

7.1 Naïve Gaussian Elimination

8.1 The LU Factorization

- Motivating $Ax=b$: Newton's method for systems of nonlinear equations (pp. 96-99)
- C&K 7.1: Naive Gaussian Elimination
 - See answered problems (1, 2, 3ab, 4, 5, 6, 7abc)
- C&K 8.1: Matrix Factorizations (see pp. 293-302 on LU factorization)
 - See answered problems (1-8)

$$Ax = b$$

Motivation: Newton's Method

(for systems of nonlinear equations, pp.96-99)

In the general case, a system of N nonlinear equations in N unknowns x_i can be displayed in the form

$$\begin{cases} f_1(x_1, x_2, \dots, x_N) = 0 \\ f_2(x_1, x_2, \dots, x_N) = 0 \\ \vdots \\ f_N(x_1, x_2, \dots, x_N) = 0 \end{cases}$$

Using vector notation, we can write this system in a more elegant form:

$$\mathbf{F}(\mathbf{X}) = \mathbf{0}$$

by defining column vectors as

$$\mathbf{F} = [f_1, f_2, \dots, f_N]^T$$

$$\mathbf{X} = [x_1, x_2, \dots, x_N]^T$$

Motivation: Newton's Method

(for systems of nonlinear equations, pp.96-99)

In general, Newton's method uses this iteration:

$$\mathbf{X}^{(k+1)} = \mathbf{X}^{(k)} - [\mathbf{F}'(\mathbf{X}^{(k)})]^{-1} \mathbf{F}(\mathbf{X}^{(k)})$$

In practice, the computational form of Newton's method does not involve inverting the Jacobian matrix but rather solves the **Jacobian linear systems**

$$[\mathbf{F}'(\mathbf{X}^{(k)})] \mathbf{H}^{(k)} = -\mathbf{F}(\mathbf{X}^{(k)}) \quad (6)$$

The next iteration of Newton's method is then

$$\mathbf{X}^{(k+1)} = \mathbf{X}^{(k)} + \mathbf{H}^{(k)} \quad (7)$$

This is **Newton's method for nonlinear systems**. The linear system (6) can be solved by procedures *Gauss* and *Solve* as discussed in Chapter 7. Small systems of order 2 can be solved easily. (See Problem 3.2.39.)

Motivation: Newton's Method

(for systems of nonlinear equations, pp.96-99)

Example 2 (p. 99)

$$\begin{cases} x + y + z = 3 \\ x^2 + y^2 + z^2 = 5 \\ e^x + xy - xz = 1 \end{cases}$$

$$\mathbf{X} = [0.1, \quad 1.2, \quad 2.5]^T$$

for $k = 1$ **to** 10 **do**

$$\mathbf{F} = \begin{bmatrix} x_1 + x_2 + x_3 - 3 \\ x_1^2 + x_2^2 + x_3^2 - 5 \\ e^{x_1} + x_1x_2 - x_1x_3 - 1 \end{bmatrix}$$

$$\mathbf{J} = \begin{bmatrix} 1 & 1 & 1 \\ 2x_1 & 2x_2 & 2x_3 \\ e^{x_1} + x_2 - x_3 & x_1 & -x_1 \end{bmatrix}$$

solve $\mathbf{JH} = \mathbf{F}$

$$\mathbf{X} = \mathbf{X} - \mathbf{H}$$

end for

A simple electrical network contains a number of resistances and a single source of electromotive force (a battery) as shown in Figure 7.1. Using Kirchhoff's laws and Ohm's law, we can write a system of linear equations that govern this circuit. If x_1 , x_2 , x_3 , and x_4 are the loop currents as shown, then the equations are

$$\begin{cases} 15x_1 - 2x_2 - 6x_3 &= 300 \\ -2x_1 + 12x_2 - 4x_3 - x_4 &= 0 \\ -6x_1 - 4x_2 + 19x_3 - 9x_4 &= 0 \\ -x_2 - 9x_3 + 21x_4 &= 0 \end{cases}$$

