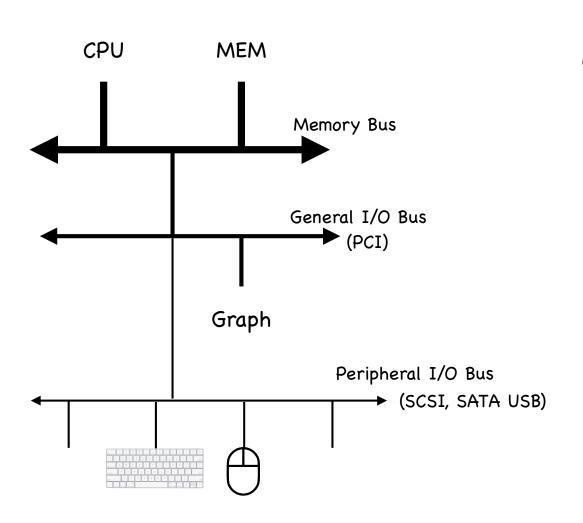
I/O Devices

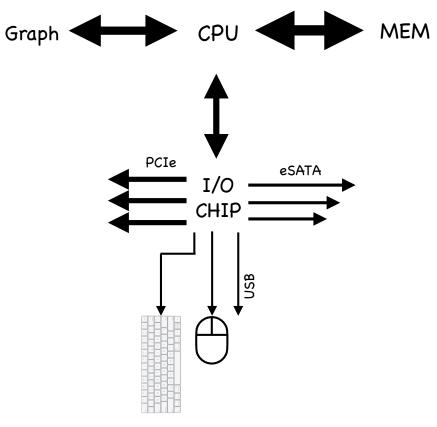
You Need to Get Out More!

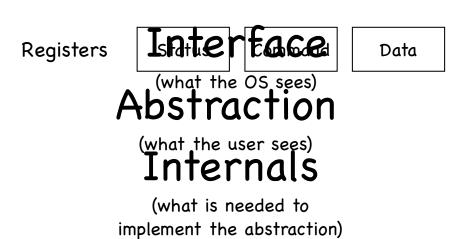
How does a computer connect with the outside world?



I/O Architecture











Data

(.....,

Internals

(what is needed to implement the abstraction)

Registers

Status

Command

Data

Microcontroller

Memory

Internals

Other device (what is needed to specific chips implement the abstraction)

Registers

Status

Command

Data

Microcontroller

Memory

Internals

Other device (what is needed to specific chips implement the abstraction)

 OS controls device by reading/writing registers

```
while (STATUS == BUSY)
  ; // wait until device is not busy
write data to DATA register
write command to COMMAND register
  // starts device and executes command
while (STATUS == BUSY)
  ; // wait until device is done with request
```

Tuning It Up

CPU is polling

```
use interrupts
```

run another process while device is busy

what if device returns very quickly?

CPU is copying all the data to and from DATA

use Direct Memory Access (DMA)

```
while (STATUS == BUSY)
  ; // wait until device is not busy
write data to DATA register
write command to COMMAND register
  // starts device and executes command
while (STATUS == BUSY)
  ; // wait until device is done with request
```

From interrupt-driven I/O to DMA

Interrupt driven I/O

Device CPU RAM

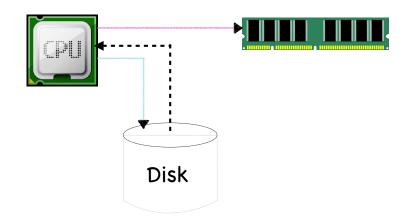
for

CPU issues read request

device interrupts CPU with

data

CPU writes data to memory



From interrupt-driven I/O to DMA

Interrupt driven I/O

Device ◆ ► CPU ◆ ► RAM

for

CPU issues read request

device interrupts CPU with

data

CPU writes data to memory

+ Direct Memory Access

Device **4**

▶RAM

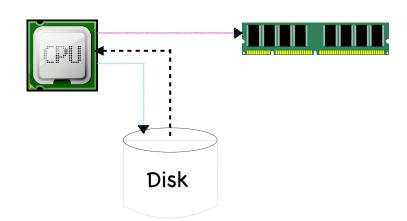
CPU sets up DMA request

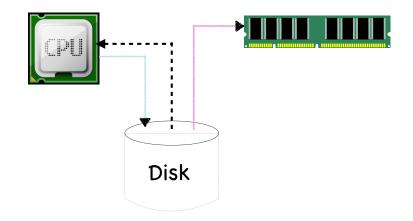
Device puts data on bus &

RAM accepts it

Device interrupts CPU

when done





Communicating with devices

- Explicit I/O instructions (privileged)in and out instructions in x86
- Memory-mapped I/O

 map device registers to memory location

 use memory load and store instructions to read/
 write to registers

How can the OS handle a multitude of devices?

Abstraction!

Encapsulate device specific interactions in a device driver

Implement device neutral interfaces above device drivers

Humans are about 70% water...

...OSs are about 70% device drivers!

File System Stack (simplified)

Application	User
POSIX API [open, read, write, close, etc]	
File System	Kerne
Generic Block Interface [block read/write]	
Block Cache	
Generic Block Layer	
Protocol-specific Block Interface	
Device Driver [SCSI, ATA, etc]	
Memory-mapped I/O, DMA, Interrupts	
Physical Device	