# CS519: Computer Networks

Lecture 1: Jan 24, 2004

Intro to Computer Networking

### Lets start at the beginning...



- What is a network for?
  - To allow two or more endpoints to communicate
- What is a network?
  - Nodes connected by links

### • • Lets start at the beginning...

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o Is this a network?



### Lets start at the beginning...

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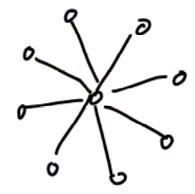
o Is this a network?



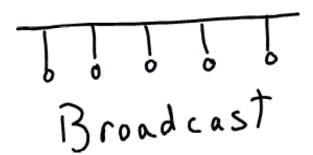
- o Of course it is!
  - Just not very interesting

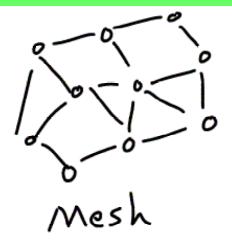
## Other "networks" (network topologies)

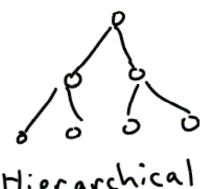
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Star, or Hub and spoke







#### What is a data network?



- The answer is NOT "a network that carries data"
  - Cause you can send "data" (e.g. a fax) over the "voice network"
- "Data network" is often a euphemism for "packet network"
  - And "voice network" is often a euphemism for "circuit network"

- a
- Historically, a circuit network was a network that literally established a physical wired connection between two points
  - With relays, plus amplifiers and stuff
- Before computers, this was the only way to do networks

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- But these days voice is modulated and digitized in numerous ways as it works through the network
  - Very few physical circuits
- So nowadays we consider a circuit network one that appears to establish a fixed "pipe" (amount of bandwidth) between two points

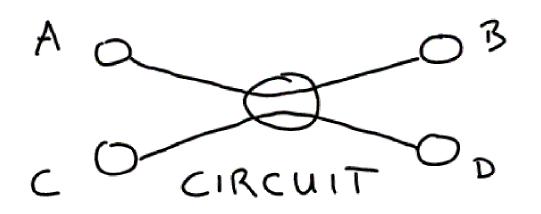
### Types of circuits



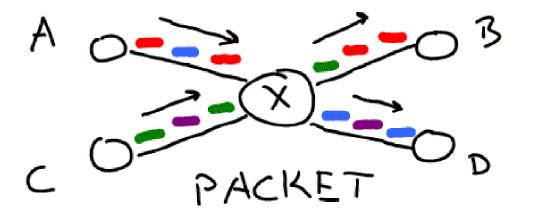
- Synchronous time-division multiplexing (STDM)
  - Each circuit is given a slice of time
- Frequency-division multiplexing (FDM)
  - Each circuit is given a transmission frequency

 By contrast, a packet network allows small units of data (packets) to be individually sent to different destinations

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C can't talk to B while A is talking to B