Systems of equations like this, even those that contain hundreds of unknowns, can be solved by using the methods developed in this chapter. The solution to the preceding system is

$$x_1 = 26.5 \quad x_2 = 9.35 \quad x_3 = 13.3 \quad x_4 = 6.13$$

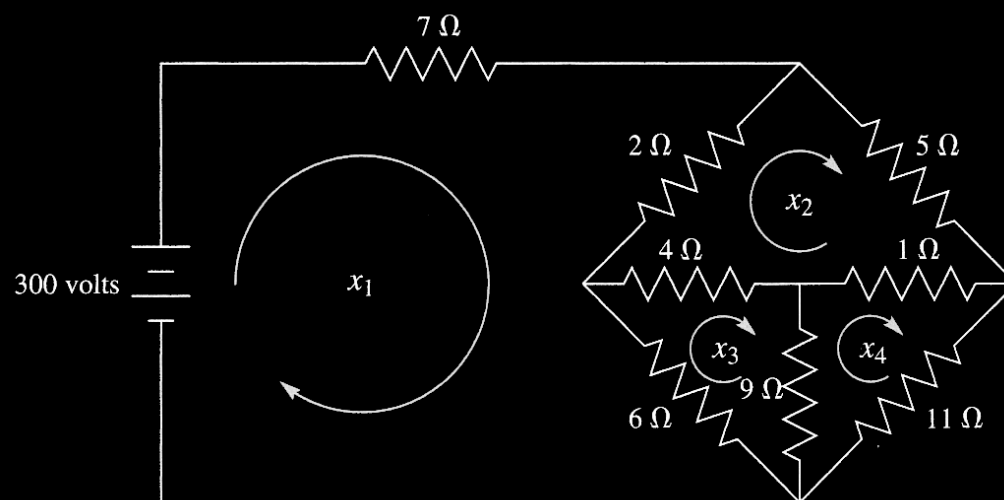


FIGURE 7.1
Electrical
network

$$Ax = b$$

$$A x = b$$

$$\left\{ \begin{array}{l} a_{11}x_1 + a_{12}x_2 + a_{13}x_3 + \cdots + a_{1n}x_n = b_1 \\ a_{21}x_1 + a_{22}x_2 + a_{23}x_3 + \cdots + a_{2n}x_n = b_2 \\ a_{31}x_1 + a_{32}x_2 + a_{33}x_3 + \cdots + a_{3n}x_n = b_3 \\ \vdots \quad \quad \quad \vdots \quad \quad \quad \vdots \quad \quad \quad \vdots \quad \quad \quad \vdots \\ a_{i1}x_1 + a_{i2}x_2 + a_{i3}x_3 + \cdots + a_{in}x_n = b_i \\ \vdots \quad \quad \quad \vdots \quad \quad \quad \vdots \quad \quad \quad \vdots \quad \quad \quad \vdots \\ a_{n1}x_1 + a_{n2}x_2 + a_{n3}x_3 + \cdots + a_{nn}x_n = b_n \end{array} \right.$$

$$A x = b$$

$$\sum_{j=1}^n a_{ij} x_j = b_i \quad (1 \leq i \leq n)$$

Naïve Gaussian Elimination

$$\begin{array}{rcl} 6x_1 - 2x_2 + 2x_3 + 4x_4 & = & 16 \\ 12x_1 - 8x_2 + 6x_3 + 10x_4 & = & 26 \\ 3x_1 - 13x_2 + 9x_3 + 3x_4 & = & -19 \\ -6x_1 + 4x_2 + x_3 - 18x_4 & = & -34 \end{array}$$

Naïve Gaussian Elimination

$$\begin{array}{rcll} 6x_1 - 2x_2 + 2x_3 + 4x_4 & = & 16 \\ -4x_2 + 2x_3 + 2x_4 & = & -6 \\ -12x_2 + 8x_3 + x_4 & = & -27 \\ 2x_2 + 3x_3 - 14x_4 & = & -18 \end{array}$$

