

# CIS 330: Applied Database Systems

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Lecture 10: Stored Procedures, Database Security

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## Two Topics

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- Stored Procedures
- Database Security

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## Stored Procedures

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- What is a stored procedure:
  - Program executed through a single SQL statement
  - Executed in the process space of the server
- Advantages:
  - Can encapsulate application logic while staying "close" to the data
  - Reuse of application logic by different users
  - Avoid tuple-at-a-time return of records through cursors

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## Stored Procedures: Examples

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```
CREATE PROCEDURE ShowNumReservations
  SELECT S.sid, S.sname, COUNT(*)
  FROM Sailors S, Reserves R
  WHERE S.sid = R.sid
  GROUP BY S.sid, S.sname
```

Stored procedures can have [parameters](#):

- Three different modes: IN, OUT, INOUT

```
CREATE PROCEDURE IncreaseRating(
  IN sailor_sid INTEGER, IN increase INTEGER)
UPDATE Sailors
  SET rating = rating + increase
  WHERE sid = sailor_sid
```

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## Stored Procedures: Examples (Contd.)

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Stored procedure do not have to be written in SQL:

```
CREATE PROCEDURE TopSailors(
  IN num INTEGER)
LANGUAGE JAVA
EXTERNAL NAME
  "file:///c:/storedProcs/rank.jar"
```

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## Calling Stored Procedures

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```
EXEC SQL BEGIN DECLARE SECTION
Int sid;
Int rating;
EXEC SQL END DECLARE SECTION
```

```
// now increase the rating of this sailor
EXEC CALL IncreaseRating(:sid,:rating);
```

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## Calling Stored Procedures (Contd.)

### JDBC:

```
CallableStatement cstmt=  
con.prepareCall("{call  
  ShowSailors});  
ResultSet rs =  
cstmt.executeQuery();  
while (rs.next()) {  
  ...  
}
```

### SQL:

```
#sql iterator  
  ShowSailors(...);  
ShowSailors showsailors;  
#sql showsailors={CALL  
  ShowSailors};  
while (showsailors.next()) {  
  ...  
}
```

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## SQL/PSM

Most DBMSs allow users to write stored procedures in a simple, general-purpose language (close to SQL) → SQL/PSM standard is a representative

### **Declare a stored procedure:**

```
CREATE PROCEDURE name(p1, p2, ..., pn)  
  local variable declarations  
  procedure code;
```

### **Declare a function:**

```
CREATE FUNCTION name (p1, ..., pn) RETURNS  
  sqlDataType  
  local variable declarations  
  function code;
```

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## Main SQL/PSM Constructs

```
CREATE FUNCTION rate Sailor  
  (IN sailorId INTEGER)  
  RETURNS INTEGER  
DECLARE rating INTEGER  
DECLARE numRes INTEGER  
SET numRes = (SELECT COUNT(*)  
  FROM Reserves R  
  WHERE R.sid = sailorId)  
IF (numRes > 10) THEN rating =1;  
ELSE rating = 0;  
END IF;  
RETURN rating;
```

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### Main SQL/PSM Constructs (Contd.)

- Local variables (DECLARE)
- RETURN values for FUNCTION
- Assign variables with SET
- Branches and loops:
  - IF (condition) THEN statements;  
ELSEIF (condition) statements;  
... ELSE statements; END IF;
  - LOOP statements; END LOOP
- Queries can be parts of expressions
- Can use cursors naturally without "EXEC SQL"

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### Two Topics

- Stored Procedures
- Database Security

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### Introduction to DB Security

- **Secrecy:** Users should not be able to see things they are not supposed to.
  - E.g., A student can't see other students' grades.
- **Integrity:** Users should not be able to modify things they are not supposed to.
  - E.g., Only instructors can assign grades.
- **Availability:** Users should be able to see and modify things they are allowed to.

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## Access Controls

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- A **security policy** specifies who is authorized to do what.
- A **security mechanism** allows us to enforce a chosen security policy.
- Two main mechanisms at the DBMS level:
  - Discretionary access control
  - Mandatory access control

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## Discretionary Access Control

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- Based on the concept of access rights or **privileges** for objects (tables and views), and mechanisms for giving users privileges (and revoking privileges).
- Creator of a table or a view automatically gets all privileges on it.
  - DBMS keeps track of who subsequently gains and loses privileges, and ensures that only requests from users who have the necessary privileges (at the time the request is issued) are allowed.

