

## CS412/CS413

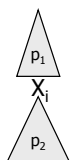
Introduction to Compilers  
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### Lecture 17: Partitioned Attribute Grammars 28 Feb 05

## Plans

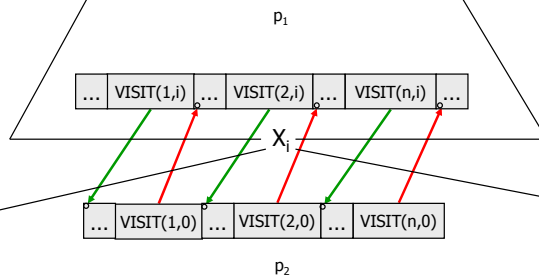
- Each production  $X_0 \rightarrow X_1 \dots X_n$  will have an associated **plan**
- A plan is a linear sequence of **instructions**.
- An instruction is one of
  - **EVAL**  $X_i.a$  evaluate attribute  $a$  of symbol  $X_i$
  - **VISIT**  $(r,i)$  visit neighbor  $i$  for the  $r$ -th time [child 0 = parent]

## Interface



- The attributes of  $X$  constitute the **interface** between plans for  $p_1$  and  $p_2$ .
  - The plan for  $p_1$  evaluates inherited attributes of  $X$
  - The plan for  $p_2$  evaluates synthesized attributes of  $X$

## Coroutine Relationship



- **VISIT** instructions act as coroutine calls

## Consistency of Plans

- The plan for  $p_1$  must be consistent with the plan for every production

$$X \rightarrow \alpha$$

- The plan for  $p_2$  must be consistent with the plans for every production

$$A \rightarrow \alpha X \beta$$

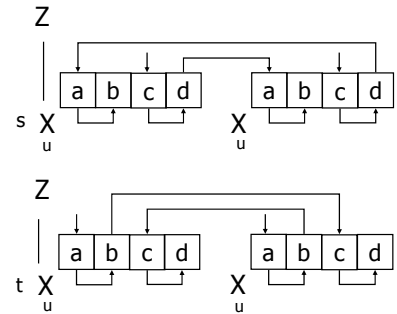
## Plans as Fragments of Topological Orders

- The plans are constructed so that for any derivation tree  $T$ , when the plan instances are "wired up" by **VISITs**, the order of **EVALS** are a topological order for  $D(T)$

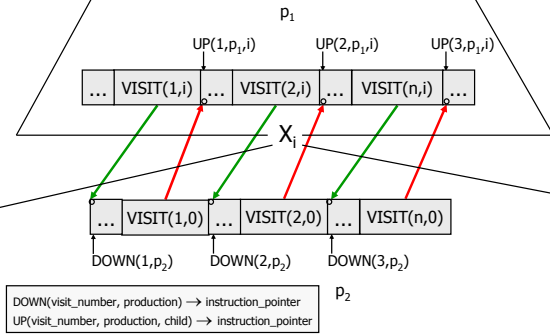
## AG for which no such plans exist

$Z \rightarrow s X_1 X_2$   
 $X_1.a = X_2.d$   
 $X_1.c = 1$   
 $X_2.a = X_1.d$   
 $X_2.c = 2$   
 $Z \rightarrow t X_1 X_2$   
 $X_1.a = 3$   
 $X_1.c = X_2.b$   
 $X_2.a = 4$   
 $X_2.c = X_1.b$   
 $X \rightarrow u$   
 $X.b = X.a$   
 $X.d = X.c$