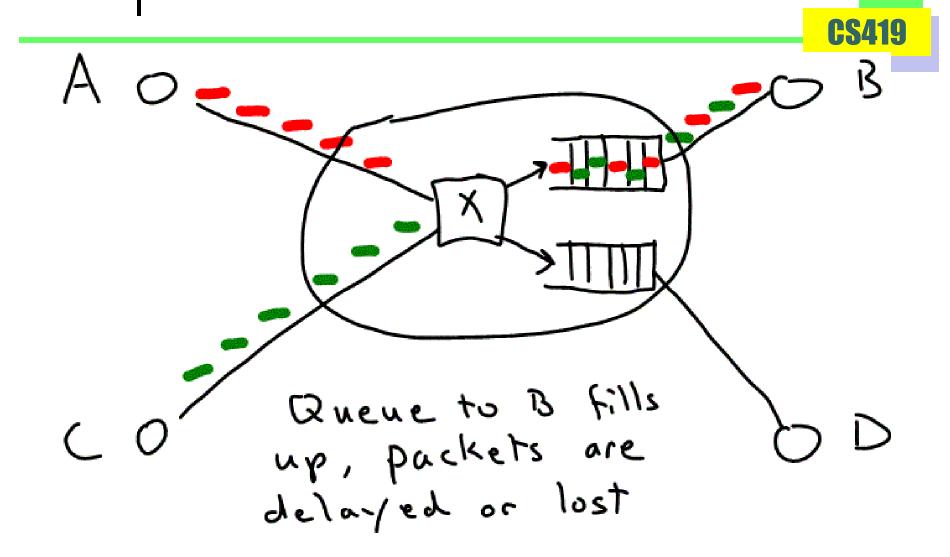


A and C can both talk to B and D

 So clearly packet switched is better than circuit switched, right?

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- So clearly packet switched is better than circuit switched, right?
- Well, as with so much in this world, it depends
- What if A and C try to talk exclusively to B at high speed at the same time?

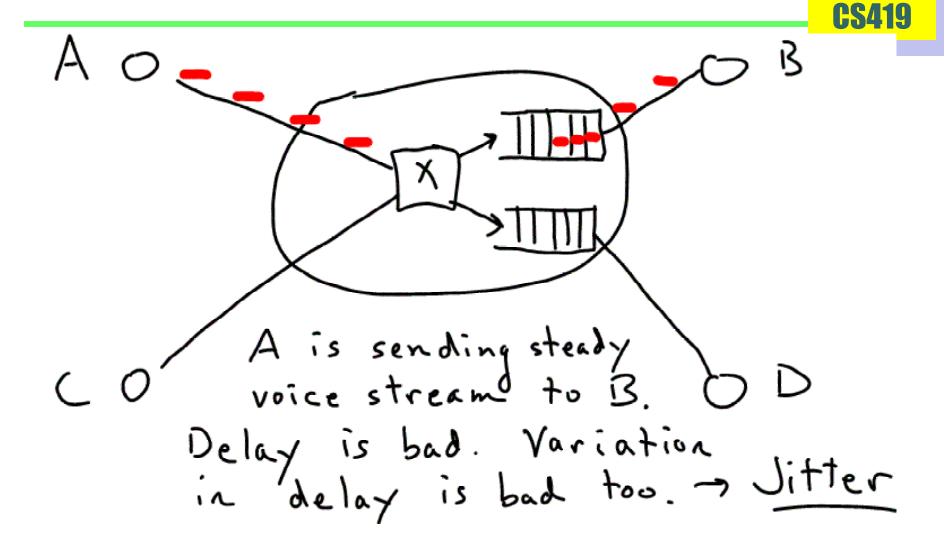
## Delay and packet loss in packet networks

