

CS 4110

Lab #24 - Lambda Calculus

Recall the Church numerals

$$\bar{0} = \lambda s. \lambda z. z$$

$$\bar{1} = \lambda s. \lambda z. s z$$

$$\bar{2} = \lambda s. \lambda z. s (s z)$$

Exercise: Define ISEVEN and DOUBLE.

$$\text{ISEVEN} \triangleq \lambda n. n \underbrace{(\lambda b. b \text{ FALSE } \text{TRUE})}_{\text{NOT}} \text{TRUE}$$

$$\text{DOUBLE} \triangleq \lambda n. \lambda s. \lambda z. n \text{ } s \text{ } (n \text{ } s \text{ } z)$$

Exercise: Define PAIR, FST, SND

$$\text{PAIR} \triangleq \lambda x. \lambda y. \lambda b. b x y$$

$$\text{FST} \triangleq \lambda p. p \text{ TRUE}$$

$$\text{SND} \triangleq \lambda p. p \text{ FALSE}$$

Exercise: Define PRED (with $\text{PRED } \bar{0} = \bar{0}$.)

$$\text{PRED} \triangleq \lambda n. \text{SND } (n \text{ } (\lambda p. \text{PAIR } (\text{SUCC } (\text{FST } p)) (\text{FST } p)) \text{ } (\text{PAIR } \bar{0} \bar{0}))$$

Here is another version that doesn't use pairs (due to Barendregt)

$PRED \triangleq \lambda n. \lambda s. \lambda z. n \quad \underbrace{(\lambda p. \lambda q. q (p s))}_{\text{BUILD}} \quad \underbrace{(\lambda x. z)}_{\text{END } \cancel{\text{START}}} \quad \underbrace{(\lambda x. x)}_{\text{START}}$

↑ continuation ↑ function so far

$PRED \bar{z}$

$= PRED \lambda s. \lambda z. s (s z)$

$=_{\beta} \lambda s. \lambda z. (\lambda p. \lambda q. q (p s)) ((\lambda p. \lambda q. q (p s)) \lambda x. z) (\lambda x. x)$

$=_{\beta} \lambda s. \lambda z. (\lambda p. \lambda q. q (p s)) (\lambda q. q ((\lambda x. z) s)) (\lambda x. x)$

$=_{\beta} \lambda s. \lambda z. (\lambda q. q ((\lambda q. q (\lambda x. z) s) s)) (\lambda x. x)$

$=_{\beta} \lambda s. \lambda z. (\lambda q. q (s z)) (\lambda x. x)$

$=_{\beta} \lambda s. \lambda z. s z$

Exercise: Define EQUAL

$EQUAL = \lambda n. \lambda m. AND \quad (ISZERO (m \ PRED \ n))$
 $\quad \quad \quad (ISZERO (n \ PRED \ m))$

where

$AND = \lambda b. \lambda c. b \ c \ FALSE$

Scott numerals

$$\left. \begin{aligned} \hat{0} &= \lambda s. \lambda z. z \\ \hat{1} &= \lambda s. \lambda z. s \hat{0} \\ \hat{2} &= \lambda s. \lambda z. s \hat{1} \end{aligned} \right\}$$

Intuition: case analysis!
- successor (of $n-1$)
- zero

Define: SUCC, PRED, ISZERO

$$\text{SUCC} = \lambda n. \lambda s. \lambda z. s \ n$$

$$\text{PRED} = \lambda n. n \ (\lambda p. p) \ \hat{0}$$

$$\text{ISZERO} = \lambda n. n \ (\lambda p. \text{FALSE}) \ \text{TRUE}$$

Define: ADD

$$\text{ADD} = \lambda n. \lambda m. n \ \text{SUCC} \ (m \ p)$$

Define: C2S and S2C

$$\text{C2S} = \lambda n. n \ \text{SUCC} \ \hat{0}$$

$$\text{S2C} = \lambda n. n \ \text{SUCC} \ \hat{0}$$