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## GRANT Command

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`GRANT privileges ON object TO users [WITH GRANT OPTION]`

- ❖ The following **privileges** can be specified:
  - ❖ **SELECT**: Can read all columns (including those added later via ALTER TABLE command).
  - ❖ **INSERT(col-name)**: Can insert tuples with non-null or non-default values in this column.
    - ❖ INSERT means same right with respect to all columns.
  - ❖ **DELETE**: Can delete tuples.
  - ❖ **REFERENCES (col-name)**: Can define foreign keys (in other tables) that refer to this column.
- ❖ If a user has a privilege with the **GRANT OPTION**, can pass privilege on to other users (with or without passing on the **GRANT OPTION**).
- ❖ Only owner can execute CREATE, ALTER, and DROP.

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## GRANT and REVOKE of Privileges

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- GRANT INSERT, SELECT ON Sailors TO Horatio
  - Horatio can query Sailors or insert tuples into it.
- GRANT DELETE ON Sailors TO Yuppy WITH GRANT OPTION
  - Yuppy can delete tuples, and also authorize others to do so.
- GRANT UPDATE (*rating*) ON Sailors TO Dustin
  - Dustin can update (only) the *rating* field of Sailors tuples.
- GRANT SELECT ON ActiveSailors TO Guppy, Yuppy
  - This does NOT allow the 'uppies to query Sailors directly!
- REVOKE: When a privilege is revoked from X, it is also revoked from all users who got it *solely* from X.

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## GRANT/REVOKE on Views

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- If the creator of a view loses the SELECT privilege on an underlying table, the view is dropped!  
(Example: Why is having access to a SELECT privilege on a view important?)
- If the creator of a view loses a privilege held with the grant option on an underlying table, (s)he loses the privilege on the view as well; so do users who were granted that privilege on the view!

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## Example

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Joe: GRANT SELECT ON Sailors TO Art WITH GRANT OPTION

Art: GRANT SELECT ON Sailors TO Bob WITH GRANT OPTION

Joe: REVOKE SELECT ON Sailors FROM Art CASCADE

CASCADE versus RESTRICT:

- CASCADE: All abandoned privileges are also revoked
- RESTRICT: Command is rejected if abandoned privileges would result

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## Another Example

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Joe: GRANT SELECT ON Sailors TO Art WITH GRANT OPTION  
Joe: GRANT SELECT ON Sailors TO Bob WITH GRANT OPTION  
Art: GRANT SELECT ON Sailors TO Bob WITH GRANT OPTION  
Joe: REVOKE SELECT ON Sailors FROM Art CASCADE

Bob retains his privilege.

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## Another Example

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Joe: GRANT SELECT ON Sailors TO Art WITH GRANT OPTION  
Joe: GRANT SELECT ON Sailors TO Art WITH GRANT OPTION  
Joe: REVOKE SELECT ON Sailors FROM Art CASCADE

Art's privilege is revoked.

Note: We can just revoke the GRANT OPTION:  
• REVOKE GRANT OPTION ON Sailors FROM Art CASCADE

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## Authorization Graphs

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Example:  
Joe: GRANT SELECT ON Sailors TO Art WITH GRANT OPTION  
Art: GRANT SELECT ON Sailors TO Bob WITH GRANT OPTION  
Bob: GRANT SELECT ON Sailors TO Art WITH GRANT OPTION  
Joe: GRANT SELECT ON Sailors TO Cal WITH GRANT OPTION  
Cal: GRANT SELECT ON Sailors TO Bob WITH GRANT OPTION  
Joe: REVOKE SELECT ON Sailors FROM Art CASCADE

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## Views and Security

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- Views can be used to present necessary information (or a summary), while hiding details in underlying relation(s).
  - Given ActiveSailors, but not Sailors or Reserves, we can find sailors who have a reservation, but not the *bid's* of boats that have been reserved.
- Creator of view has a privilege on the view if (s)he has the privilege on all underlying tables.
- Together with GRANT/REVOKE commands, views are a very powerful access control tool.