Naïve Gaussian Elimination

$$6x_1 - 2x_2 + 2x_3 + 4x_4 = 16$$

$$-4x_2 + 2x_3 + 2x_4 = -6$$

$$2x_3 - 5x_4 = -9$$

$$4x_3 - 13x_4 = -21$$

Naïve Gaussian Elimination

$$\begin{array}{rcl} 6x_1 - 2x_2 + 2x_3 + 4x_4 & = & 16 \\ -4x_2 + 2x_3 + 2x_4 & = & -6 \\ 2x_3 - 5x_4 & = & -9 \\ -3x_4 & = & -3 \end{array}$$

This system is said to be in **upper triangular** form.

Back Substitution

This completes the first phase (**forward elimination**) in the Gaussian algorithm. The second phase (**back substitution**) will solve System (5) for the unknowns *starting at the bottom*. Thus, from the fourth equation, we obtain the last unknown

$$x_4 = \frac{-3}{-3} = 1$$

Putting $x_4 = 1$ in the third equation gives us

$$2x_3 - 5 = -9$$

and we find the next to last unknown

$$x_3 = \frac{-4}{2} = -2$$

and so on. The solution is

$$x_1 = 3 \quad x_2 = 1 \quad x_3 = -2 \quad x_4 = 1$$

Naïve Gaussian Elimination

$$\begin{bmatrix} a_{11} & a_{12} & a_{13} & \cdots & a_{1n} \\ a_{21} & a_{22} & a_{23} & \cdots & a_{2n} \\ a_{31} & a_{32} & a_{33} & \cdots & a_{3n} \\ \vdots & \vdots & \vdots & & \vdots \\ a_{i1} & a_{i2} & a_{i3} & \cdots & a_{in} \\ \vdots & \vdots & \vdots & & \vdots \\ a_{n1} & a_{n2} & a_{n3} & \cdots & a_{nn} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ \vdots \\ x_i \\ \vdots \\ x_n \end{bmatrix} = \begin{bmatrix} b_1 \\ b_2 \\ b_3 \\ \vdots \\ b_i \\ \vdots \\ b_n \end{bmatrix}$$

Naïve Gaussian Elimination

After the first step,

$$\begin{bmatrix} a_{11} & a_{12} & a_{13} & \cdots & a_{1n} \\ 0 & a_{22} & a_{23} & \cdots & a_{2n} \\ 0 & a_{23} & a_{33} & \cdots & a_{3n} \\ \vdots & \vdots & \vdots & & \vdots \\ 0 & a_{i2} & a_{i3} & \cdots & a_{in} \\ \vdots & \vdots & \vdots & & \vdots \\ 0 & a_{n2} & a_{n3} & \cdots & a_{nn} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ \vdots \\ x_i \\ \vdots \\ x_n \end{bmatrix} = \begin{bmatrix} b_1 \\ b_2 \\ b_3 \\ \vdots \\ b_i \\ \vdots \\ b_n \end{bmatrix}$$

Naïve Gaussian Elimination

first equation as the first **pivot equation** and to a_{11} as the first **pivot element**. For each of the remaining equations ($2 \leq i \leq n$), we compute

$$\begin{cases} a_{ij} \leftarrow a_{ij} - \left(\frac{a_{i1}}{a_{11}} \right) a_{1j} & (1 \leq j \leq n) \\ b_i \leftarrow b_i - \left(\frac{a_{i1}}{a_{11}} \right) b_1 \end{cases}$$

The symbol $\leftarrow$ indicates a *replacement*. Thus, the content of the memory location allocated to a_{ij} is replaced by $a_{ij} - (a_{i1}/a_{11})a_{1j}$, and so on. This is accomplished by the following line of pseudocode:

$$a_{ij} \leftarrow a_{ij} - (a_{i1}/a_{11})a_{1j}$$

Note that the quantities (a_{i1}/a_{11}) are the **multipliers**. The new coefficient of x_1 in the i th equation will be 0 because $a_{i1} - (a_{i1}/a_{11})a_{11} = 0$.

