CS 4110 – Programming Languages and Logics Lecture #14: More λ -calculus



1 Lambda calculus evaluation

There are many different evaluation strategies for the λ -calculus. The most permissive is *full* β *reduction*, which allows any *redex*—i.e., any expression of the form $(\lambda x. e_1) e_2$ —to step to $e_1\{e_2/x\}$ at any time. It is defined formally by the following small-step operational semantics rules:

$$\frac{e_1 \to e_1'}{e_1 \ e_2 \to e_1' \ e_2} \qquad \frac{e_2 \to e_2'}{e_1 \ e_2 \to e_1 \ e_2'} \qquad \frac{e_1 \to e_1'}{\lambda x. \ e_1 \to \lambda x. \ e_1'} \qquad \beta \frac{(\lambda x. \ e_1) \ e_2 \to e_1 \{e_2/x\}}{(\lambda x. \ e_1) \ e_2 \to e_1 \{e_2/x\}}$$

The *call by value* (CBV) strategy enforces a more restrictive strategy: it only allows an application to reduce after its argument has been reduced to a value (i.e., a λ -abstraction) and does not allow evaluation under a λ . It is described by the following small-step operational semantics rules (here we show a left-to-right version of CBV):

$$\frac{e_1 \to e_1'}{e_1 e_2 \to e_1' e_2} \qquad \frac{e_2 \to e_2'}{v_1 e_2 \to v_1 e_2'} \qquad \beta \frac{}{(\lambda x. e_1) v_2 \to e_1 \{v_2/x\}}$$

Finally, the *call by name* (CBN) strategy allows an application to reduce even when its argument is not a value but does not allow evaluation under a λ . It is described by the following small-step operational semantics rules:

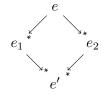
$$\frac{e_1 \to e_1'}{e_1 e_2 \to e_1' e_2} \qquad \beta \frac{}{(\lambda x. e_1) e_2 \to e_1 \{e_2/x\}}$$

2 Confluence

It is not hard to see that the full β reduction strategy is non-deterministic. This raises an interesting question: does the choices made during the evaluation of an expression affect the final result? The answer turns out to be no: full β reduction is *confluent* in the following sense:

Theorem (Confluence). If $e \to^* e_1$ and $e \to^* e_2$ then there exists e' such that $e_1 \to^* e'$ and $e_2 \to^* e'$.

Confluence can be depicted graphically as follows:



Confluence is often also called the Church-Rosser property.

3 Substitution

Each of the evaluation relations for λ -calculus has a β defined in terms of a substitution operation on expressions. Because the expressions involved in the substitution may share some variable names (and because we are working up to α -equivalence) the definition of this operation is slightly subtle and defining it precisely turns out to be tricker than might first appear.

As a first attempt, consider an obvious (but incorrect) definition of the substitution operator. Here we are substituting e for x in some other expression:

$$y\{e/x\} = \begin{cases} e & \text{if } y = x \\ y & \text{otherwise} \end{cases}$$

$$(e_1 e_2)\{e/x\} = (e_1\{e/x\}) (e_2\{e/x\})$$

$$(\lambda y.e_1)\{e/x\} = \lambda y.e_1\{e/x\} \quad \text{where } y \neq x$$

The intuitive idea is that the last rule relies on α -equivalence to "rewrite" abstractions that use x so they do not conflict. Unfortunately, this definition produces the wrong results when we substitute an expression with free variables under a λ . For example,

$$(\lambda y.x)\{y/x\} = (\lambda y.y)$$

To fix this problem, we need to revise our definition so that when we substitute under a λ we do not accidentally bind variables in the expression we are substituting. The following definition correctly implements *capture-avoiding substitution*:

$$\begin{array}{lcl} y\{e/x\} &=& \left\{ \begin{array}{ll} e & \text{if } y=x \\ y & \text{otherwise} \end{array} \right. \\ (e_1\ e_2)\{e/x\} &=& \left(e_1\{e/x\}\right)\left(e_2\{e/x\}\right) \\ (\lambda y.e_1)\{e/x\} &=& \lambda y.(e_1\{e/x\}) & \text{where } y\neq x \text{ and } y\not\in fv(e) \end{array}$$

Note that in the case for λ -abstractions, we require that the bound variable y be different from the variable x we are substituting for and that y not appear in the free variables of e, the expression we are substituting. Because we work up to α -equivalence, we can always pick y to satisfy these side conditions. For example, to calculate $(\lambda z.x\ z)\{(w\ y\ z)/x\}$ we first rewrite $\lambda z.x\ z$ to $\lambda u.x\ u$ and then apply the substitution, obtaining $\lambda u.(w\ y\ z)\ u$ as the result.