

CS 3410 Lab 3

Fall 2026



Agenda

1 C Memory Management

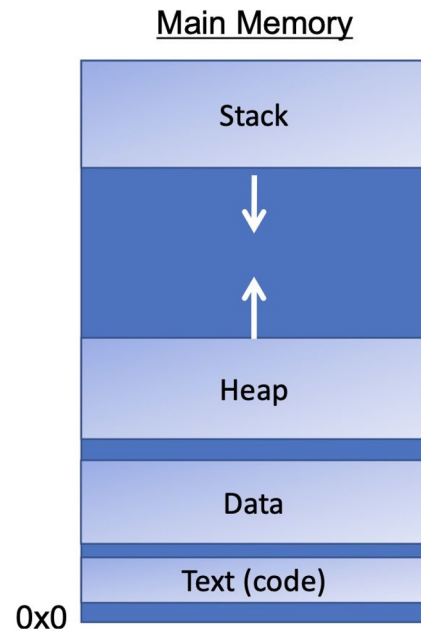
2 Priority Queue

3 A3 Tips

C Memory Management

The Heap

- The heap is an area of memory below the stack that grows up towards higher addresses
- Unlike the stack, where memory goes away when a function finishes, the heap provides memory that persists until the caller is done with it



How do we access memory on the heap?

- `malloc()`: Request a pointer to a contiguous block of memory on the heap
- `free()`: Release or deallocate the allocated memory back to the operating system



malloc() and free() syntax

```
int main() {  
    int *ptr = (int*)malloc(sizeof(*ptr) * 10);  
    // ... do some things  
    free(ptr);  
}
```

Rule of Thumb

Every call to `malloc()` should have a corresponding call to `free()`

- An allocation that is not freed by the time the program ends is called a **Memory Leak**

Dual Meaning of * Symbol: Type Declaration

```
int c = 10;  
int d = 5;  
int* a = (int *)malloc(sizeof(int));  
*a = c;  
int** b = &a;  
**b = d;  
free(a);
```

Type decl:
type of “a” is a pointer to an int

Dual Meaning of * Symbol: Dereference Operator

```
int c = 10;  
int d = 5;  
int* a = (int *)malloc(sizeof(int));  
*a = c;  
int** b = &a;  
**b = d;  
free(a);
```

*** is a dereference operator**

It says:
update the **memory** location
pointed to by “a” with the **value**
stored in “c”

Exercise: Draw out the references between memory blocks.

