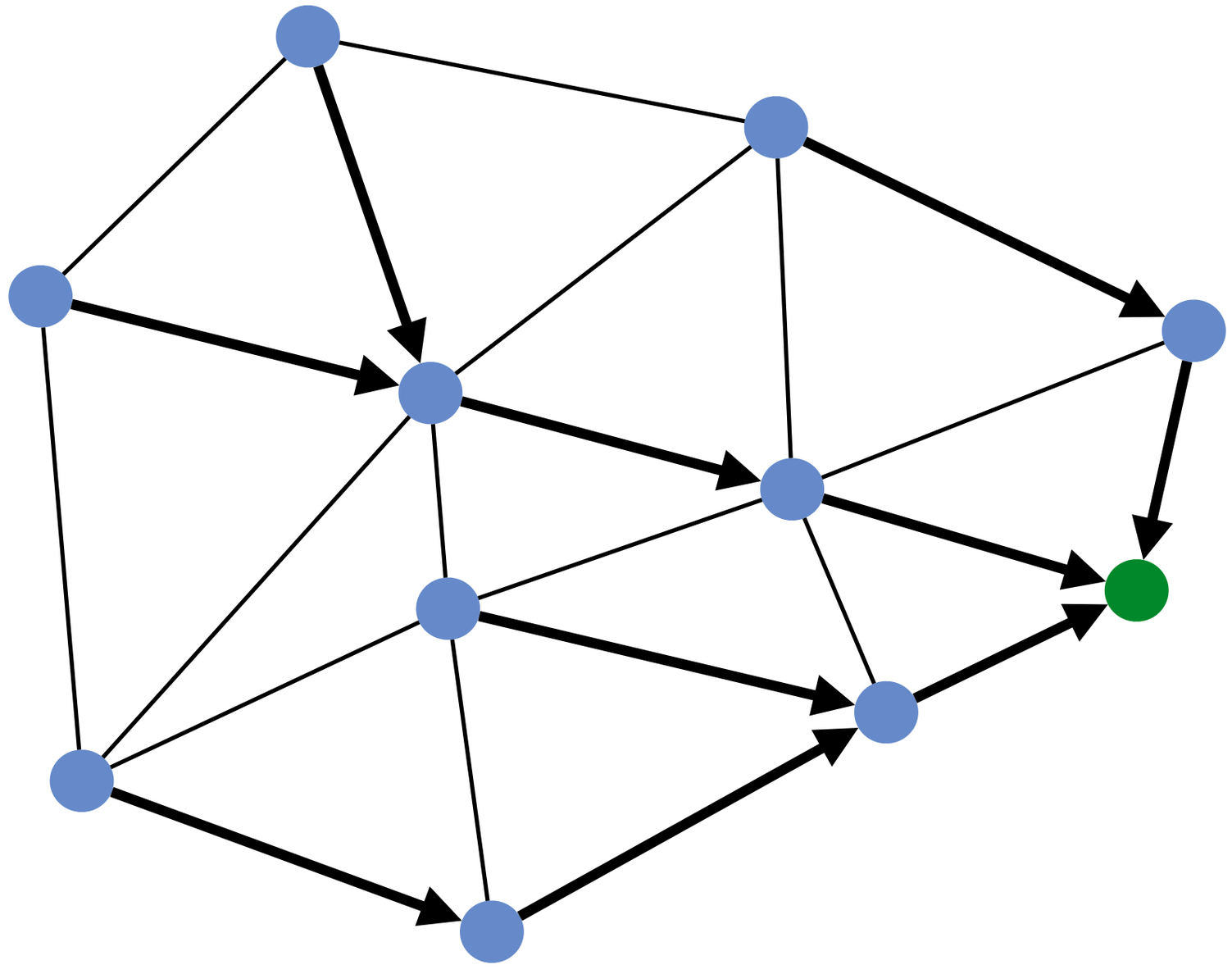
Example 1



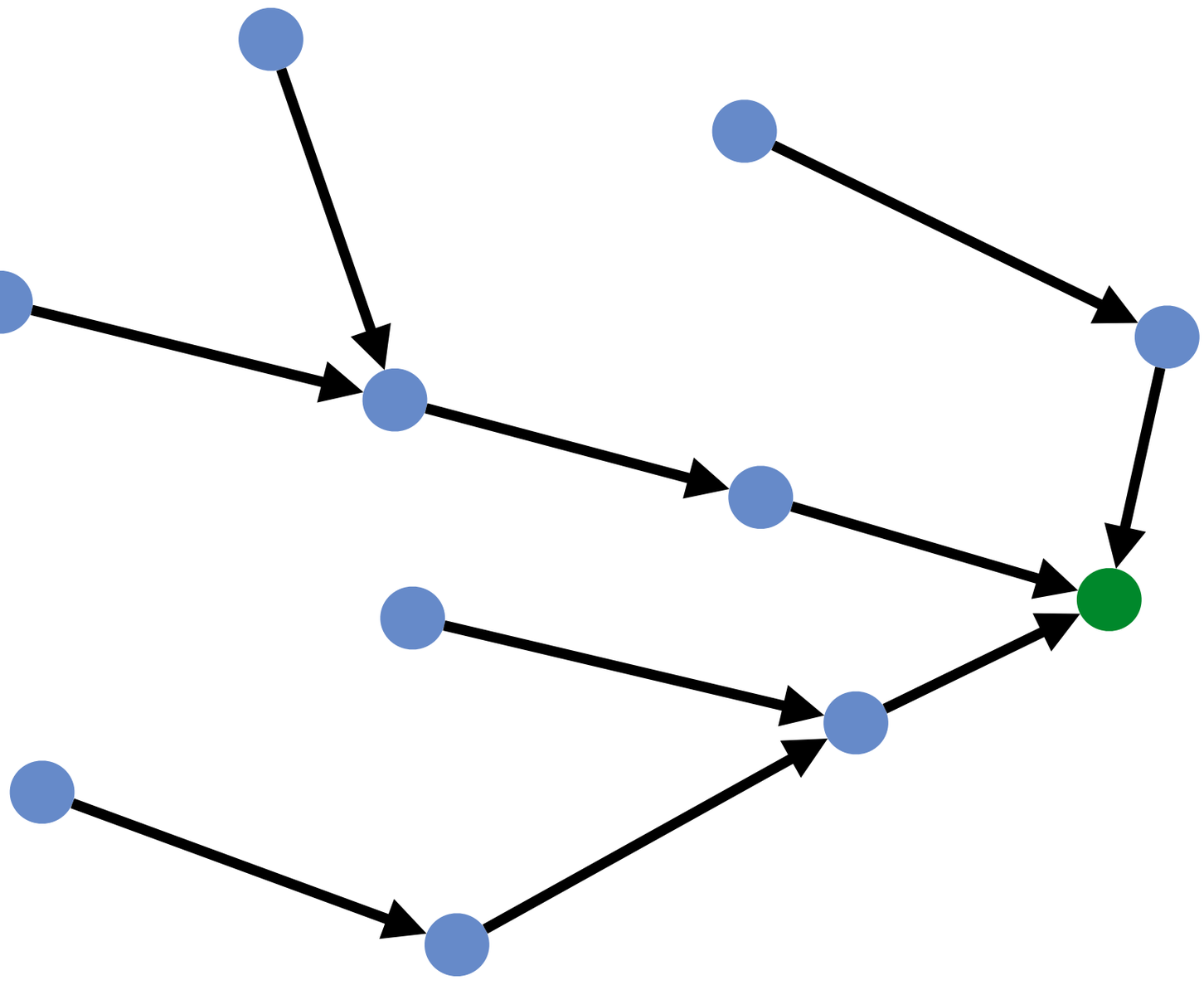
