

CS 312

Spring 2008

Lecture 28

Logic programming

Wrap-up

Declarative vs. imperative

- Imperative programming: tell computer how to change its state to accomplish a result
- **Declarative** programming: tell computer what you want computed, without specifying state changes
 - Avoids side effects, enables analysis and optimization
 - functional programming: give an expression equal to the desired result
 - **logic programming**: give a logical formula describing what should be true of the result
 - a simple version: database queries

Logic programming in Prolog

Programmer defines **boolean**-valued **predicates**

- Language figures out all ways to make predicates true.
- Example (syntax modified from Prolog)

```
parent(X,Y) <= father(X,Y).
```

```
parent(X,Y) <= mother(X,Y).
```

```
father(bob, alice).           i.e., <= true
```

```
?- parent(bob, X).
```

```
X = alice
```

```
?- parent(X, X).
```

```
No
```

```
sibling(X,Y) <= parent(Z,X), parent(Z,Y).
```

```
father(bob, charlie).
```

```
?- sibling(alice, X).
```

```
X = alice
```

```
X = charlie
```

Concatenating lists

- Goal: define predicate `join(L1, L2, L3)` meaning `L1@L2 = L3`.
- If `T1 @ L2 = T3`, then `H1::T1@L2 = H1::T3`. So:

```
join([], L2, L2).
```

```
join(H1::T1, L2, H1::T3) <= join(T1,L2,T3).
```

```
?- join([1,2,3], [4,5,6], X)
```

```
  X = [1,2,3,4,5,6]
```

```
?- join([1,X,3], 4::Y, [1,2,Z,W,5,6])
```

```
  X = 2, Y = [5,6], Z = 3, W = 4
```

```
?- join([1,X,X], [Y,Y], [X,X,Y,Y,Y])
```

What did we cover?

Goal: better software design and implementation

- New programming paradigms
 - higher-order functions, pattern matching, polymorphism, concurrency, ...
- Specifying functions and data abstractions
- Reasoning about correctness
 - using specifications, logic
- Reasoning about performance
 - asymptotic complexity, recurrences, amortized complexity, locality
- Important data structures and algorithms
 - balanced binary trees, hash tables, splay trees, B-trees, functional impls

Life after 312

- SML is fun and ML variants (SML, OCaml, Haskell) are used in some “real-world” apps.
- Functional *style* is useful in almost any language.
- Most course material is not specific to SML:
 - Specifications, AF, RI, logic and verification
 - Recurrences and complexity analysis
 - Data structures and algorithms
- What if you miss functional programming?

Simulating functions with objects

- First-class functions can be simulated with first-class objects.

`val f: t -> t' = fn(x:t) => e` is similar to:

```
class Fn {  
    t' apply(t x) { return e; }  
}
```

```
Fn f = new Fn();
```

`f(x)` is translated to `f.apply(x)`

- Java nested classes can even mention variables from containing scope.
- C# supports first-class functions directly (delegates)

Pattern matching

- Pattern matching is not supported by object-oriented languages
- Problem: matching type T requires knowing exactly what T is.
 - Doesn't work with abstract types -- conflicts with data abstraction
 - Could not expose pattern matching in SML signatures
- Can we have a pattern-matching mechanism that works with objects and data abstraction?

JMatch: Java + pattern matching

- JMatch supports **predicate methods** with multiple **modes** capturing directions of computation

```
class List {  
  Object head; List tail;  
  List(Object h, List t) returns (h, t)  
    (head = h & tail = t)  
}
```

- **Forward mode:** creates an object. **Backward mode:** pattern matches, binds h and t:

```
switch (lst) {  
  case List(l, List(Object x, List rest)):  
    return List(x, f(rest))  
}
```

JMatch logic programming

- A limited form of logic programming!

```
List join(List x, List y) returns(x) returns(y) (  
    x = List(hx, tx) &  
    tr = join(tx, y) &  
    result = List(hx, tr)  
)
```

```
let List(1, List(2, null)) = join(prefix, List y);  
... use y here ...
```

Rebalancing a red-black tree in JMatch

```
static RBNode balance(int color, int value,
                    RBTree left, RBTree right) {

    if (color == BLACK) {
        switch (value, left, right) {
            case int z, RBNode(RED, int y,
                RBNode(RED, int x, RBTree a, RBTree b), RBTree c),
                RBTree d:
            case z, RBNode(RED, x, a, RBNode(RED, y, b, c)), d:
            case x, c, RBNode(RED, z, RBNode(RED, y, a, b), d):
            case x, a, RBNode(RED, y, b, RBNode(RED, z, c, d)):
                return RBNode(RED, y,
                    RBNode(BLACK, x, a, b), RBNode(BLACK, z, c, d));
        }
    }
    return RBNode(color, value, left, right);
}
```

Iteration

- Logic programming has iteration built in.

```
class RBNode implements IntCollection, Tree {
    RBTREE left, right; int value; boolean color;
    boolean contains(int x) iterates(x) (
        x < value && left.contains(value) ||
        x = value ||
        x > value && right.contains(value)
    )
}
```

- Forward mode: usual BST lookup
- Backward mode: in-order tree traversal!

```
foreach (tree.contains(int x) & x < 10) {
    ... use x ...
}
```

The tree iterator in Java

```
class Treeliterator implements Iterator {

    Iterator subiterator;
    boolean hasNext;
    Object current;
    int state;
    // states:
    // 1. Iterating through left child.
    // 2. Just yielded current node value
    // 3. Iterating through right child

    Treeliterator() {
        subiterator = RBTree.this.left.iterator();
        state = 1;
        preloadNext();
    }

    public boolean hasNext() {
        return hasNext;
    }

    public Object next() {
        if (!hasNext) throw new NoSuchElementException();
        Object ret = current;

        private void preloadNext() {
            loop: while (true) {
                switch (state) {
                    case 1:
                    case 3:
                        hasNext = true;
                        if (subiterator.hasNext()) {
                            current = subiterator.next();
                            return;
                        } else {
                            if (state == 1) {
                                state = 2;
                                current = RBTree.this.value;
                                return;
                            } else {
                                hasNext = false;
                                return;
                            }
                        }
                    case 2:
                        subiterator = RBTree.right.iterator();
                        state = 3;
                        continue loop;
                }
            }
        }
    }
}
```


Conclusions

- Object-oriented languages are incorporating many functional programming language features (higher-order functions, polymorphism, lexical scoping...)
- Pattern matching may show up too!

Final exam

- May 13, 7-9:30pm, Phillips 203
- Open book
- Cumulative

Follow-on courses

- Complexity: CS 381
- Understanding programming paradigms and language features: CS 411, CS 611
- Language implementation: CS 412/413
- Algorithms and algorithm design: CS 482
- Logic: CS 486
- Think about participating in 312 (and in other courses) as a course consultant