Note: Assume all the addresses are of 2 bytes and the integer defaults to 4 bytes.

```
int c = 10;
int d = 5;
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*a = c;
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```



Stack		
Address	Name	Value
0xFFFF	c	10
0xFFFE		
0xFFFD		
0xFFFC		
0xFFFB	d	5
0xFFFA		
0xFFF9		
0xFFF8		
0xFFF7		
0xFFF6		
0xFFF5		

0x0008		
0x0007		
0x0006		
0x0005		
0x0004		
0x0003		
0x0002		
0x0001		
0x0000		
Address	Name	Value
Heap		

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Stack		
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0xFFFFF	c	10
0xFFFFE		
0xFFFFD		
0xFFFFC		
0xFFFFB	d	5
0xFFFFA		
0xFFFF9		
0xFFFF8		
0xFFFF7		
0xFFFF6		
0xFFFF5		

0x0008		
0x0007		
0x0006		
0x0005		
0x0004		
0x0003		
0x0002		
0x0001		
0x0000		
Address	Name	Value
Heap		

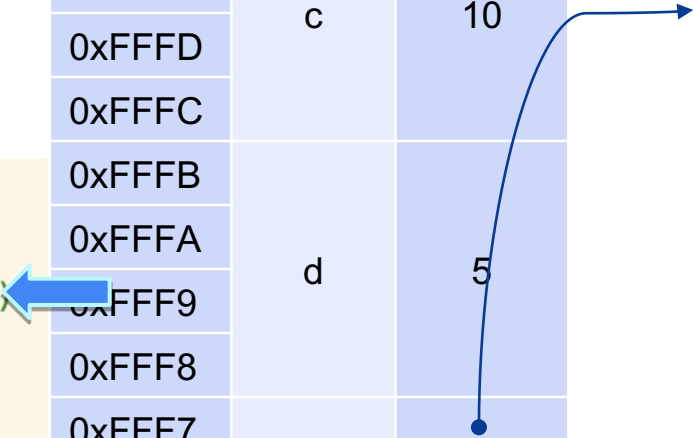
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Address	Name	Value
0xFFFF	c	10
0xFFFE		
0xFFFD		
0xFFFC		
0xFFFB	d	5
0xFFFA		
0xFF99		
0xFF98		
0xFF97	a	0x0005
0xFF96		
0xFF95		
0xFF94		

0x0008	c copy	10
0x0007		
0x0006		
0x0005		
0x0004		
0x0003		
0x0002		
0x0001		
0x0000		
Address	Name	Value
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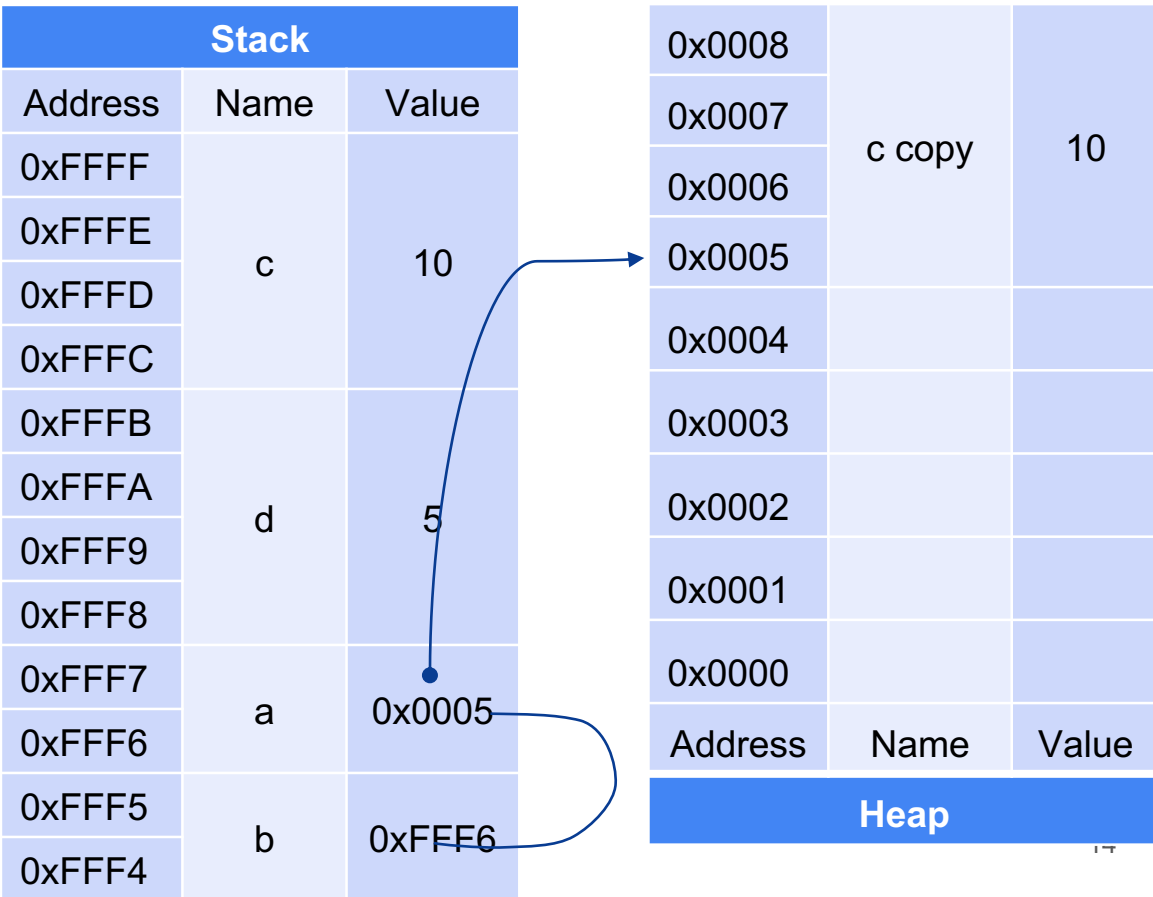
Stack		
Address	Name	Value
0xFFFF	c	10
0xFFFE		
0xFFFD		
0xFFFC		
0xFFFB	d	5
0xFFFA		
0xFFF9		
0xFFF8		
0xFFF7	a	0x0005
0xFFF6		
