

#### Verification in Coq

Prof. Clarkson Fall 2017

Today's music: Check Yo Self by Ice Cube

#### Review

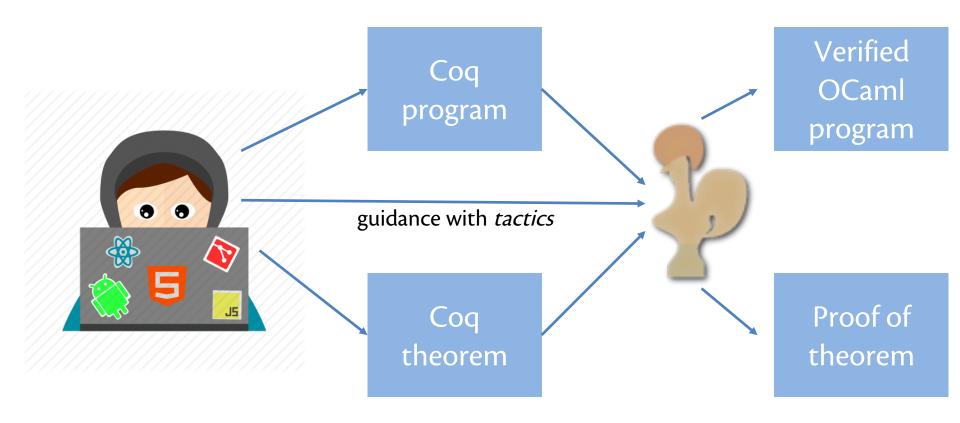
#### Previously in 3110:

- Functional programming in Coq
- Logic in Coq
- Curry-Howard correspondence (proofs are programs)
- Induction in Coq

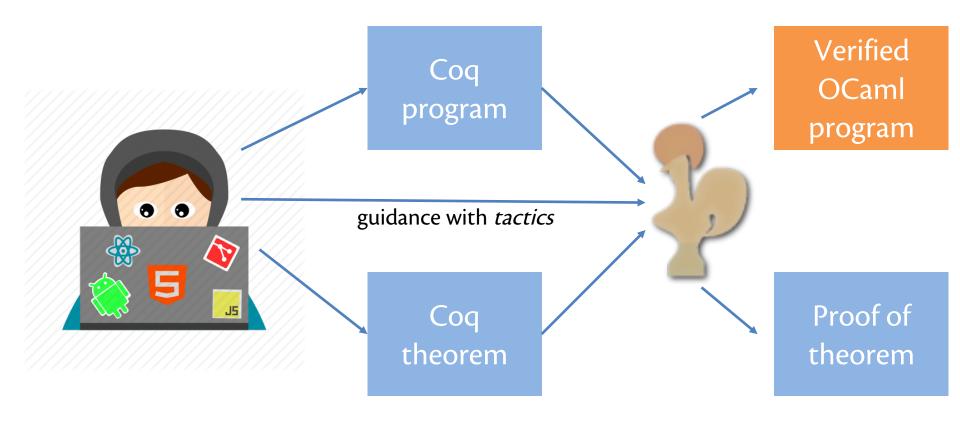
#### **Today:** Verification of...

