

Crash Recovery

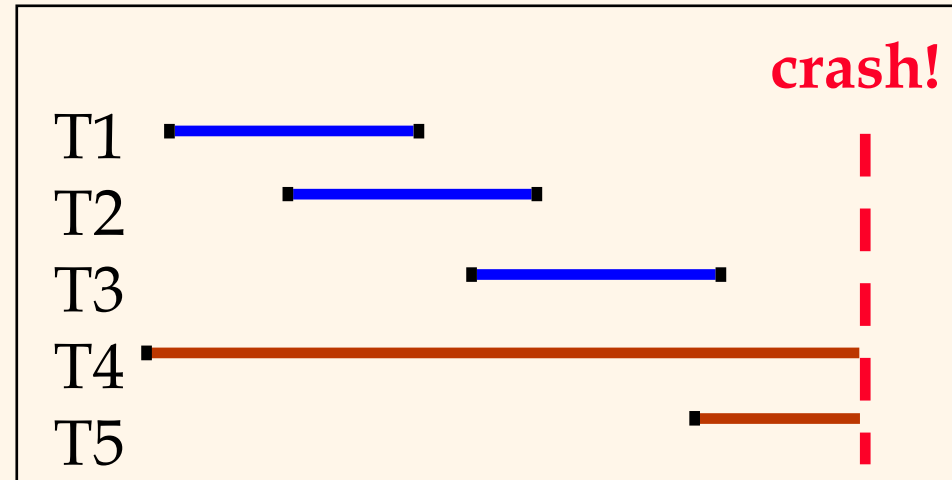
[R&G] Chapter 18

The ACID properties

- ❖ **A**tomicity: All actions in the Xact happen, or none happen.
- ❖ **C**onsistency: If each Xact is consistent, and the DB starts consistent, it ends up consistent.
- ❖ **I**solation: Execution of one Xact is isolated from that of other Xacts.
- ❖ **D**urability: If a Xact commits, its effects persist.
- ❖ The **Recovery Manager** guarantees Atomicity & Durability.

Motivation

- ❖ Atomicity:
 - Transactions may abort (“Rollback”).
- ❖ Durability:
 - What if DBMS stops running? (Causes?)
- ❖ Desired Behavior after system restarts:
 - T1, T2 & T3 should be durable.
 - T4 & T5 should be aborted (effects not seen).



Assumptions

- ❖ Concurrency control is in effect.
 - **Strict 2PL**, in particular.
- ❖ Updates are happening “in place”.
 - i.e. data is overwritten on (deleted from) the disk.
- ❖ A simple scheme to guarantee Atomicity & Durability?

Handling the Buffer Pool

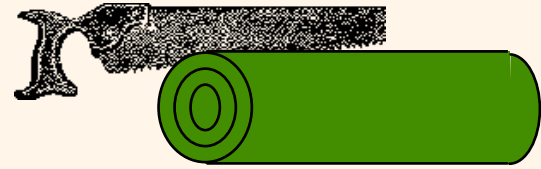
- ❖ **Force** every write to disk?
 - Poor response time.
 - But provides durability.
- ❖ **Steal** buffer-pool frames from uncommitted Xacts?
 - If not, poor throughput.
 - If so, how can we ensure atomicity?

	No Steal	Steal
Force	Trivial	
No Force		Desired

More on Steal and Force

- ❖ STEAL (why enforcing Atomicity is hard)
 - *To steal frame F*: Current page in F (say P) is written to disk; some Xact holds lock on P.
 - What if the Xact with the lock on P aborts?
 - Must remember the old value of P at steal time (to support **UNDO**ing the write to page P).
- ❖ NO FORCE (why enforcing Durability is hard)
 - What if system crashes before a modified page is written to disk?
 - Write as little as possible, in a convenient place, at commit time, to support **REDO**ing modifications.

Basic Idea: Logging



- ❖ Record REDO and UNDO information, for every update, in a *log*.
 - Sequential writes to log (put it on a separate disk).
 - Minimal info (diff) written to log, so multiple updates fit in a single log page.
- ❖ Log: An ordered list of REDO/UNDO actions
 - Log record contains:
 - <XID, pageID, offset, length, old data, new data>
 - and additional control info (which we'll see soon).

Write-Ahead Logging (WAL)

- ❖ The Write-Ahead Logging Protocol:

- ① Must **force** the **log record** for an update before the corresponding **data page** gets to disk.
- ② Must **write all log records** for a Xact before commit.