Example 1: Pick Destination



Example 1: Put Arrows on Outgoing Ports

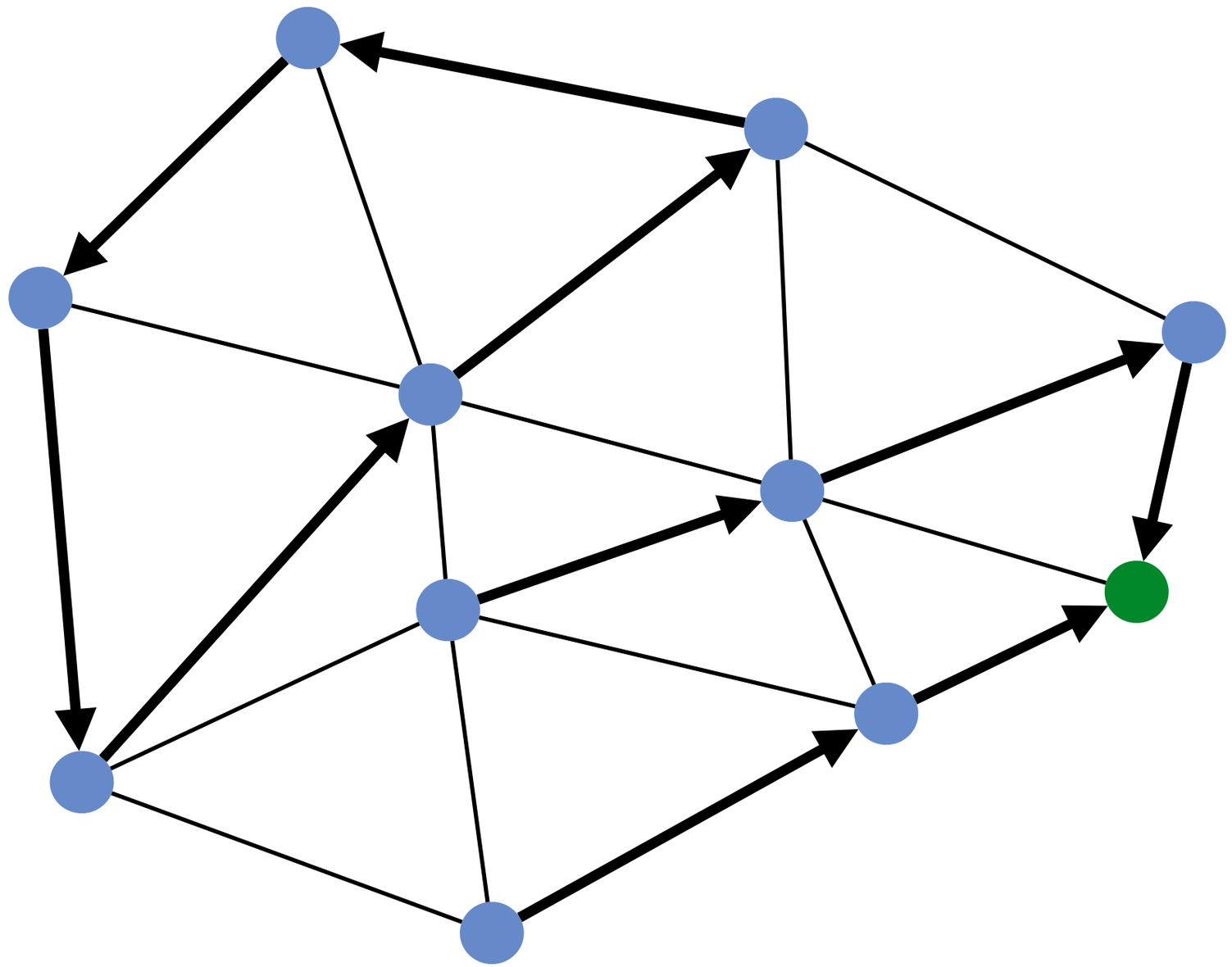


