

CS 3110

Mutable Data Types

A New Despair
Mutability Strikes Back
Return of Imperative Programming

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Today's music: *The Imperial March*
from the soundtrack to *Star Wars, Episode V: The Empire Strikes Back*

Review

Previously in 3110:

- Advanced data structures
 - Streams and laziness
 - Balanced binary trees



Today: THE DARK SIDE ARRIVES

- Mutable data types: refs, mutable fields, (arrays)

REFS

Demo

References

- Aka “refs” or “ref cell”
- **Pointer** to a typed location in memory
- The binding of a variable to a pointer is **immutable** but the **contents of the memory** may change

References

- **Syntax:** **ref** **e**
- **Evaluation:**
 - Evaluate **e** to a value **v**
 - Allocate a new *location* **loc** in memory to hold **v**
 - Store **v** in **loc**
 - Return **loc**
 - Note: locations are values; can pass and return from functions
- **Type checking:**
 - New type constructor: **t ref** where **t** is a type
 - Note: **ref** is used as keyword in type and as keyword in value
 - **ref e : t ref** if **e : t**

References

- **Syntax:** $e1 := e2$
- **Evaluation:**
 - Evaluate $e2$ to a value $v2$
 - Evaluate $e1$ to a location loc
 - Store $v2$ in loc
 - Return $()$
- **Type checking:**
 - If $e2 : t$
 - and $e1 : t \text{ ref}$
 - then $e1 := e2 : \text{unit}$

References

- **Syntax: $!e$**
 - note: not negation
- **Evaluation:**
 - Evaluate e to loc
 - Return contents of loc
- **Type checking:**
 - If $e : t \text{ ref}$
 - then $!e : t$

Aliases

References may create **aliases**:

```
let x = ref 42
let y = ref 42
let z = x
let () = x := 43
let w = (!y) + (!z)
```

z and **x** are aliases

Equality

- Suppose we have two refs...
 - `let r1 = ref 3110`
 - `let r2 = ref 3110`
- Double equals is *physical equality*
 - `r1 == r1`
 - `r1 != r2`
- Single equals is *structural equality*
 - `r1 = r1`
 - `r1 = r2`
 - `ref 3110 <> ref 2110`
- You usually want single equals

EXAMPLE: COUNTER

Semicolon

- **Syntax:** $e1; e2$
- **Evaluation:**
 - Evaluate $e1$ to a value $v1$
 - Then **throw away** that value
(note: $e1$ could have side effects)
 - evaluate $e2$ to a value $v2$
 - return $v2$
- **Type checking:**
 - If $e1 : \text{unit}$
 - and $e2 : t$
 - then $e1; e2 : t$

Scope matters

```
(* correct *)  
let next_val =  
  let counter = ref 0  
  in fun () ->  
    incr counter;  
    !counter
```

```
(* faulty *)  
let next_val = fun () ->  
  let counter = ref 0  
  in incr counter;  
  !counter
```

MUTABLE FIELDS

Demo

Implementing refs

Ref cells are essentially syntactic sugar:

```
type 'a ref = { mutable contents: 'a }  
let ref x = { contents = x }  
let ( ! ) r = r.contents  
let ( := ) r newval = r.contents <- newval
```

- That type is declared in **Pervasives**
- The functions are compiled down to something equivalent



BEWARE

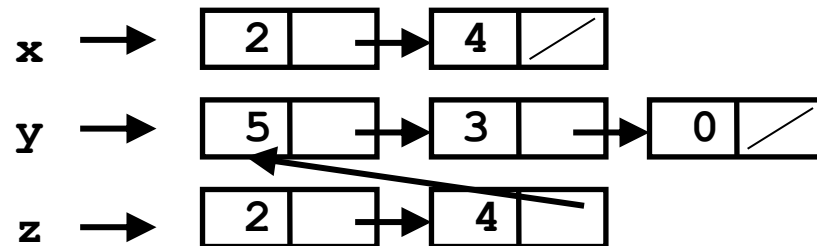
Immutable lists

We have never needed to worry about aliasing with lists!

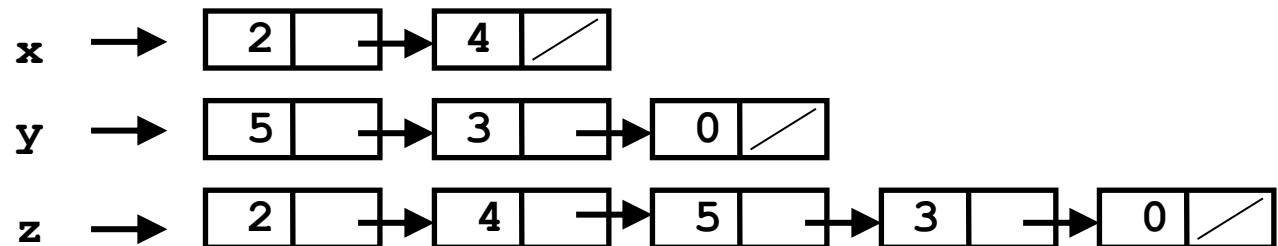
```
let x = [2;4]
```

```
let y = [5;3;0]
```

```
let z = x @ y
```



vs.



(no code you write could ever tell, but OCaml implementation uses the first one)

OCaml:

blissfully unaware of aliasing

Java:

obsession with aliasing

Faulty code

```
class ProtectedResource {
    private Resource theResource = ...;
    private String[] allowedUsers = ...;
    public String[] getAllowedUsers() {
        return allowedUsers;
    }
    public String currentUser() { ... }
    public void useTheResource() {
        for(int i=0; i < allowedUsers.length; i++) {
            if(currentUser().equals(allowedUsers[i])) {
                ... // access allowed: use it
                return;
            }
        }
        throw new IllegalAccessException();
    }
}
```

Discussion: Can you find the security fault?

Have to make copies

The problem:

```
p.getAllowedUsers()[0] = p.currentUser();  
p.useTheResource();
```

The fix:

```
public String[] getAllowedUsers() {  
    ... return a copy of allowedUsers ...  
}
```

Similar errors as recent as Java 1.7beta

Benefits of immutability

- Programmer doesn't have to think about aliasing; can concentrate on other aspects of code
- Language implementation is free to use aliasing, which is cheap
- Often easier to reason about whether code is correct
- Perfect fit for concurrent programming

But there are downsides:

- I/O is fundamentally about mutation
- Some data structures (hash tables, arrays, ...) are more efficient if imperative

Try not to abuse your new-found power!

Upcoming events

- N/A

This is (reluctantly) imperative.

THIS IS 3110