## AG for which no such plans exist



## UP and DOWN



## Evaluator

```

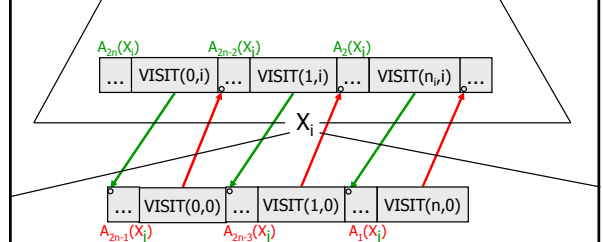
node := root;
ip := MAPDOWN(1, root.rule);
repeat
  case state of
    Xi.a: {
      evaluate Xi.a;
      increment ip;
    }
    VISIT(r,i), i>0: { /* child visit */
      ip := DOWN(r, Xi.rule);
      node := Xi;
    }
    VISIT(r,0): { /* parent visit */
      ip := UP(r, node.parent.rule, node.child_number);
      node = node.parent;
    }
  }
end case
until node = root and instruction at ip = VISIT(1,0);

```

## Partitions

- If such plans are to exist, there must exist, for each nonterminal  $X$ , a partition of  $A(X)$  into classes  $A_{2n}(X)$ ,  $A_{2n-1}(X), \dots, A_2(X), A_1(X)$ , where
    - even  $A_i(X)$  are subsets of  $IA(X)$
    - odd  $A_i(X)$  are subsets of  $SA(X)$
- s.t., for every derivation tree  $T$ , and every nonterminal instance  $X$  in  $T$ , the attribute instances of  $X$  can be evaluated in the order  $A_{2n}(X), A_{2n-1}(X), \dots, A_2(X), A_1(X)$ . Within each  $A_i(X)$ , the order of evaluation is unconstrained and may differ from plan to plan.

## Partitions in Plans



- VISIT instructions act as co-routine calls

## Partitioned Attribute Grammar

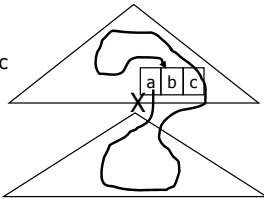
- If, in addition to the existence of the partitions, the grammar is locally acyclic, then it is called **partitioned**.

## Computing Plans from Partitions

- Suppose we know valid partitions for every nonterminal  $X_i$
- To compute the plan for production  $p$   
$$X_0 \rightarrow X_1 \dots X_n$$
  - Start with  $D_p$ , the direct dependency graph of  $p$
  - Collapse all input attribute occurrence of the same class  $A_j(X_i)$  into one node labeled  $VISIT(*,i)$  merging edges
  - Add edges between consecutive partition classes of the given  $X_i$
  - Topological sort the resulting graph and fill in visit numbers in place of \*s

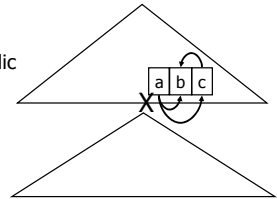
## Ordered Attribute Grammars

- For each nonterminal  $X$ 
  - Construct graph  $DS(X) = \langle A(X), E \rangle$  that over-approximates the transitive dependences that may arise among the attributes of  $X$  in some derivation tree
  - Defer how.
  - If  $DS(X)$  is cyclic for any  $X$  give up.



## OAG: step 1

- For each nonterminal  $X$ 
  - Construct a graph  $DS(X) = \langle A(X), E \rangle$  that over-approximates the transitive dependences that may arise among the attributes of  $X$  in some derivation tree
  - Defer how.
  - If  $DS(X)$  is cyclic for any  $X$  give up.



## OAG: step 2

- Attempt to compute a partition from  $DS(X)$  without reference to the productions in which  $X$  occurs, as follows:
  - Topological sort  $DS(X)$  minimizing alternations between  $IA(X)$  and  $SA(X)$ .
  - Each switch from inherited to synthesized (or vice versa) is a boundary between classes of the partition.

## OAG: step 3

- Use the given method for finding a plan from the partitions. If this fails (because topological sort discovers a cycle), then fail.

## OAG: step 1, cont.

- To compute  $DS(X)$  for all  $X$ 
  - Simultaneously take transitive closures of the direct dependence graphs  $D_p$  for all  $p$ , and whenever an edge between two attributes of the same nonterminal occurrence is created, add it to **every** occurrence of  $X$ .
  - When finished, choose the attributes and edges of an arbitrary occurrence of each nonterminal  $X$  as  $DS(X)$