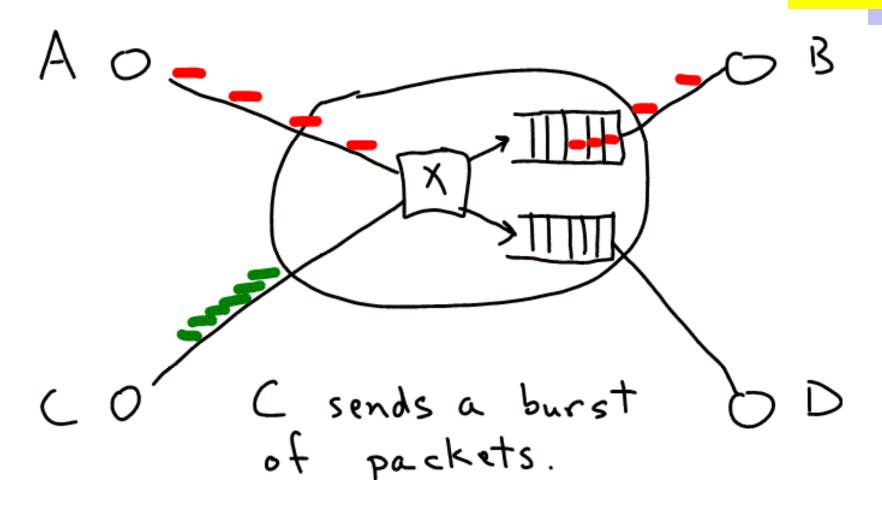


## Delay and packet loss in packet networks

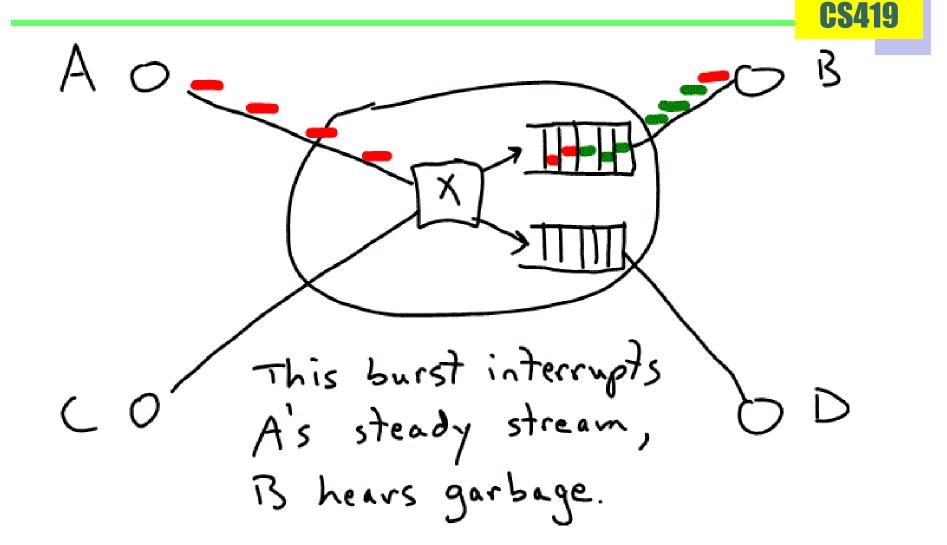
- Can happen any time multiple links feed into a single link
  - And incoming volume exceeds outgoing volume
- Larger queues can reduce packet loss at the expense of more delay
- Ultimately the sources have to slow down (congestion control)
- By contrast, circuit networks can block (busy tone)

### • • Also Jitter





#### Also Jitter



### • Circuits versus packets



- Circuits are an all or nothing proposition
  - Give good quality, if you can get yourself a circuit in the first place
  - Efficient only if the application keeps the circuit full (I.e. a voice stream)
- Packets are more flexible
  - Can send a little or a lot
  - But other traffic can interfere at any time
  - More efficient when traffic is bursty

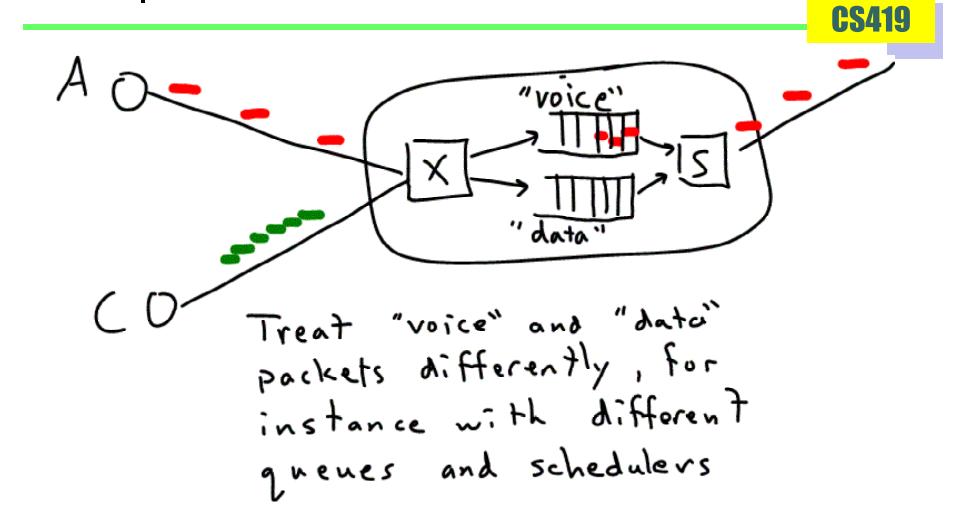
## Can a packet network emulate a circuit?