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## Role-Based Authorization

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- In SQL-92, privileges are actually assigned to **authorization ids**, which can denote a single user or a group of users.
- In SQL:1999 (and in many current systems), privileges are assigned to **roles**.
  - Roles can then be granted to users and to other roles.
  - Reflects how real organizations work.
  - Illustrates how standards often catch up with "de facto" standards embodied in popular systems.

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## Security to the Level of a Field!

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- Can create a view that only returns one field of one tuple. (How?)
- Then grant access to that view accordingly.
- Allows for *arbitrary* granularity of control
  - A bit clumsy to specify.
  - Can be hidden under a good UI.

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## Mandatory Access Control

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- Based on system-wide policies that cannot be changed by individual users.
  - Each DB object is assigned a security class.
  - Each subject (user or user program) is assigned a clearance for a security class.
  - Rules based on security classes and clearances govern who can read/write which objects.
- Most commercial systems do not support mandatory access control. Versions of some DBMSs do support it; used for specialized (e.g., military) applications.

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## Why Mandatory Control?

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- Discretionary control has some flaws, e.g., the *Trojan horse* problem:
  - Dick creates Horsie and gives INSERT privileges to Justin (who doesn't know about this).
  - Dick modifies the code of an application program used by Justin to additionally write some secret data to table Horsie.
  - Now, Justin can see the secret info.
- The modification of the code is beyond the DBMSs control, but it can try and prevent the use of the database as a channel for secret information.

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## Bell-LaPadula Model

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- Objects (e.g., tables, views, tuples)
- Subjects (e.g., users, user programs)
- Security classes:
  - Top secret (TS), secret (S), confidential (C), unclassified (U):  $TS > S > C > U$
- Each object and subject is assigned a class.
  - Subject S can read object O only if  $class(S) \geq class(O)$  (Simple Security Property)
  - Subject S can write object O only if  $class(S) \leq class(O)$  (\*-Property)

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## Intuition

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- Idea is to ensure that information can never flow from a higher to a lower security level.
- E.g., If Dick has security class C, Justin has class S, and the secret table has class S:
  - Dick's table, Horsie, has Dick's clearance, C.
  - Justin's application has his clearance, S.
  - So, the program cannot write into table Horsie.
- The mandatory access control rules are applied in addition to any discretionary controls that are in effect.

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## Multilevel Relations

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bid	bname	color	class
101	Salsa	Red	S
102	Pinto	Brown	C

- Users with S and TS clearance will see both rows; a user with C will only see the 2<sup>nd</sup> row; a user with U will see no rows.
- If user with C tries to insert <101,Pasta,Blue,C>:
  - Allowing insertion violates key constraint
  - Disallowing insertion tells user that there is another object with key 101 that has a class > C!
  - Problem resolved by treating class field as part of key.

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## Statistical DB Security

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- Statistical DB: Contains information about individuals, but allows only aggregate queries (e.g., average age, rather than Joe's age).
- New problem: It may be possible to **infer** some secret information!
  - E.g., If I know Joe is the oldest sailor, I can ask "How many sailors are older than X?" for different values of X until I get the answer 1; this allows me to infer Joe's age.
- Idea: Insist that each query must involve at least N rows, for some N. Will this work? (No!)

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## Why Minimum N is Not Enough

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- By asking "How many sailors older than X?" until the system rejects the query, can identify a set of N sailors, including Joe, that are older than X; let X=55 at this point.
- Next, ask "What is the sum of ages of sailors older than X?" Let result be S1.
- Next, ask "What is sum of ages of sailors other than Joe who are older than X, plus my age?" Let result be S2.
- S1-S2 is Joe's age!

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## Summary

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- Three main security objectives: secrecy, integrity, availability.
- DB admin is responsible for overall security.
  - Designs security policy, maintains an **audit trail**, or history of users' accesses to DB.
- Two main approaches to DBMS security: discretionary and mandatory access control.
  - Discretionary control based on notion of privileges.
  - Mandatory control based on notion of security classes.
- Statistical DBs try to protect individual data by supporting only aggregate queries, but often, individual information can be inferred.

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