Continuing to second equation...

With the second equation as the pivot equation, we compute for each remaining equation ($3 \leq i \leq n$)

$$\begin{cases} a_{ij} \leftarrow a_{ij} - \left(\frac{a_{i2}}{a_{22}} \right) a_{2j} & (2 \leq j \leq n) \\ b_i \leftarrow b_i - \left(\frac{a_{i2}}{a_{22}} \right) b_2 \end{cases}$$

k^{th} step of forward elimination

$$\begin{bmatrix} a_{11} & a_{12} & a_{13} & \cdots & \cdots & \cdots & a_{1n} \\ 0 & a_{22} & a_{23} & \cdots & \cdots & \cdots & a_{2n} \\ 0 & 0 & a_{33} & \cdots & \cdots & \cdots & a_{3n} \\ \vdots & \vdots & \vdots & \ddots & & & \vdots \\ 0 & 0 & 0 & \cdots & a_{kk} & \cdots & a_{kn} \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 0 & \cdots & a_{ik} & \cdots & a_{in} \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\ 0 & 0 & 0 & \cdots & a_{nk} & \cdots & a_{nn} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ \vdots \\ x_k \\ \vdots \\ x_i \\ \vdots \\ x_n \end{bmatrix} = \begin{bmatrix} b_1 \\ b_2 \\ b_3 \\ \vdots \\ b_k \\ \vdots \\ b_i \\ \vdots \\ b_n \end{bmatrix}$$

Naïve GE: Pseudocode

```
integer  $i, j, k$ ;   real array  $(a_{ij})_{1:n \times 1:n}, (b_i)_{1:n}$   
for  $k = 1$  to  $n - 1$  do  
    for  $i = k + 1$  to  $n$  do  
        for  $j = k$  to  $n$  do  
             $a_{ij} \leftarrow a_{ij} - (a_{ik}/a_{kk})a_{kj}$   
        end for  
         $b_i \leftarrow b_i - (a_{ik}/a_{kk})b_k$   
    end for  
end for
```

Storing Multipliers (instead of zeros)

$$a_{ik} \leftarrow a_{ik} - (a_{ik}/a_{kk})a_{kk}$$

Since we expect this to be 0, no purpose is served in computing it. The location where the 0 is being created is a good place to store the multiplier. If these remarks are put into practice, the pseudocode will look like this:

```
integer i, j, k;  real xmult;  real array (aij)1:n × 1:n, (bi)1:n  
for k = 1 to n - 1 do  
  for i = k + 1 to n do  
    xmult ← aik / akk  
    aik ← xmult  
    for j = k + 1 to n do  
      aij ← aij - (xmult)akj  
    end for  
    bi ← bi - (xmult)bk  
  end for  
end for
```

Back Substitution

At the beginning of the back substitution phase, the linear system is of the form

$$\left\{ \begin{array}{llll} a_{11}x_1 + a_{12}x_2 + a_{13}x_3 + \cdots & \cdots + & a_{1n}x_n = b_1 \\ & \cdots + & a_{2n}x_n = b_2 \\ & & + & a_{3n}x_n = b_3 \\ & & & \vdots \\ & & & \vdots \\ & a_{ii}x_i + a_{i,i+1}x_{i+1} + & \cdots + & a_{in}x_n = b_i \\ & & & \vdots \\ & & & \vdots \\ & & a_{n-1,n-1}x_{n-1} + a_{n-1,n}x_n = b_{n-1} \\ & & & a_{nn}x_n = b_n \end{array} \right.$$