- Functions
- Data structures
- Compilers

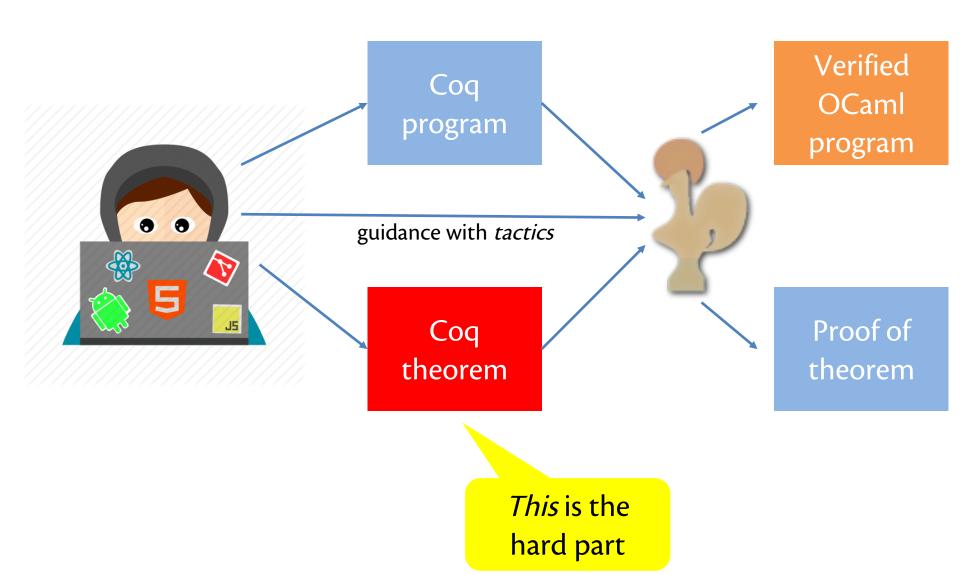
## Coq for program verification



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## Coq for program verification



#### Theorems and test cases

- Do I have the right ones?
- Do I have enough?
- What am I missing?

... there are no great answers to these questions, only methodologies that help

Prove that precondition implies postcondition

#### **VERIFICATION OF A FUNCTION**

#### **Factorial**

- Precondition: n >= 0
- Postcondition: fact n = n!

Problem: how to express! in Coq?

#### **Factorial**

```
Fixpoint fact (n:nat) :=
  match n with
  | 0 => 1
  | S k => n * (fact k)
  end.
```

Theorem fact\_correct : forall n, fact n = fact n.

vacuous

#### Tail-recursive factorial

```
Fixpoint fact_tr_acc (n:nat) (acc:nat) :=
  match n with
  | 0 => acc
  | S k => fact_tr_acc k (n * acc)
  end.
```

Definition fact\_tr (n:nat) :=
 fact\_tr\_acc n 1.

Precondition: n >= 0
Postcondition: fact\_tr n = fact n

actually
unnecessary
because nat
already
implies it

# **Verify factorial**

```
Lemma helper: forall (n acc: nat),
  fact tr acc n acc = (fact n) * acc.
Proof.
  intros n.
  induction n as [ | k IH]; intros acc.
  - simpl. ring.
  - simpl. rewrite IH. ring.
Qed.
Theorem fact tr correct : forall n:nat,
  fact tr n = fact n.
Proof.
  intros n. unfold fact tr. rewrite helper. ring.
Qed.
```

## **Verify factorial**

Generalized inductive hyopothesis: not all variables introduced

```
fact tr acc n acc = (f \cdot c \cdot n) * acc.
Proof.
  intros n.
  induction n as [ | k IH]; intros acc.
  - simpl. ring.
                                   Verify efficient impl equiv.
  - simpl. rewrite IH. ring.
                                     to "obviously correct"
Qed.
                                       inefficient impl.
Theorem fact tr correct: forall n:nat,
  fact tr n = fact n.
Proof.
  intros n. unfold fact tr. rewrite helper. ring.
Qed.
```

unfold tactic instantiates definition

#### **Extract verified factorial**

Coq nat becomes
OCamlint

Extract Inductive nat

Extract Inlined Constant Init.Nat.mul

Extraction "fact.ml" fact tr.

Coq \* becomes OCaml \*

Extract Coq to OCaml

Prove that equations hold for operations

# VERIFICATION OF A DATA STRUCTURE

#### Stack

```
module type Stack = sig
  type 'a t
  val empty
              : 'a t
  val is empty : 'a t -> bool
               : 'a t -> int
  val size
               : 'a t -> 'a option
  val peek
               : 'a -> 'a t -> 'a t
  val push
               : 'a t -> 'a t option
  val pop
end
```

## Categories of operations

- Creator: creates value of type "from scratch" without any inputs of that type
- Producer: takes value of type as input and returns value of type as output
- Observer: takes value of type as input but does not return value of type as output
- (Mutator: takes value of type as input and mutates the value)

#### Stack

```
module type Stack = sig
                            creator
  type 'a t
                 : 'a t
  val empty
  val is empty
                : 'a t -> bool
                                   observers
                 : 'a t -> int
  val size
                 : 'a t -> 'a option
  val peek
                 : 'a -> 'a t -> 'a t
  val push
                 : 'a t -> 'a t option
  val pop
end
                          producers
```