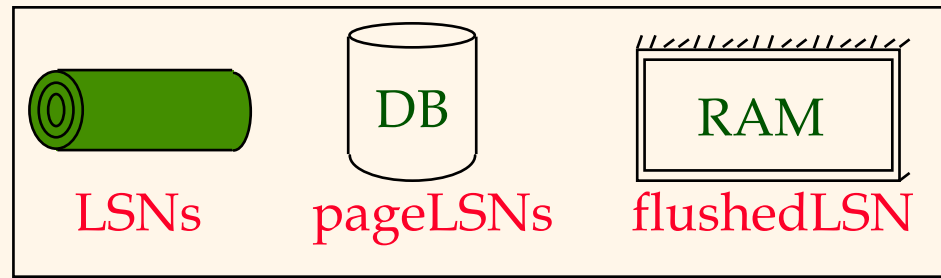
- ❖ #1 guarantees Atomicity.

- ❖ #2 guarantees Durability.

- ❖ Exactly how is logging (and recovery!) done?

- We'll study the ARIES algorithms.

WAL & the Log



- ❖ Each log record has a unique **Log Sequence Number (LSN)**.

- LSNs always increasing.

- ❖ Each data page contains a **pageLSN**.

- The LSN of the most recent *log record* for an update to that page.

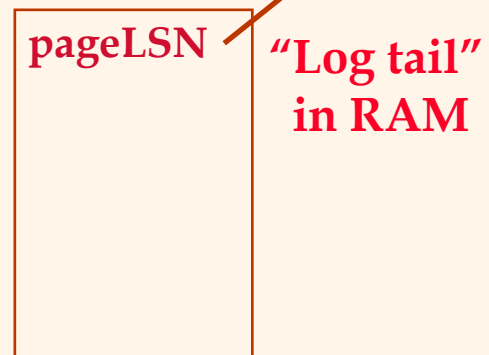
- ❖ System keeps track of **flushedLSN**.

- The max LSN flushed so far.

- ❖ WAL: *Before* a page is written,

- $\text{pageLSN} \leq \text{flushedLSN}$

Log records
flushed to disk



Log Records

LogRecord fields:

**update
records
only** {
 prevLSN
 XID
 type
 pageID
 length
 offset
 before-image
 after-image

Possible log record types:

- ❖ **Update**
- ❖ **Commit**
- ❖ **Abort**
- ❖ **End** (signifies end of commit or abort)
- ❖ **Compensation Log Records (CLRs)**
 - for UNDO actions

Other Log-Related State

❖ Transaction Table:

- One entry per active Xact.
- Contains **XID**, **status** (running/committed/aborted), and **lastLSN**.

❖ Dirty Page Table:

- One entry per dirty page in buffer pool.
- Contains **recLSN** -- the LSN of the log record which first caused the page to be dirty.

Normal Execution of an Xact

- ❖ Series of reads & writes, followed by commit or abort.
 - We will assume that write is atomic on disk.
 - In practice, additional details to deal with non-atomic writes.
- ❖ Strict 2PL.
- ❖ STEAL, NO-FORCE buffer management, with Write-Ahead Logging.

Checkpointing

- ❖ Periodically, the DBMS creates a checkpoint, in order to minimize the time taken to recover in the event of a system crash. Write to log:
 - **begin_checkpoint** record: Indicates when chkpt began.
 - **end_checkpoint** record: Contains current *Xact table* and *dirty page table*. This is a **'fuzzy checkpoint'**:
 - Other Xacts continue to run; so these tables accurate only as of the time of the **begin_checkpoint** record.
 - No attempt to force dirty pages to disk; effectiveness of checkpoint limited by oldest unwritten change to a dirty page. (So it's a good idea to periodically flush dirty pages to disk!)
 - Store LSN of chkpt record in a safe place (**master** record).