Back Substitution

The back substitution starts by solving the n th equation for x_n :

$$x_n = \frac{b_n}{a_{nn}}$$

Then, using the $(n - 1)$ th equation, we solve for x_{n-1} :

$$x_{n-1} = \frac{1}{a_{n-1,n-1}} (b_{n-1} - a_{n-1,n}x_n)$$

Back Substitution

We continue working upward, recovering each x_i by the formula

$$x_i = \frac{1}{a_{ii}} \left(b_i - \sum_{j=i+1}^n a_{ij} x_j \right) \quad (i = n - 1, n - 2, \dots, 1)$$

Here is pseudocode to do this:

```
integer  $i, j, n$ ;   real  $sum$ ;   real array  $(a_{ij})_{1:n \times 1:n}, (x_i)_{1:n}$   
 $x_n \leftarrow b_n / a_{nn}$   
for  $i = n - 1$  to  $1$  step  $-1$  do  
     $sum \leftarrow b_i$   
    for  $j = i + 1$  to  $n$  do  
         $sum \leftarrow sum - a_{ij} x_j$   
    end for  
     $x_i \leftarrow sum / a_{ii}$   
end for
```


Summary

(1) The basic **forward elimination** procedure using equation k to operate on equations $k + 1, k + 2, \dots, n$ is

$$\begin{cases} a_{ij} \leftarrow a_{ij} - (a_{ik}/a_{kk})a_{kj} & (k \leq j \leq n, k < i \leq n) \\ b_i \leftarrow b_i - (a_{ik}/a_{kk})b_k \end{cases}$$

Here we assume $a_{kk} \neq 0$. The basic **back substitution** procedure is

$$x_i = \frac{1}{a_{ii}} \left(b_i - \sum_{j=i+1}^n a_{ij}x_j \right) \quad (i = n - 1, n - 2, \dots, 1)$$

(2) When solving the linear system $A\mathbf{x} = \mathbf{b}$, if the true or exact solution is $\mathbf{x}$ and the approximate or computed solution is $\tilde{\mathbf{x}}$, then important quantities are

error vectors	$\mathbf{e} = \tilde{\mathbf{x}} - \mathbf{x}$
residual vectors	$\mathbf{r} = A\tilde{\mathbf{x}} - \mathbf{b}$

LU Factorization Example (supplements chalkboard)

Numerical Example

The system of Equations (2) of Section 7.1 can be written succinctly in matrix form:

$$\begin{bmatrix} 6 & -2 & 2 & 4 \\ 12 & -8 & 6 & 10 \\ 3 & -13 & 9 & 3 \\ -6 & 4 & 1 & -18 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} = \begin{bmatrix} 16 \\ 26 \\ -19 \\ -34 \end{bmatrix} \quad (2)$$

Furthermore, the operations that led from this system to Equation (5) of Section 7.1, that is, the system

$$\begin{bmatrix} 6 & -2 & 2 & 4 \\ 0 & -4 & 2 & 2 \\ 0 & 0 & 2 & -5 \\ 0 & 0 & 0 & -3 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} = \begin{bmatrix} 16 \\ -6 \\ -9 \\ -3 \end{bmatrix} \quad (3)$$

Matrix Interpretation

$$MAx = Mb \quad (4)$$

where M is a matrix chosen so that MA is the coefficient matrix for System (3). Hence, we have

$$MA = \begin{bmatrix} 6 & -2 & 2 & 4 \\ 0 & -4 & 2 & 2 \\ 0 & 0 & 2 & -5 \\ 0 & 0 & 0 & -3 \end{bmatrix} \equiv U$$

which is an upper triangular matrix.