 After all, our STDM circuit sent data over the wire in "chunks"

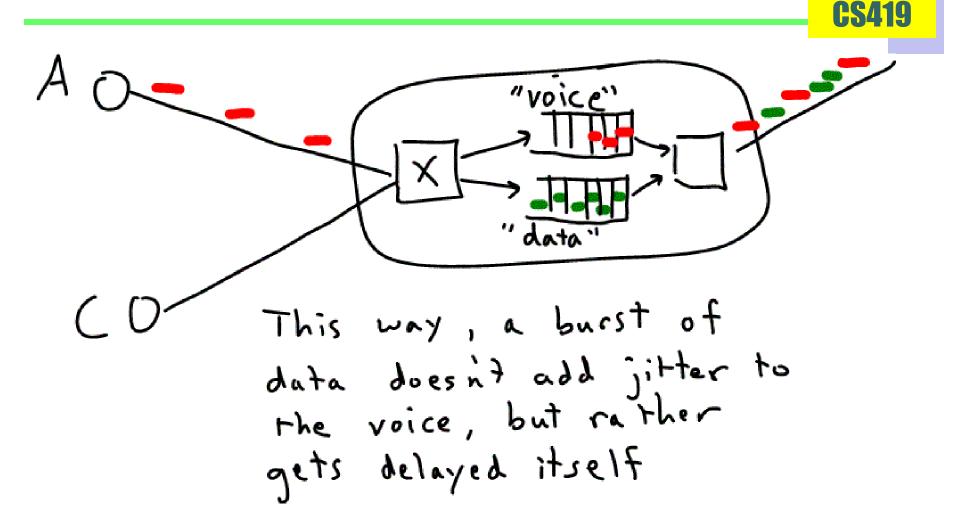
## Can a packet network emulate a circuit?

- After all, our STDM circuit sent data over the wire in "chunks"
- The answer is yes, it can
- And indeed, the first packet networks offered "services" that very much emulated circuits

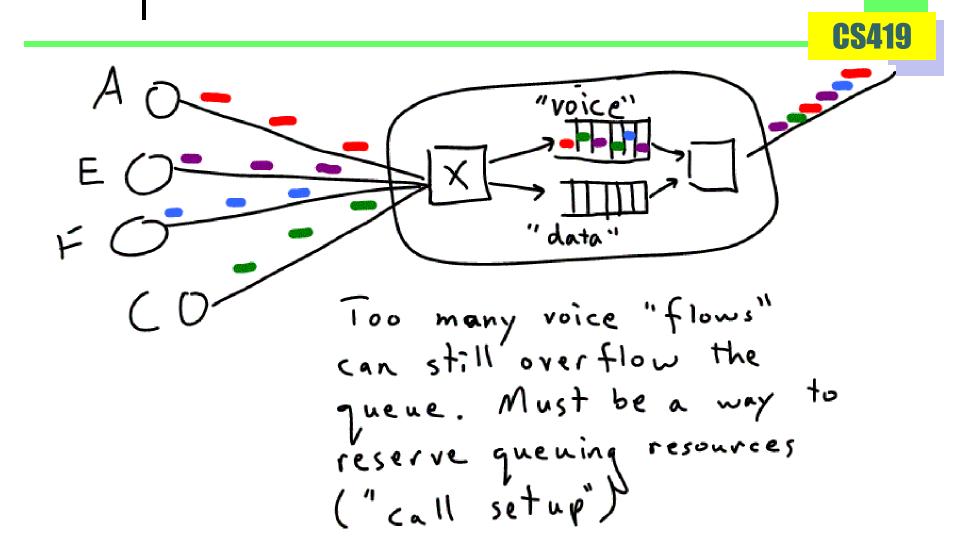
### One way to do it



#### One way to do it



#### But this has complications too



### "Datagram" versus "virtual circuit" networks



- Both are packet networks
  - (We won't discuss pure circuit networks any more in this course)
- Virtual circuit networks have the notion of call setup and blocking
  - But much more complex traffic models than our simple two-queue example
- Datagram networks is how the Internet ultimately got built!