Example 1: Remove unused Links

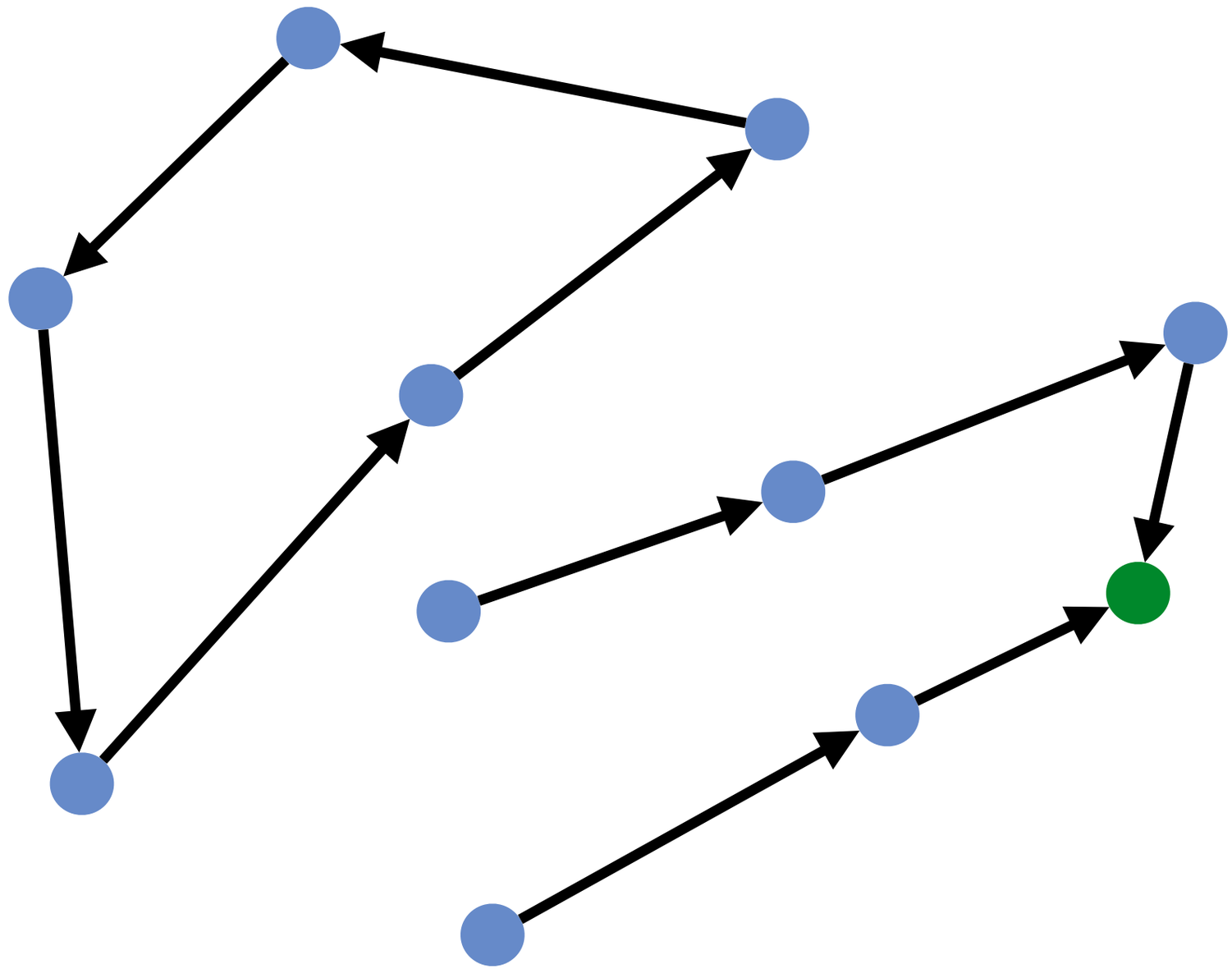


Leaves Spanning Tree: Valid

Example 2:

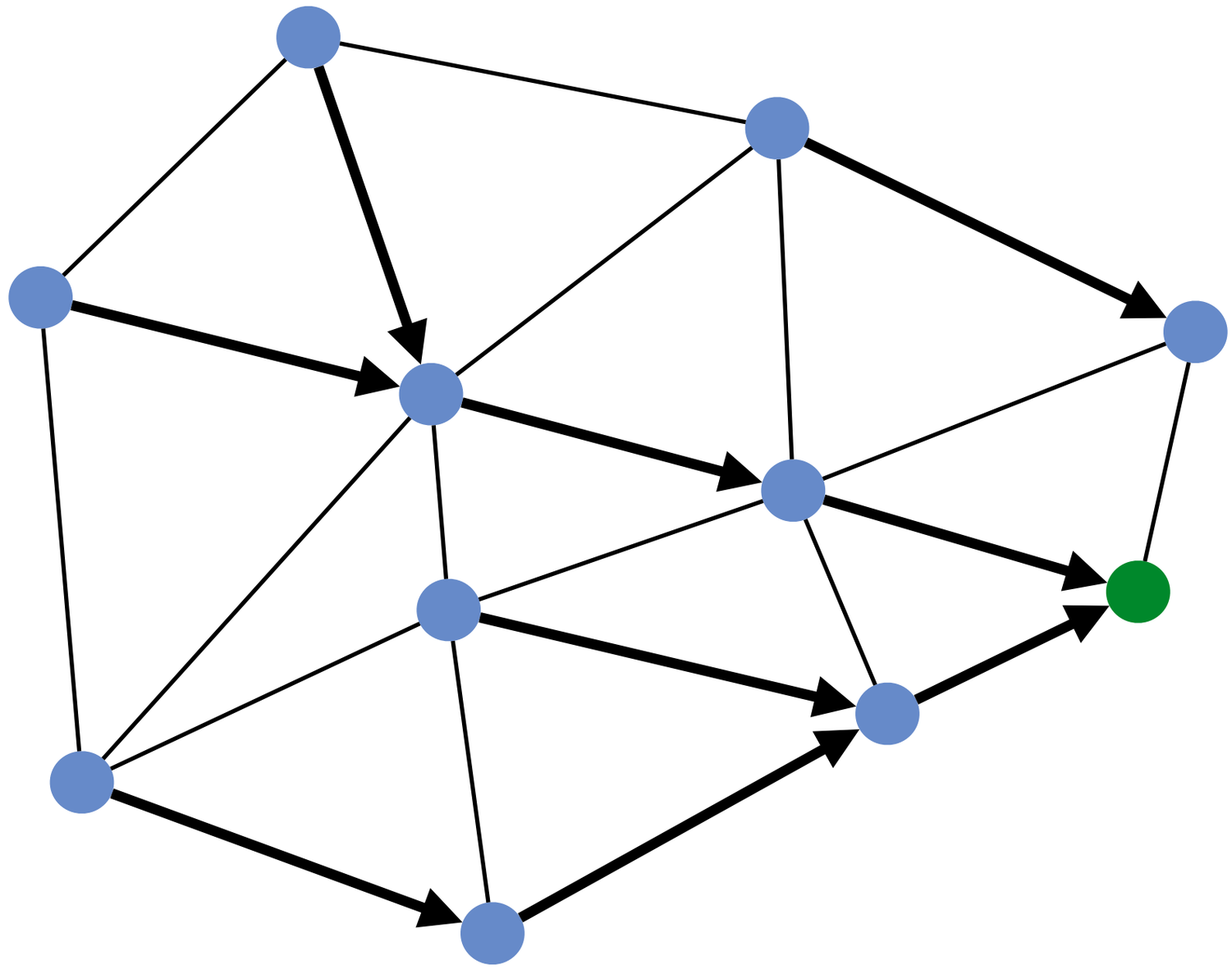


Example 2:

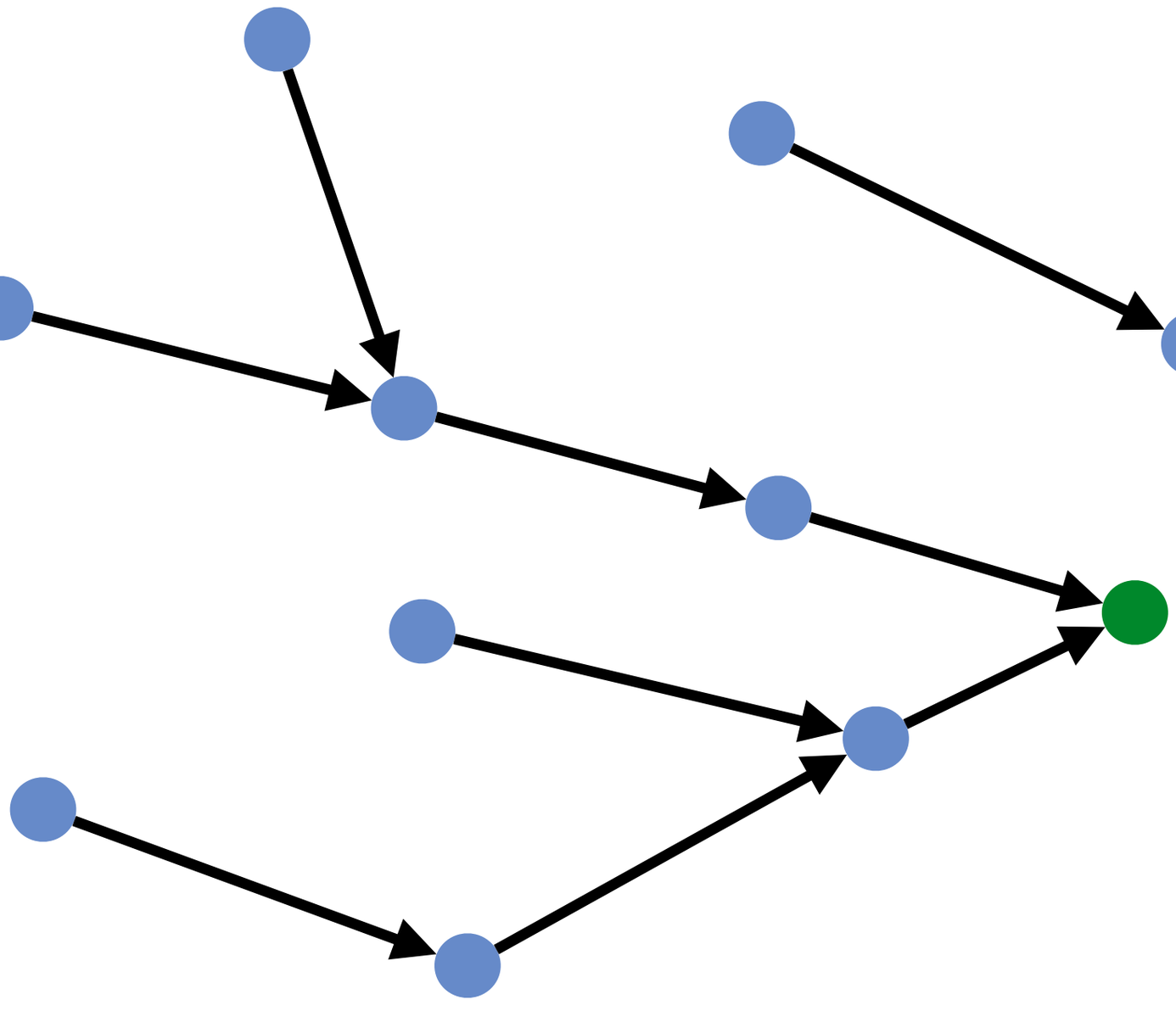


Is this valid?