The Big Picture: What's Stored Where



LogRecords

prevLSN
XID
type
pageID
length
offset
before-image
after-image



Data pages

each
with a
pageLSN

master record



Xact Table

lastLSN
status

Dirty Page Table

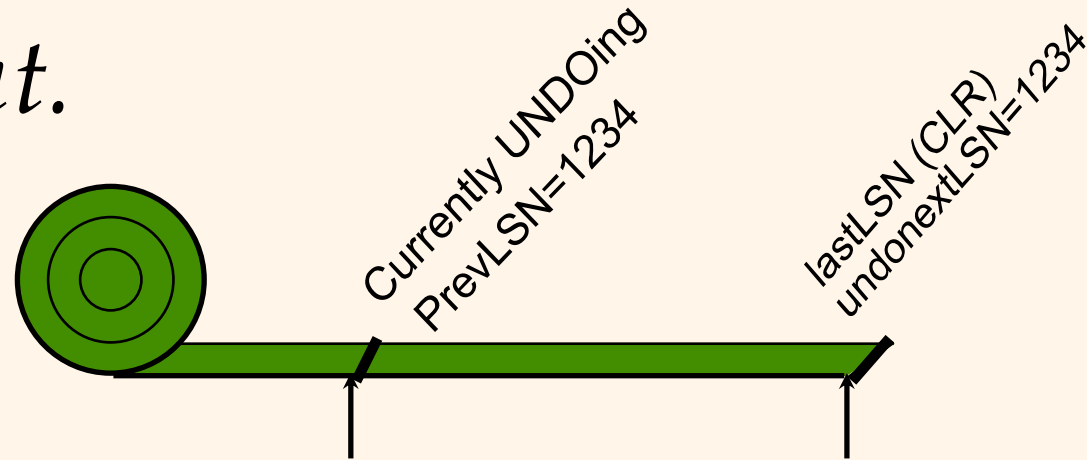
recLSN

flushedLSN

Simple Transaction Abort

- ❖ For now, consider an explicit abort of a Xact.
 - No crash involved.
- ❖ We want to “play back” the log in reverse order, UNDOing updates.
 - Get **lastLSN** of Xact from Xact table.
 - Can follow chain of log records backward via the **prevLSN** field.
 - Before starting UNDO, write an **Abort log record**.
 - For recovering from crash during UNDO!

Abort, cont.

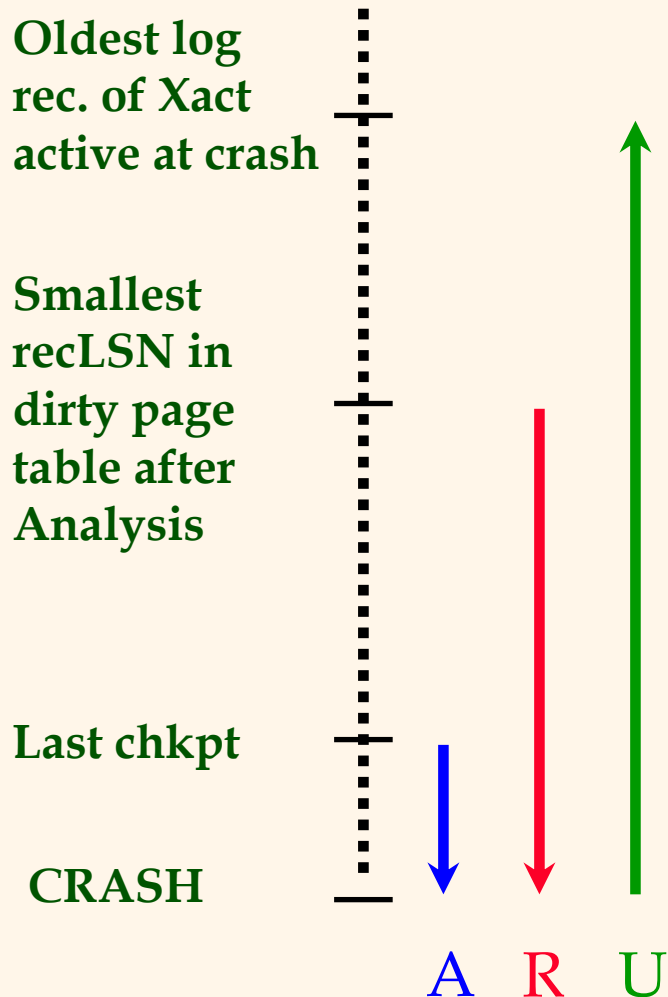


- ❖ To perform UNDO, must have a lock on data!
 - No problem!
- ❖ Before restoring old value of a page, write a CLR:
 - You continue logging while you UNDO!!
 - CLR has one extra field: **undonextLSN**
 - Points to the next LSN to undo (i.e. the prevLSN of the record we're currently undoing).
 - CLR's *never* Undone (but they might be Redone when repeating history: guarantees Atomicity!)
- ❖ At end of UNDO, write an “end” log record.