## But virtual circuit networks still important

- We don't see virtual circuit networks to our desktop
  - Though this was the vision for many folks
- But virtual circuit networks formed the unpinning of the Internet
  - Something called ATM
  - Being replaced with MPLS

## This class focuses on the Internet



- Which is a datagram network
- One big topic will be how queues in the Internet manage not to become hopelessly overloaded
  - Many of you know, the answer is TCP, but we'll look at this in detail

### Some terms introduced so far



- Network, node, link, queue
- Circuit and packet networks
  - a.k.a. data and voice networks
- Virtual circuit and datagram networks
- Delay, latency, loss, drop, jitter, blocking

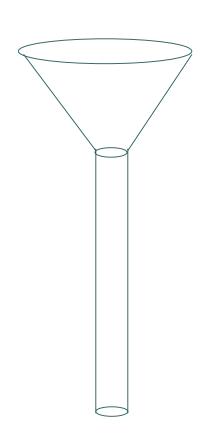
#### Bandwidth and Latency

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- We looked at delay due to queuing
- But there are three main components to delay:
  - Propagation delay
  - Transmit delay
  - Queuing delay

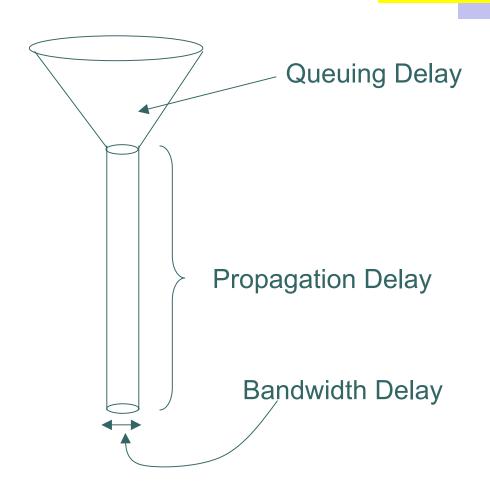
## Queuing, transmit, and propagation delays





## Queuing, transmit, and propagation delays

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### Total latency



- o Total latency =
  - Propagation + Transmit + Queue
- Propagation =
  - Distance / Speed of light
- o Transmit =
  - Packet size / Bandwidth

### Delay x Bandwidth Product

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- Refers to the number of bits you can have "in the pipe" at the same time
  - Or, how many bits you can stuff in the pipe before the first bit comes out the other end
  - Like hot water getting from the water heater to your shower!
- As bandwidth increases (and distance doesn't change) this is becoming an issue

## An extreme (but realistic) Delay x Bandwidth Example

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- Coast-to-coast propagation delay = 15ms
- OC192 link = 10 Gbps
- 10 Gbps x 15ms = 150,000,000 bits =
   19 Mbytes = 7 songs (MP3 files)
- You could stuff 7 songs into an OC192 pipe at Boston before the first song starting arriving in LA!!!

## A more common Delay x Bandwidth Example



- 50ms coast to coast delay (mainly from queuing)
- 100 Mbps Ethernet
- This is about 600Kbytes...still a decent sized file
- Delay x Bandwidth is starting to dominate our thinking about protocol performance

## Common provider bandwidth units



- DSO = 64 Kbps
- o DS1 = 1.544 Mbps
- o DS3 = 44.736 Mbps
- $\circ$  OC3 = 155.52 Mbps
- $\circ$  OC12 = 622.08 Mbps
- $\circ$  OC48 = 2.488 Gbps
- $\circ$  OC192 = 9.953 Gbps
- $\circ$  OC768 = 39.813 Gbps

## Bandwidth and throughput and goodput

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- Bandwidth is the maximum theoretical speed of a pipe
- Throughput is the actual measured speed
  - Vague term because depends on where you measure
- Goodput is the throughput seen by the application
  - Throughput over the pipe can be more than goodput because of dropped and retransmitted packets, control packets, and headers