Column 1

The first step of naive Gaussian elimination results in Equation (3) of Section 7.1 or the system

$$\begin{bmatrix} 6 & -2 & 2 & 4 \\ 0 & -4 & 2 & 2 \\ 0 & -12 & 8 & 1 \\ 0 & 2 & 3 & -14 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} = \begin{bmatrix} 16 \\ -6 \\ -27 \\ -18 \end{bmatrix}$$

This step can be accomplished by multiplying (1) by a lower triangular matrix M_1 :

$$M_1 A x = M_1 b$$

where

$$M_1 = \begin{bmatrix} 1 & 0 & 0 & 0 \\ -2 & 1 & 0 & 0 \\ -\frac{1}{2} & 0 & 1 & 0 \\ 1 & 0 & 0 & 1 \end{bmatrix}$$

Column 2

To continue, step 2 resulted in Equation (4) of Section 7.1 or the system

$$\begin{bmatrix} 6 & -2 & 2 & 4 \\ 0 & -4 & 2 & 2 \\ 0 & 0 & 2 & -5 \\ 0 & 0 & 4 & -13 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} = \begin{bmatrix} 16 \\ -6 \\ -9 \\ -21 \end{bmatrix}$$

which is equivalent to

$$\mathbf{M}_2 \mathbf{M}_1 \mathbf{A} \mathbf{x} = \mathbf{M}_2 \mathbf{M}_1 \mathbf{b}$$

where

$$\mathbf{M}_2 = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & -3 & 1 & 0 \\ 0 & \frac{1}{2} & 0 & 1 \end{bmatrix}$$

Again, $\mathbf{M}_2$ differs from an identity matrix by the presence of the negatives of the multipliers in the second column from the diagonal down.

Column 3

$$\mathbf{M}_3\mathbf{M}_2\mathbf{M}_1\mathbf{A}\mathbf{x} = \mathbf{M}_3\mathbf{M}_2\mathbf{M}_1\mathbf{b}$$

where

$$\mathbf{M}_3 = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & -2 & 1 \end{bmatrix}$$

Now the forward elimination phase is complete, and with

$$\mathbf{M} = \mathbf{M}_3\mathbf{M}_2\mathbf{M}_1$$

we have the upper triangular coefficient System (3).

Recognizing “L”

$$\begin{aligned} A &= M^{-1}U \\ &= M_1^{-1}M_2^{-1}M_3^{-1}U \\ &= LU \end{aligned}$$

Since each M_k has such a special form, its inverse is obtained by simply changing the sign of the negative multiplier entries! Hence, we have

$$\begin{aligned} L &= \begin{bmatrix} 1 & 0 & 0 & 0 \\ 2 & 1 & 0 & 0 \\ \frac{1}{2} & 0 & 1 & 0 \\ -1 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 3 & 1 & 0 \\ 0 & -\frac{1}{2} & 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 2 & 1 \end{bmatrix} \\ &= \begin{bmatrix} 1 & 0 & 0 & 0 \\ 2 & 1 & 0 & 0 \\ \frac{1}{2} & 3 & 1 & 0 \\ -1 & -\frac{1}{2} & 2 & 1 \end{bmatrix} \end{aligned}$$

It is somewhat amazing that L is a unit lower triangular matrix composed of the multipliers. Notice that in forming L , we did not determine M first and then compute $M^{-1} = L$. (Why

Verifying $LU=A$

It is easy to verify that

$$\begin{aligned}LU &= \begin{bmatrix} 1 & 0 & 0 & 0 \\ 2 & 1 & 0 & 0 \\ \frac{1}{2} & 3 & 1 & 0 \\ -1 & -\frac{1}{2} & 2 & 1 \end{bmatrix} \begin{bmatrix} 6 & -2 & 2 & 4 \\ 0 & -4 & 2 & 2 \\ 0 & 0 & 2 & -5 \\ 0 & 0 & 0 & -3 \end{bmatrix} \\ &= \begin{bmatrix} 6 & -2 & 2 & 4 \\ 12 & -8 & 6 & 10 \\ 3 & -13 & 9 & 3 \\ -6 & 4 & 1 & -18 \end{bmatrix} = A\end{aligned}$$

We see that A is **factored** or **decomposed** into a unit lower triangular matrix L and an upper triangular matrix U .