Transaction Commit

- ❖ Write **commit** record to log.
- ❖ All log records up to Xact's **lastLSN** are flushed.
 - Guarantees that **flushedLSN** \geq **lastLSN**.
 - Note that log flushes are sequential, synchronous writes to disk.
 - Many log records per log page.
- ❖ Commit() returns.
- ❖ Write **end** record to log.

Crash Recovery: Big Picture



- ❖ Start from a **checkpoint** (found via **master** record).
- ❖ Three phases. Need to:
 - Figure out which Xacts committed since checkpoint, which failed (**Analysis**).
 - **REDO** *all* actions.
 - ◆ (repeat history)
 - **UNDO** effects of failed Xacts.

Recovery: The Analysis Phase

- ❖ Reconstruct state at checkpoint.
 - via **end_checkpoint** record.
- ❖ Scan log forward from checkpoint.
 - **End** record: Remove Xact from Xact table.
 - **Other records**: Add Xact to Xact table, set **lastLSN=LSN**, change Xact status on **commit**.
 - **Update** record: If P not in Dirty Page Table,
 - Add P to D.P.T., set its **recLSN=LSN**.

Recovery: The REDO Phase

- ❖ We *repeat History* to reconstruct state at crash:
 - Reapply *all* updates (even of aborted Xacts!), redo CLR's.
- ❖ Scan forward from log rec containing smallest *recLSN* in D.P.T. For each CLR or update log rec *LSN*, REDO the action unless:
 - Affected page is not in the Dirty Page Table, or
 - Affected page is in D.P.T., but has *recLSN* > *LSN*, or
 - *pageLSN* (in DB) \geq *LSN*.
- ❖ To REDO an action:
 - Reapply logged action.
 - Set *pageLSN* to *LSN*. No additional logging!

Recovery: The UNDO Phase

ToUndo = { l | l a lastLSN of a “loser” Xact }

Repeat:

- Choose largest LSN among ToUndo.
- If this LSN is a CLR and undonextLSN == NULL
 - Write an End record for this Xact.
- If this LSN is a CLR, and undonextLSN != NULL
 - Add undonextLSN to ToUndo
- Else this LSN is an update. Undo the update, write a CLR, add prevLSN to ToUndo.

Until ToUndo is empty.

Example of Recovery



Xact Table

lastLSN

status

Dirty Page Table

recLSN

flushedLSN

ToUndo

LSN	LOG
-----	-----

00	begin_checkpoint
----	------------------

05	end_checkpoint
----	----------------

10	update: T1 writes P5
----	----------------------

20	update T2 writes P3
----	---------------------

30	T1 abort
----	----------

40	CLR: Undo T1 LSN 10
----	---------------------

45	T1 End
----	--------

50	update: T3 writes P1
----	----------------------

60	update: T2 writes P5
----	----------------------

✗ CRASH, RESTART

prevLSNs

Example: Crash During Restart!