0xFFF5		
0xFFF4		

0x0008	c copy	10
0x0007		
0x0006		
0x0005		
0x0004		
0x0003		
0x0002		
0x0001		
0x0000		
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0xFFF6		
0xFFF5	b	0xFFF6
0xFFF4		

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0x0006		
0x0005		
0x0004		
0x0003		
0x0002		
0x0001		
0x0000		
Address	Name	Value
Heap		

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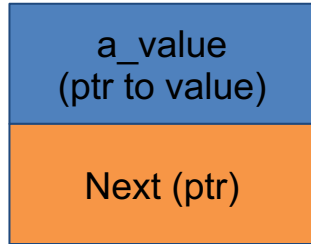
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int d = 5;
int* a = (int *)malloc(sizeof(int));
*a = c;
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**b = d;
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Stack		
Address	Name	Value
0xFFFF	c	10
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0xFFFB	d	5
0xFFFA		
0xFFF9		
0xFFF8		
0xFFF7	a	
0xFFF6		
0xFFF5	b	0xFFF6
0xFFF4		

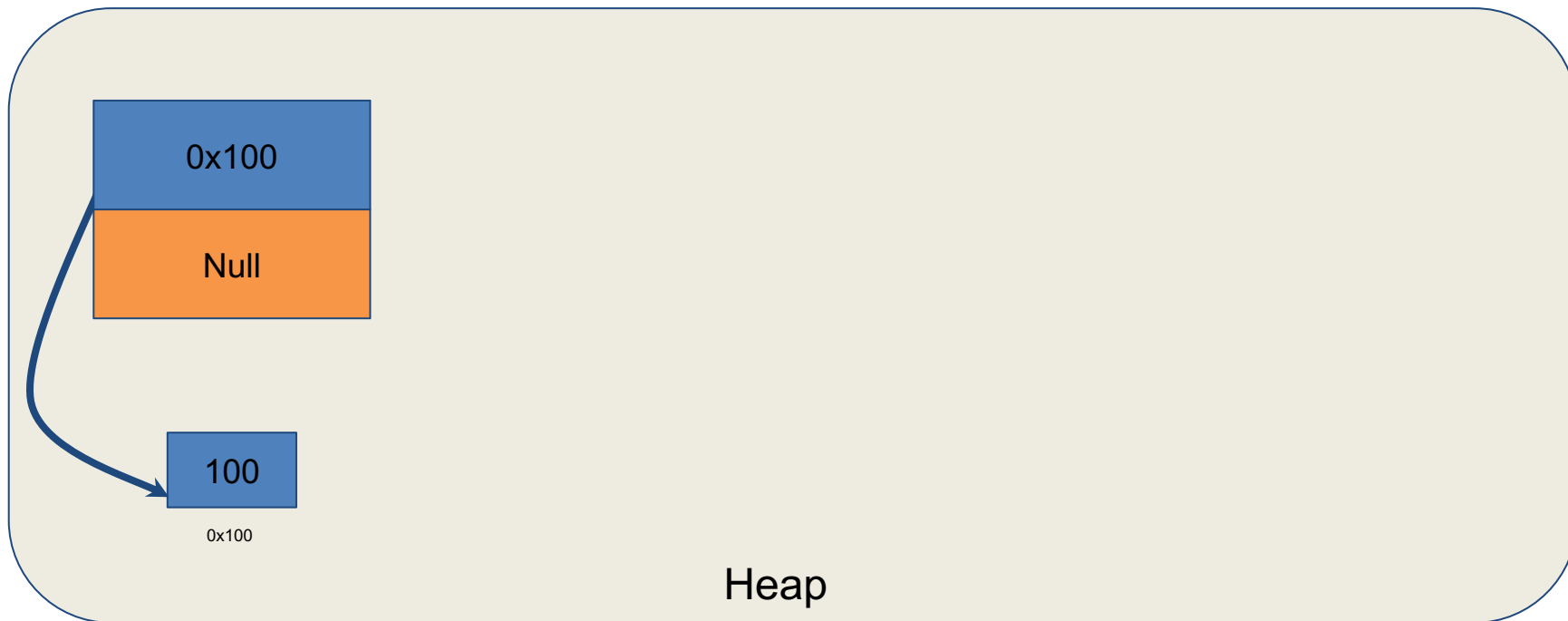
0x0008	e-copy	5
0x0007		
0x0006		
0x0005		
0x0004		
0x0003		
0x0002		
0x0001		
0x0000		
Address	Name	Value
Heap		

Lab: Priority Queue

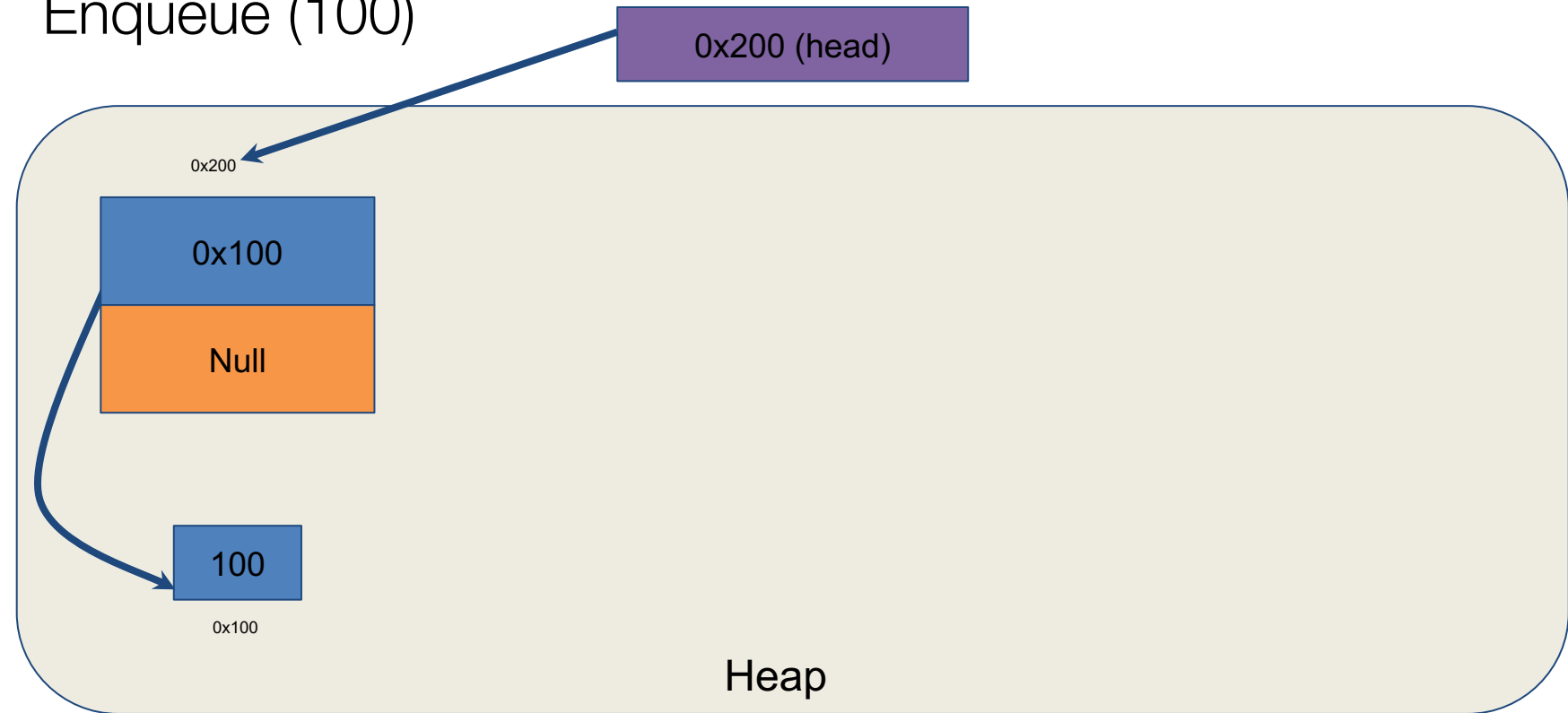
Malloc PQNode



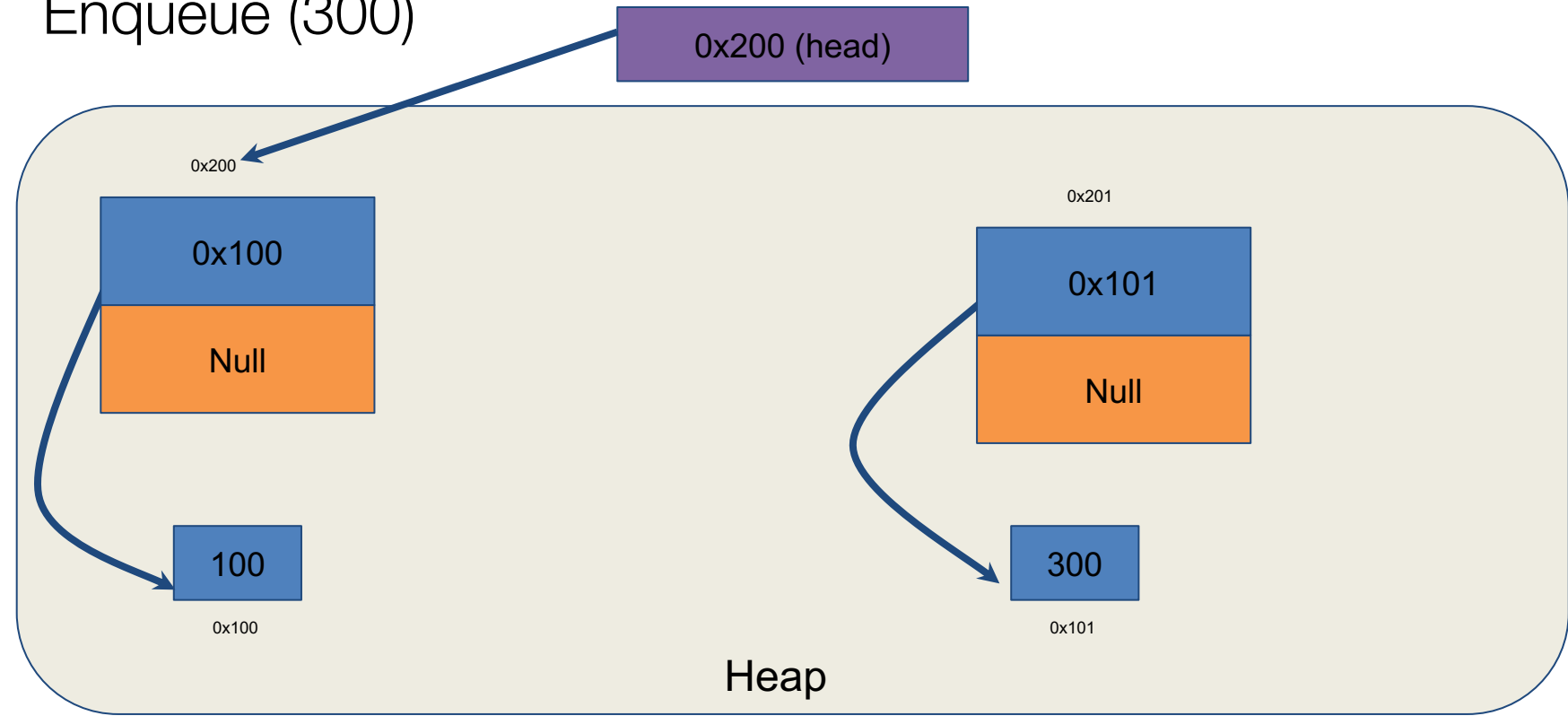
Malloc PQNode



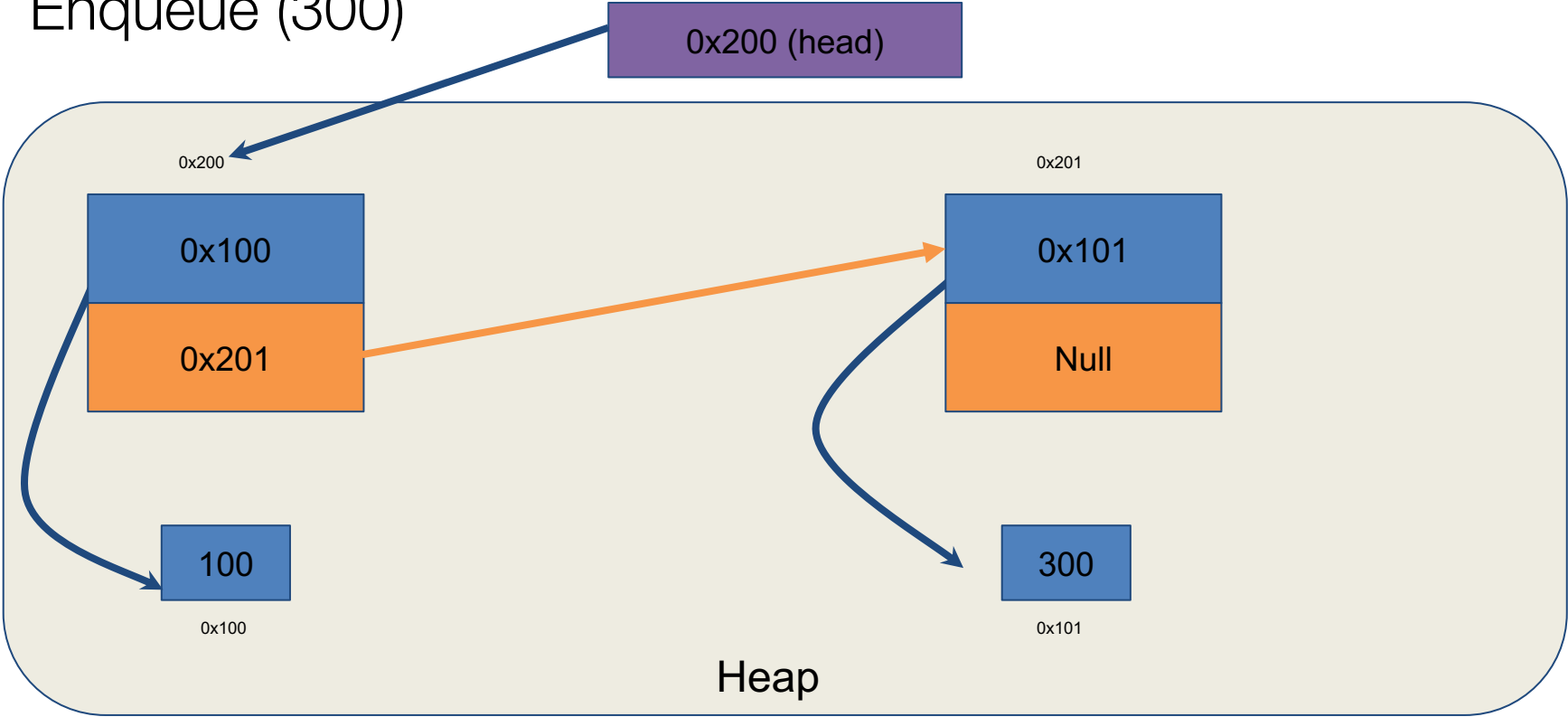
Enqueue (100)



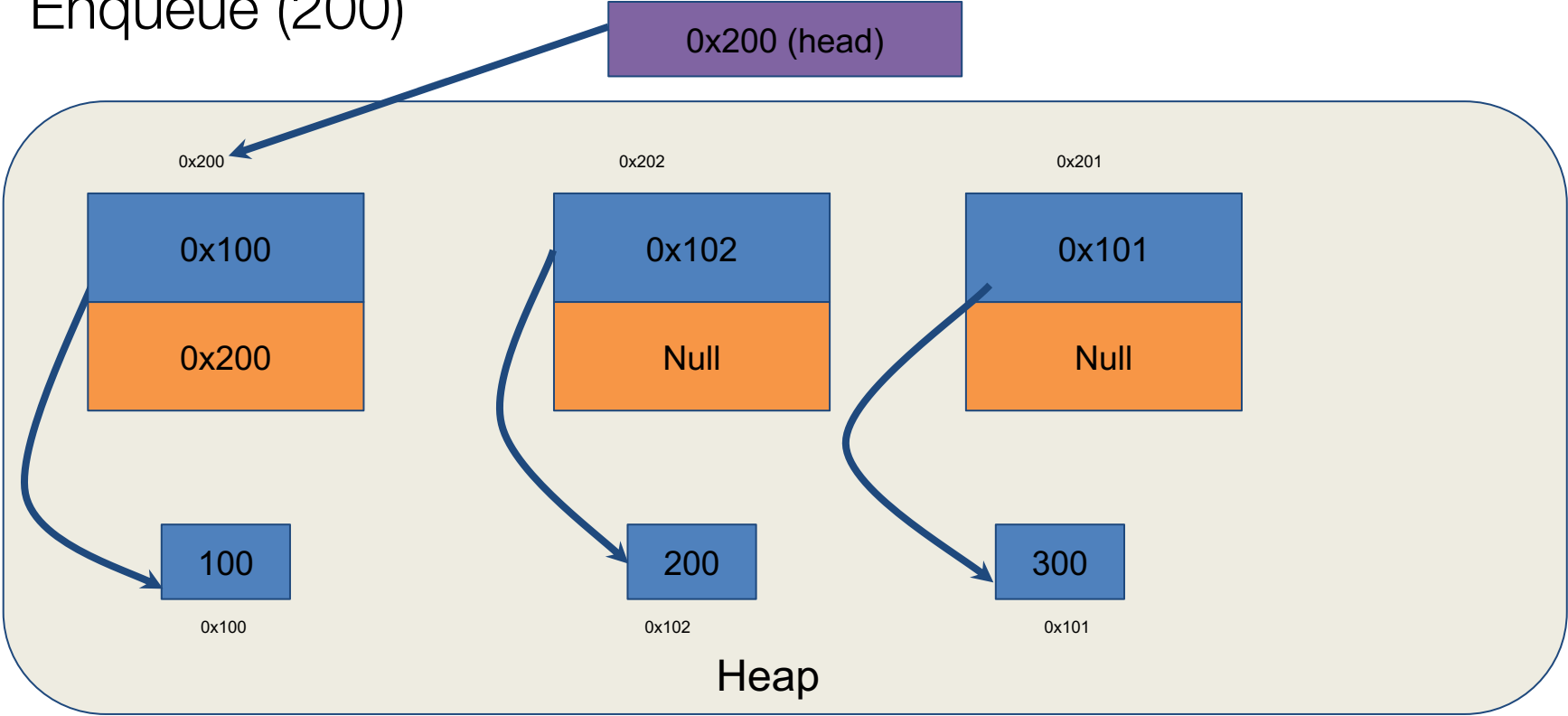
Enqueue (300)



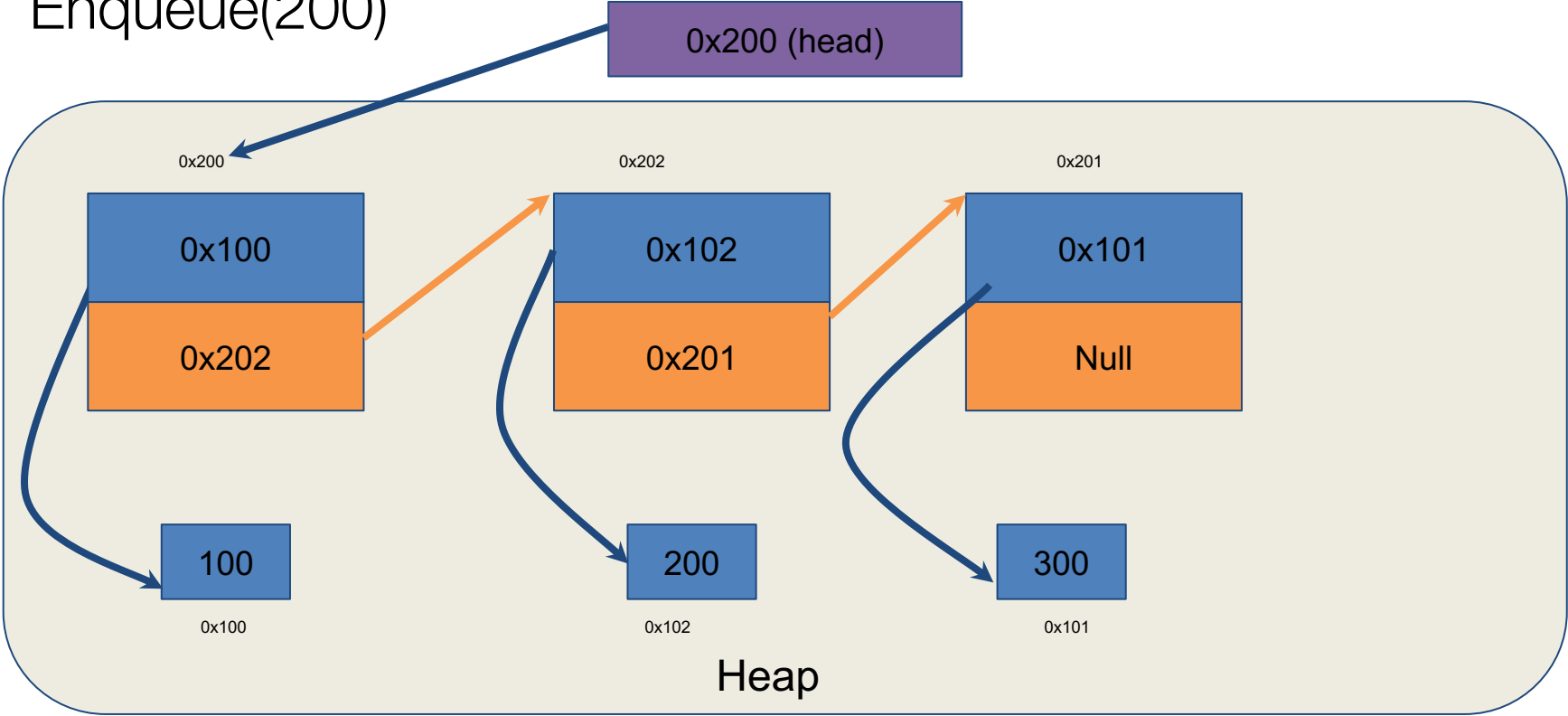
Enqueue (300)



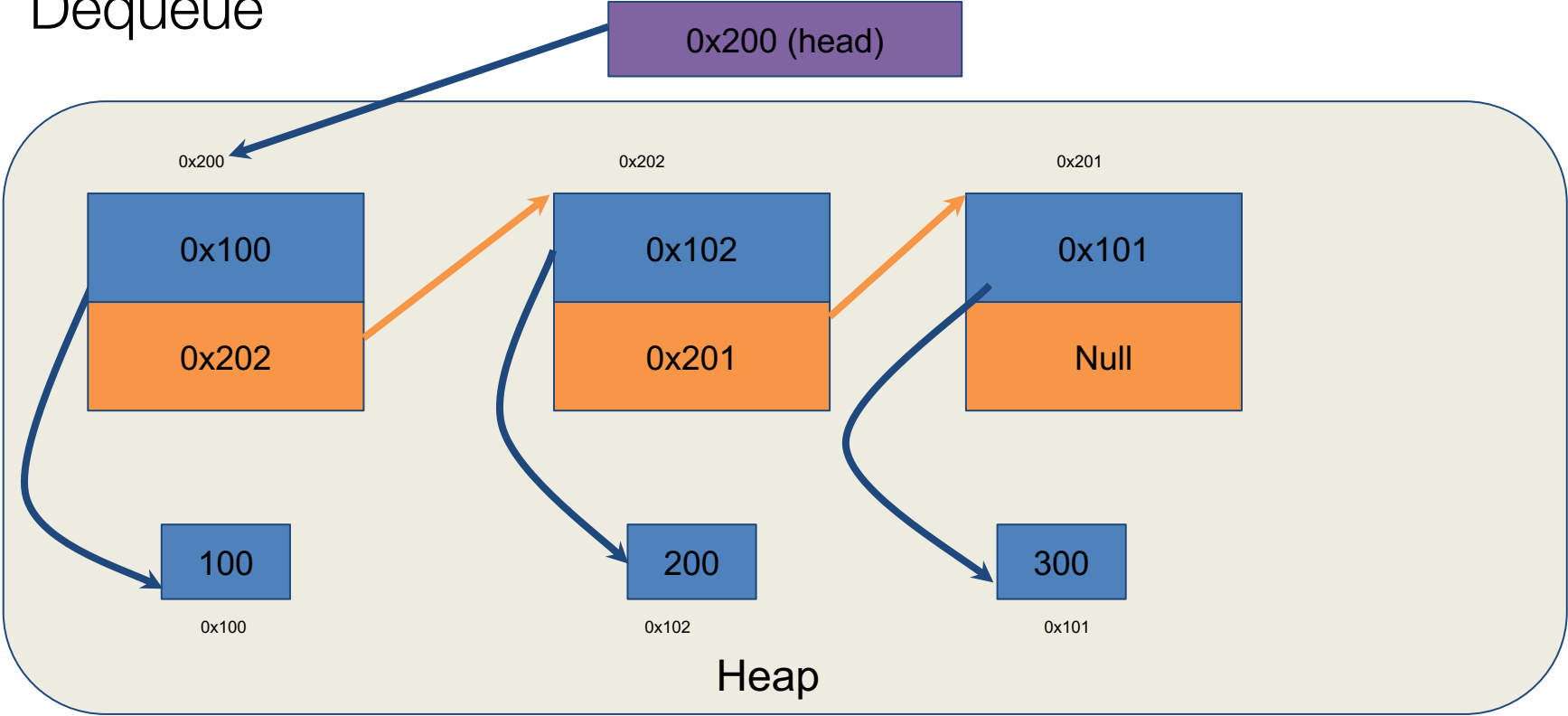
Enqueue (200)



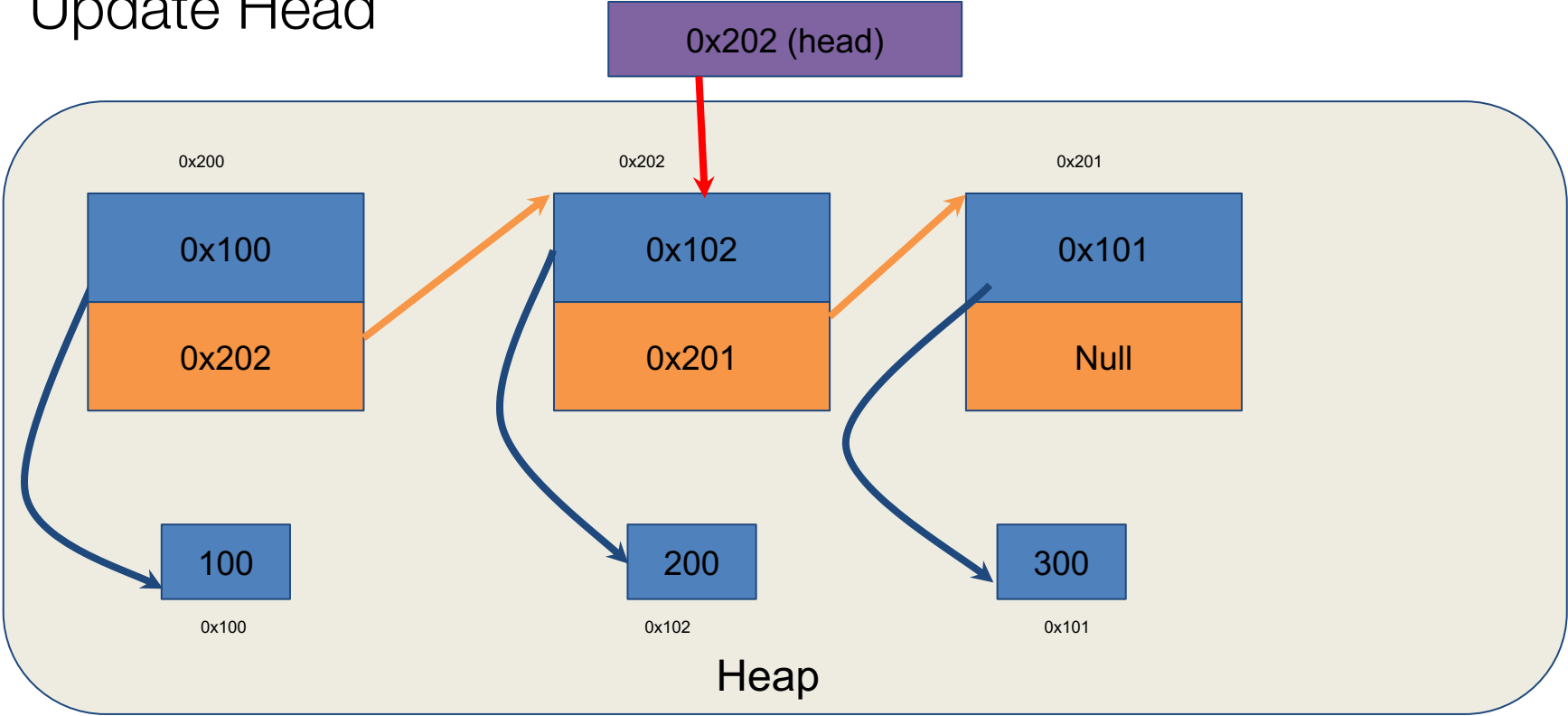