## Stack eqn. specification

```
• is empty empty = true
• is empty (push ) = false
• peek empty = None
• peek (push x ) = Some x
• size empty = 0
• size (push s) = 1 + \text{size s}
pop empty = None
• pop (push s) = Some s
```

## **Equational specification**

- aka algebraic specification
- Set of equations
- Describes interactions between:
  - observers and creators
  - observers and producers
  - producers and creators
  - producers and other producers
- Might not have equation for every possible interaction, because some might not be meaningful

#### Stack as list

```
Module MyStack.
Definition stack (A:Type) := list A.
Definition empty {A:Type} : stack A :=
 nil.
Definition is_empty {A:Type} (s : stack A)
  : bool
:=
 match s with
  nil => true
  :: => false
  end.
```

#### Stack as list

```
Definition push {A:Type} (x : A) (s : stack A)
  : stack A
:=
  X::S.
Definition peek {A:Type} (s : stack A)
  : option A
:=
 match s with
   nil => None
  | x::_ => Some x
  end.
```

#### Stack as list

```
Definition pop {A:Type} (s : stack A)
  : option (stack A)
:=
 match s with
   nil => None
  ::xs => Some xs
  end.
Definition size {A:Type} (s : stack A)
  : nat
:=
  length s.
End MyStack.
```

# Verify stack as list

```
Theorem empty is empty: forall (A:Type),
  @is empty A empty = true.
Proof. auto. Oed.
Theorem push not empty: forall (A:Type) (x:A) (s: stack A),
  is\_empty(push x s) = false.
Proof. auto. Oed.
Theorem peek empty: forall (A:Type),
  @peek A empty = None.
Proof. auto. Oed.
Theorem peek push : forall (A:Type) (x:A) (s : stack A),
 peek(push x s) = Some x.
Proof. auto. Oed.
```

# Verify stack as list

```
Theorem pop empty: forall (A:Type),
  @pop A empty = None.
Proof. auto. Oed.
Theorem pop push : forall (A:Type) (x:A) (s : stack A),
 pop(push x s) = Some s.
Proof. auto. Oed.
Theorem size empty: forall (A:Type),
  @size A empty = 0.
Proof. auto. Oed.
Theorem size push: forall (A:Type) (x:A) (s: stack A),
  size(push x s) = 1 + size s.
Proof. auto. Oed.
```

#### **Extract verified stack**

```
Extract Inductive bool => "bool" [ "true" "false" ].
Extract Inductive option => "option" [ "Some" "None" ].
Extract Inductive list => "list" [ "[]" "(::)" ].
Extract Inductive nat => int [ "0" "succ" ].

Extraction "stacks.ml" MyStack.
```

Coq bool, option, list, nat become OCaml equiv.

Prove that meaning is preserved

#### **VERIFICATION OF A COMPILER**

## **Expressions**

```
Inductive expr : Type :=
   Const : nat -> expr
  | Plus : expr -> expr -> expr.
Fixpoint eval expr (e : expr) : nat :=
  match e with
   Const n \Rightarrow n
  | Plus e1 e2 =>
    plus (eval expr e1) (eval expr e2)
  end.
```

## Stack programs

Definition prog := list instr.

Definition stack := list nat.

## Stack programs

```
Fixpoint eval prog
  (p : prog) (s : stack)
  : option stack
:=
 match p,s with
  (PUSH n)::p', s =>
   eval prog p' (n::s)
  ADD::p', x::y::s' =>
    eval prog p' ((x+y)::s')
   nil, s => Some s
  , => None
  end.
```

## Compiler

```
Fixpoint compile (e : expr) : prog :=
  match e with
  | Const n => [PUSH n]
  | Plus e1 e2 =>
    compile e2 ++ compile e1 ++ [ADD]
  end.
```

## Verify the compiler

```
Theorem compile_correct:
  forall (e:expr),
    eval_prog (compile e) []
    = Some [eval_expr e].
```

Proof in lecture code.

### **Extract verified compiler**

```
Extract Inlined Constant app => "(@)".
Extraction "compiler.ml" compile.
```

## **Upcoming events**

- [Thur or Fri] A5 out
- [Thur Tue] Design review meetings

This is verified.

**THIS IS 3110**