Example 3:



Example 3:



Is this valid?

Checking Validity of a Routing State

- Simple to check validity of routing state for a particular destination
- Dead ends: nodes without arrows
- Loops: obvious, disconnected from destination and rest of the graph

Two Questions

- How can we **verify** given routing state is valid?

- How can we **produce** valid routing state?

Creating Valid Routing State

- Easy to avoid dead ends
- Avoiding loops is hard
- **The key difference between routing protocols is how they avoid loops!**
- Try to think a loop avoidance design for five minutes

Four flavors of protocols

- **Create Tree, route on tree**
    - E.g., Spanning tree protocol (switched Ethernet)
    - **Good:** easy, no (persistent) loops, no dead ends
    - **Not-so-good:** unnecessary processing, high latency, low bandwidth

- **Obtain a global view:**
    - E.g., Link state

- **Distributed route computation:**
    - E.g., Distance vector
    - E.g., Border Gateway Protocol

Routing Metrics

- Routing goals: compute paths with minimum X
    - X = number of “hops” (nodes in the middle)
    - X = latency
    - X = weight
    - X = failure probability
    - …
- Generally assume every link has “cost” associated with it

- We want to minimize the cost of the entire path
    - **We will focus on a subset of properties X, where:**
    - **Cost of a path = sum of costs of individual links/nodes on the path**
    - E.g., number of hops and latency

#1: Create a Tree

#1: Create a Tree Out of Topology

- Remove enough links to create a tree containing all nodes
- Sounds familiar? Spanning trees!
- If the topology has no loops, then just make sure not sending packets back from where they came
    - That causes an immediate loop
- Therefore, if no loops in topology and no formation of immediate loops ensures valid routing
- However… three challenges
    - Unnecessary host resources used to process packets
    - High latency
    - Low bandwidth (utilization)

Global view

Two Aspects of Global View Method

- **Protocol**: What we focus on today
    - Where to create global view
    - How to create global view
    - Disseminating route computation (if necessary)
    - When to run route computation

- **Algorithm**: computing loop-free paths on graph
    - Straightforward to compute lowest cost paths
        - Using Dijkstra's algorithm (please study; algorithms course)
    - We won't spend time on this

Where to create global view?

- One option: Central server
    - Collects a global view
    - Computes the routing table for each node
    - “Installs” routing tables at each node
    - **Software-defined Networks: later in course**

- Second option: At each router
    - Each router collects a global view
    - Computes its own routing table using Link-state protocol

- **Link-state routing protocol**
    - OSPF is a specific implementation of link-state protocol
        - IETF RFC 2328 (IPv4) or 5340 (IPv6)