Enqueue(200)



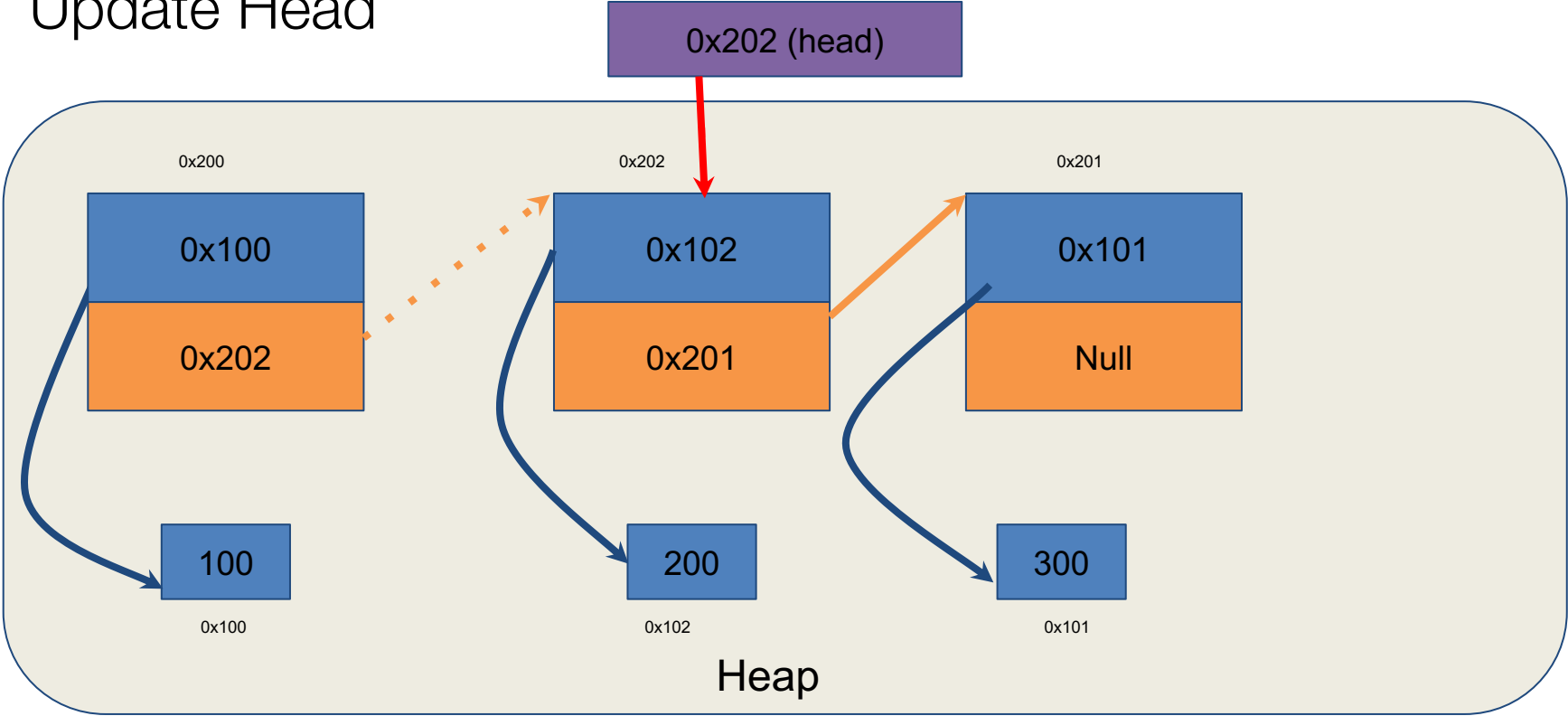
Dequeue



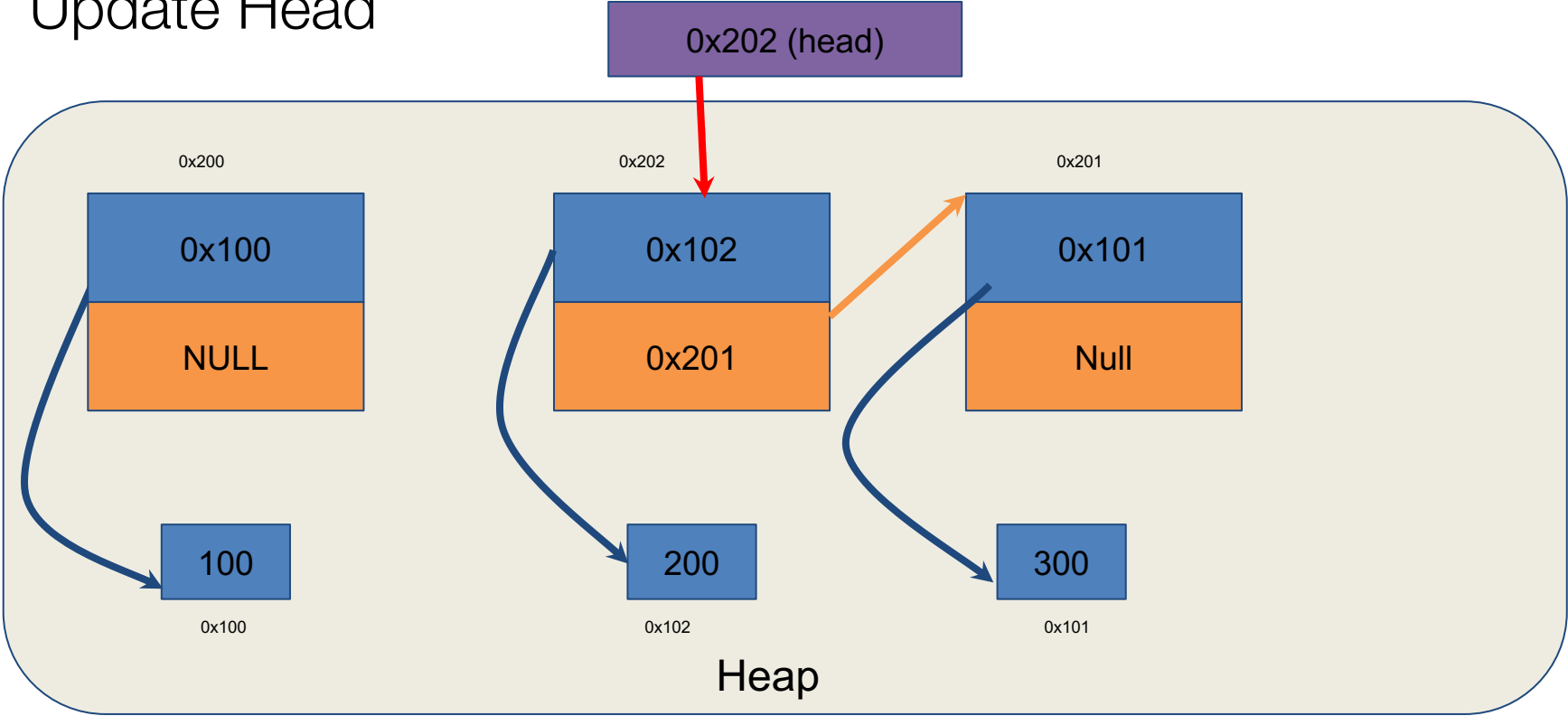
Update Head



Update Head



Update Head



Lab Hints

- You can implement `stack_push` by calling `pq_enqueue` and passing in `NULL` as the compare function. If you do this, you'll need to think carefully about your implementation for `pq_enqueue`.
- You'll also need to write a compare function to order the nodes in your priority queue. Nodes are sorted in *ascending* order, first by frequency, then by ASCII value. Follow this convention when implementing the function:
 - `cmp_fn(a, b) < 0` $\rightarrow$ a is ordered before b
 - `cmp_fn(a, b) >= 0` $\rightarrow$ a is ordered after b



A3 Tips

Start early, start early, start early!

- This assignment has many intricacies. To give yourself enough time to test the functionality end-to-end, you'll want to start early.
- **Test thoroughly!** We've provided a unit-testing framework for you to use with test files already started for you. Add more tests as you need them to make sure your code is correct.
 - This applies to `huffman.c` as the complete test suite for Priority Queue is already given



Test your code!

- Testing before you get to Task 2 will be crucial. We've provided you with a simple unit-testing library called `cu_unit`.
- The structure of a unit test is as follows:

```
int _test_equality() {  
    cu_start(); // We must call this at the start of every test  
    //-----  
    int x = 4;  
    int a_x = &x;  
    cu_check(x == *a_x); // We can call cu_check() as many times as we want  
    within a test  
    // -----  
    cu_end(); // We must call this at the end of every test  
}
```

Test your code!

- We can run the test by adding it to `main()`:

```
int _test_equality() {  
    cu_start(); // We must call this at the start of every test  
    //-----  
    int x = 4;  
    int a_x = &x;  
    cu_check(x == *a_x); // We can call cu_check() as many times as we want  
    within a test  
    // -----  
    cu_end(); // We must call this at the end of every test  
}  
  
int main() {  
    cu_run(_test_equality); // Run the test  
    return 0;  
}
```



Test your code!

- Tests will not be graded for this assignment, but it's a good idea to test your code before moving on to the next task, as we'll be grading your code for each section individually.
- Test your priority queue with different types, comparison functions, etc.
- Test your Huffman tree on different files (C programs, for instance), and with a variety of characters.



Remember!

- **In A3 (not this lab), you must** use our wrapper functions `my_malloc()` and `my_free()` for any allocations or deallocations.
- They follow the exact same syntax as `malloc()` and `free()` and perform the same kind of allocation/deallocation.
- The only thing extra is that they log the operation, which will help you and us detect and pinpoint memory leaks in your implementation.



Good luck!