Xact Table

lastLSN

status

Dirty Page Table

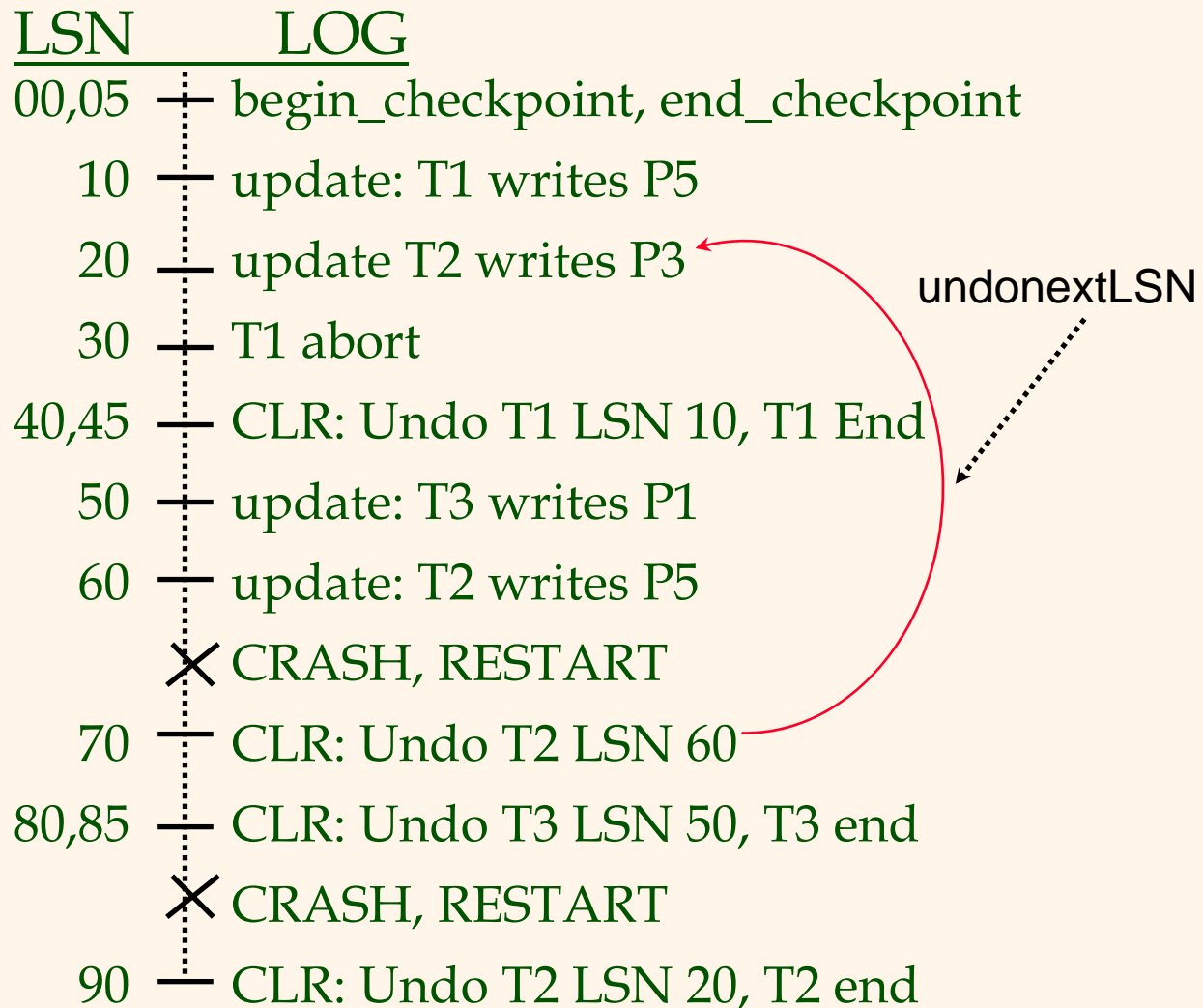
recLSN

flushedLSN

ToUndo

LSN	LOG
00,05	begin_checkpoint, end_checkpoint
10	update: T1 writes P5
20	update T2 writes P3
30	T1 abort
40,45	CLR: Undo T1 LSN 10, T1 End
50	update: T3 writes P1
60	update: T2 writes P5
×	CRASH, RESTART
70	CLR: Undo T2 LSN 60
80,85	CLR: Undo T3 LSN 50, T3 end
×	CRASH, RESTART
90	CLR: Undo T2 LSN 20, T2 end

undonextLSN



The diagram illustrates the state of the log after a crash and restart. A red curved arrow starts from the "undonextLSN" label, points to LSN 20, and then loops back to LSN 70, indicating a redo operation. A dotted arrow points from "undonextLSN" to LSN 40, indicating the next log record to be processed.

Additional Crash Issues

- ❖ What happens if system crashes during Analysis? During REDO?
- ❖ How do you limit the amount of work in REDO?
 - Flush asynchronously in the background.
 - Watch “hot spots”!
- ❖ How do you limit the amount of work in UNDO?
 - Avoid long-running Xacts.

Summary of Logging/Recovery

- ❖ **Recovery Manager** guarantees Atomicity & Durability.
- ❖ Use WAL to allow STEAL/NO-FORCE w/o sacrificing correctness.
- ❖ LSNs identify log records; linked into backwards chains per transaction (via prevLSN).
- ❖ pageLSN allows comparison of data page and log records.

Summary, Cont.

- ❖ **Checkpointing:** A quick way to limit the amount of log to scan on recovery.
- ❖ Recovery works in 3 phases:
 - **Analysis:** Forward from checkpoint.
 - **Redo:** Forward from oldest recLSN.
 - **Undo:** Backward from end to first LSN of oldest Xact alive at crash.
- ❖ Upon Undo, write CLR's.
- ❖ Redo “repeats history”: Simplifies the logic!