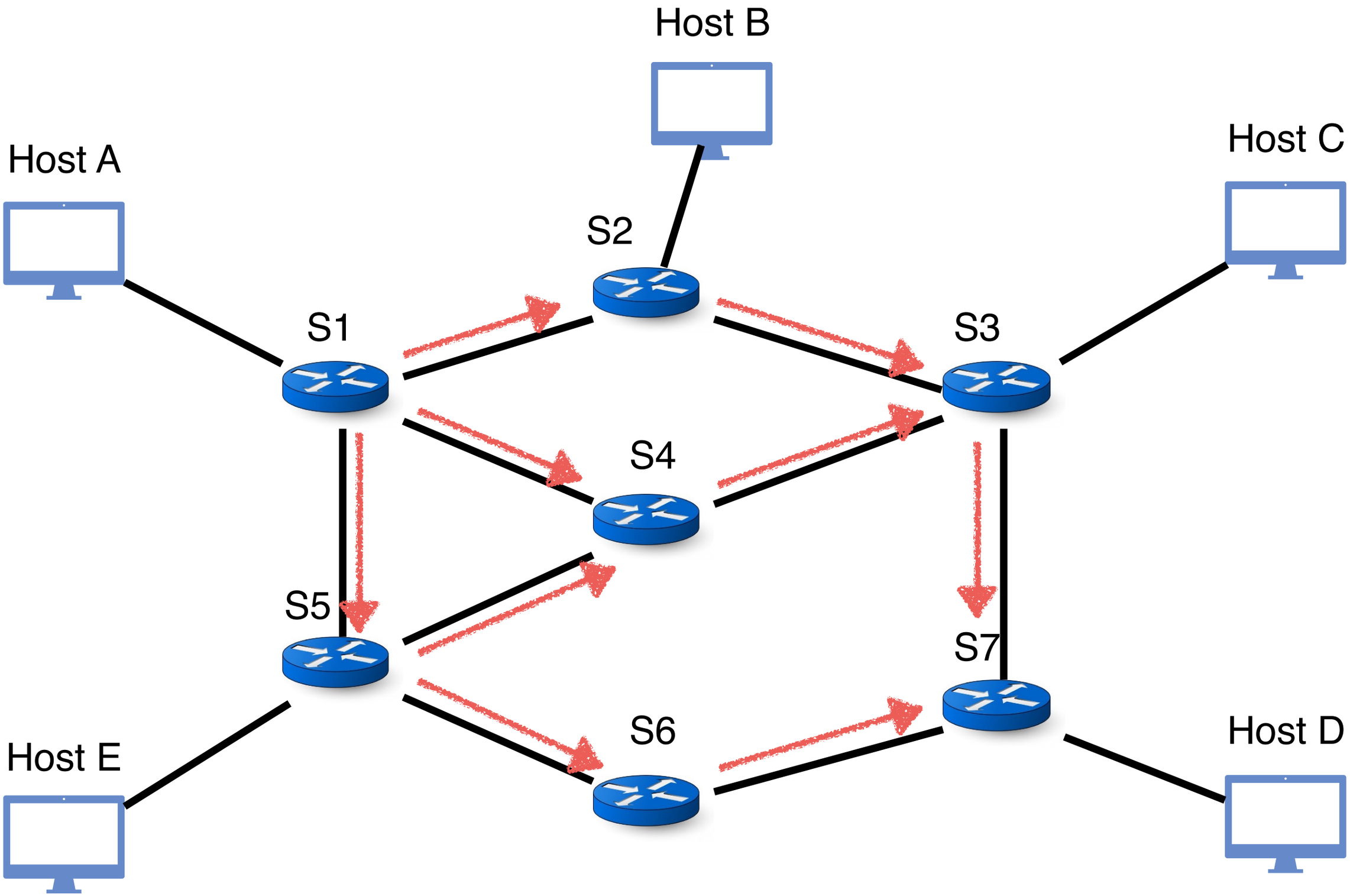
Overview of Link-State Routing

- **Every router knows its local "link state"**
    - Knows state of links to neighbors
    - Up/down, and associated cost
- **A router floods its link state to all other routers**
    - Uses a special packet — Link State Announcements (LSA)
    - Announcement is delivered to all nodes (next slide)
    - Hence, every router learns the entire network graph
- **Runs route computation locally**
    - Computing least cost paths from them to all other nodes
    - E.g., using Dijkstra's algorithm

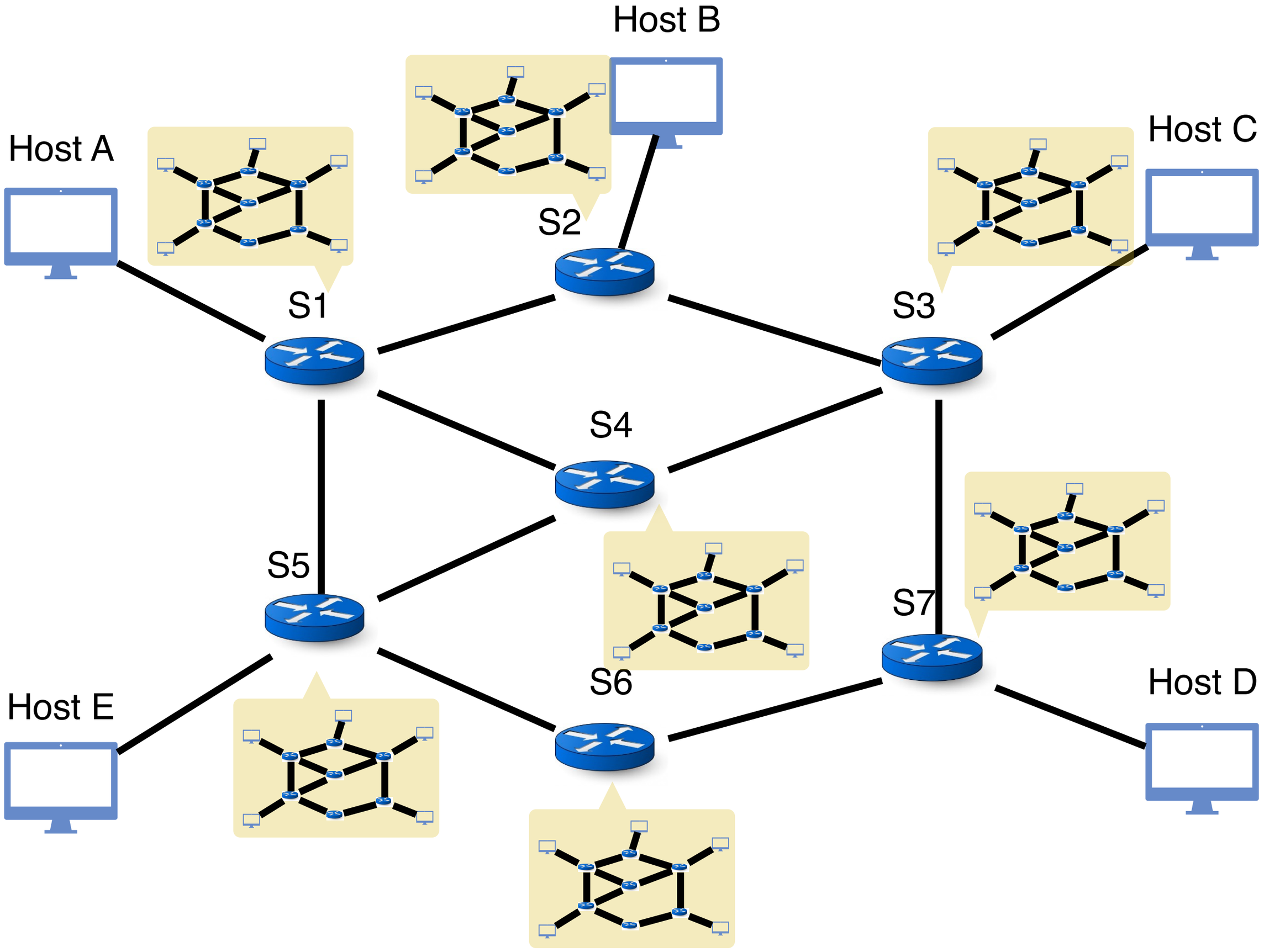
How does Flooding Work?

- "Link state announcement" (LSA) arrives on a link at a router
- That router:
    - Remembers the packet
    - Forwards the packet out all **other links**
    - Does **not** send it out the incoming link
        - Why?
- If a previously received announcement arrives again…
    - Router drops it (no need to forward again)

Link-State Routing



Each Node Then has a Global View



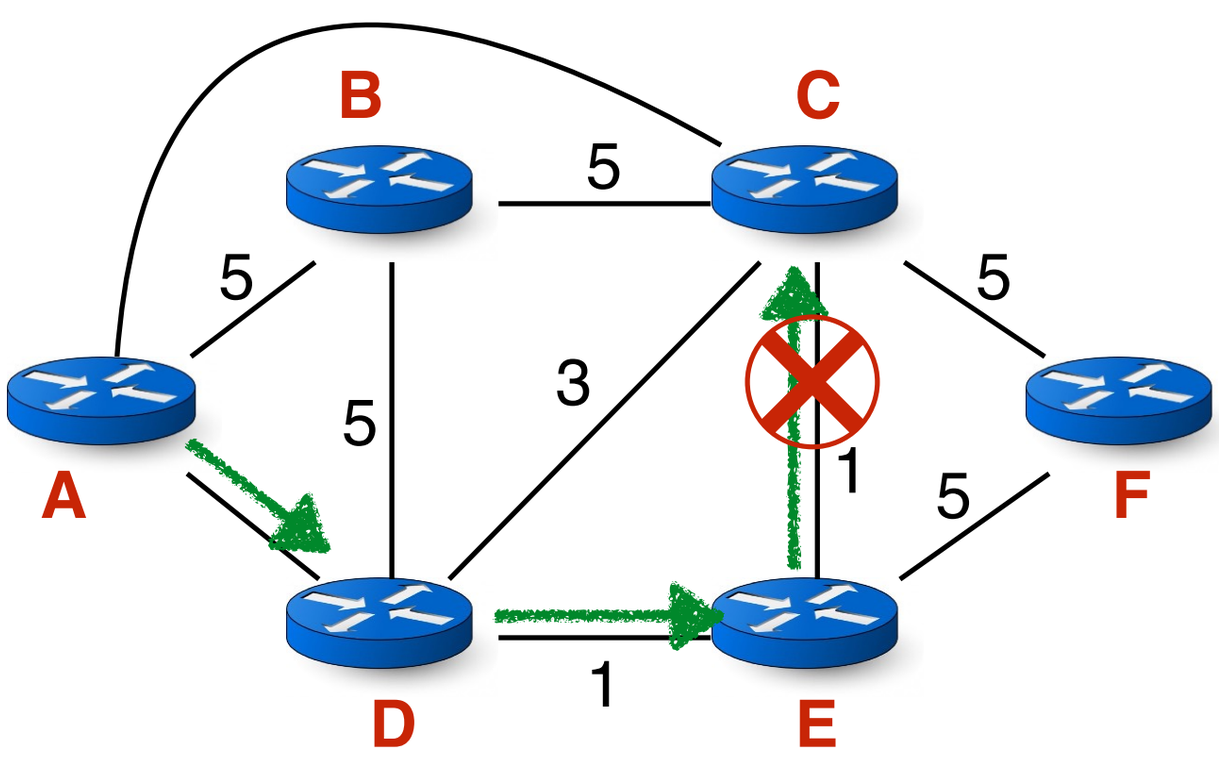
When to Initiate Flooding of announcements?

- **Topology change**
    - Link failures
    - Link recovery
- **Configuration change**
    - Link cost change (why would one change link cost?)
- **Periodically**
    - Refresh the link-state information
    - Typically (say) 30 minutes
    - Corrects for possible corruption of data

Making Floods Reliable

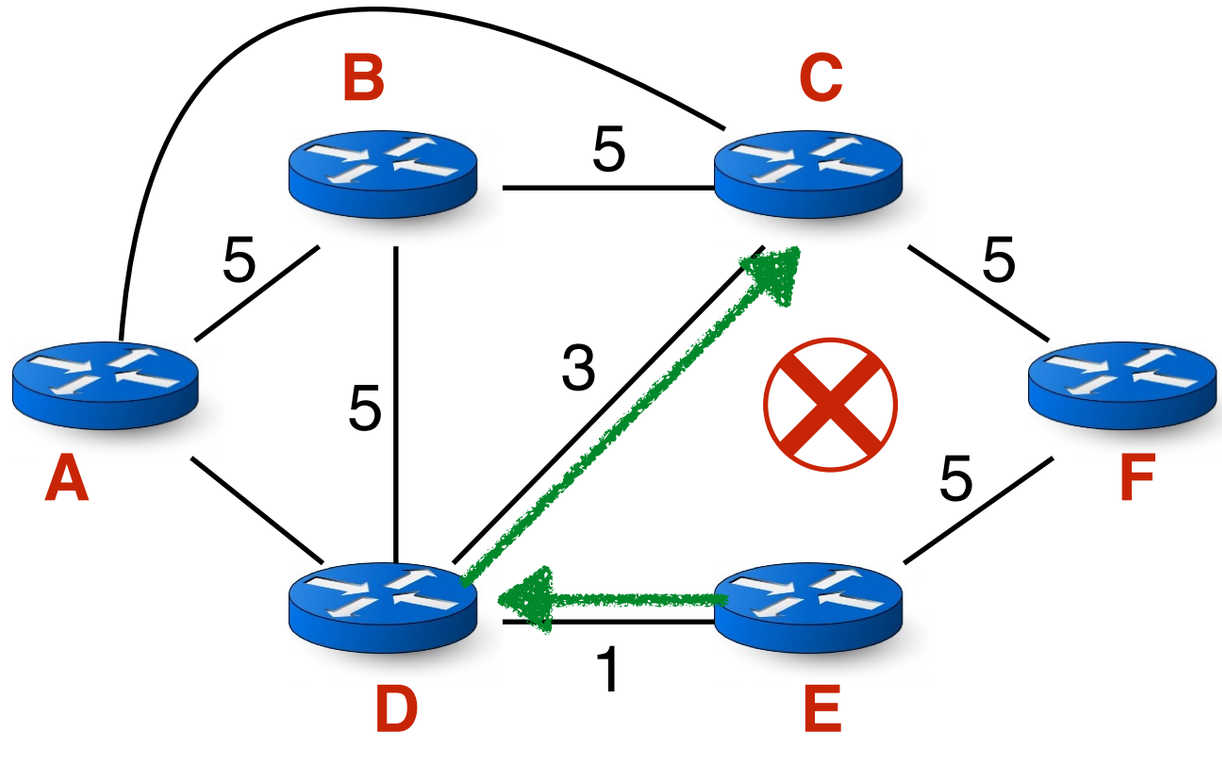
- Reliable Flooding
    - Ensure all nodes receive same link state announcements
        - No announcements dropped
    - Ensure all nodes use the latest version
- Suppose we can implement reliable flooding. How can it still fail?
- Can you ever have loops with link-state routing?
- **Again: Can you ever have loops with link-state routing?**

Are Loops Still Possible?



A and D think this is the path to C

E-C link fails, but D doesn't know yet

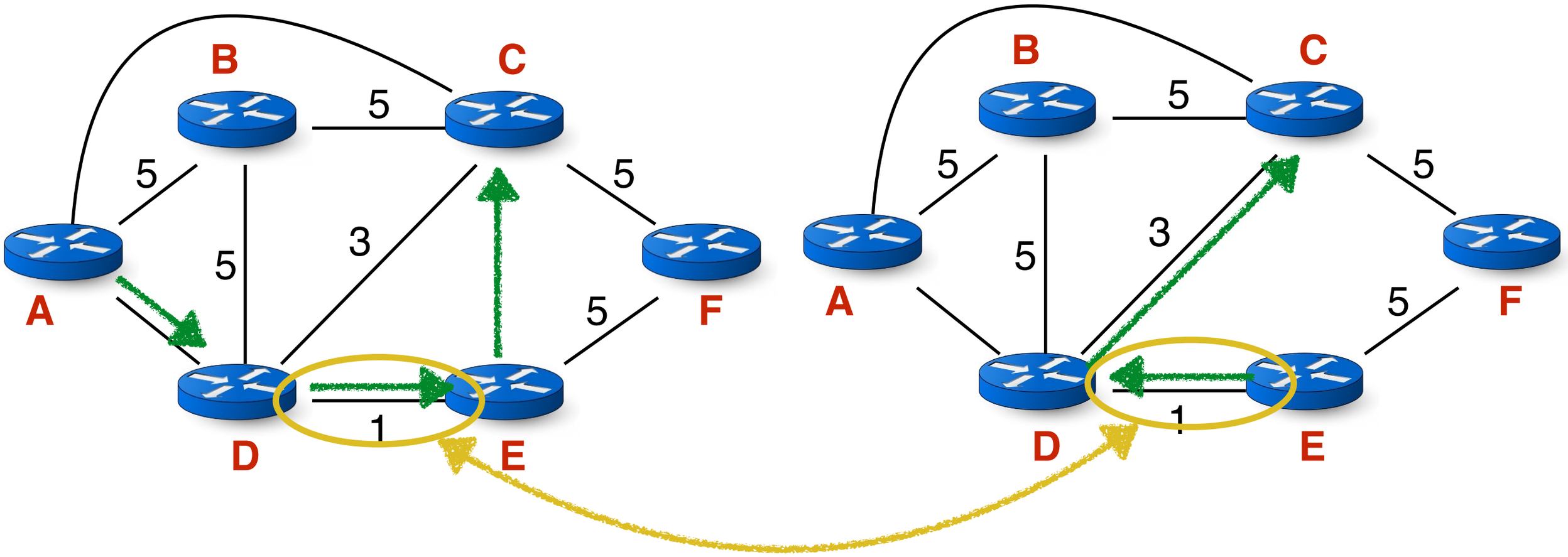


E thinks that this the path to C

E reaches C via D, D reaches C via E

Loop!

Transient Disruptions



- Inconsistent link-state views
  - Some routers know about failure before others
  - The shortest paths are no longer consistent
  - Can cause **transient forwarding loops**
    - **Transient loops are still a problem!**

Convergence

- All routers have consistent routing information
    - E.g., all nodes having the same link-state database
- Forwarding is consistent after convergence
    - All nodes have the same link-state database
    - All nodes forward packets on same paths
- **But while still converging, bad things can happen**

Time to Reach Convergence

- Sources of convergence delay?
    - Time to detect failure
    - Time to flood link-state information (~longest RTT)
    - Time to recompute forwarding tables
- Performance problems during convergence period?
    - Dead ends
    - Looping packets
    - And some more we'll see later ….

Link State is Conceptually Simple

- Everyone floods links information
- Everyone then knows graph of the network
- Everyone independently computes paths on the graph
- **All the complexity is in the details**

#3: Distributed Route Computation

- Often getting a global view of the network is infeasible
    - Distributed algorithms to compute feasible route

- **Approach A**: Finding optimal route for maximizing/minimizing a metric

- **Approach B**: Finding feasible route via exchanging paths among switches

Next lecture! Happy February Break!!