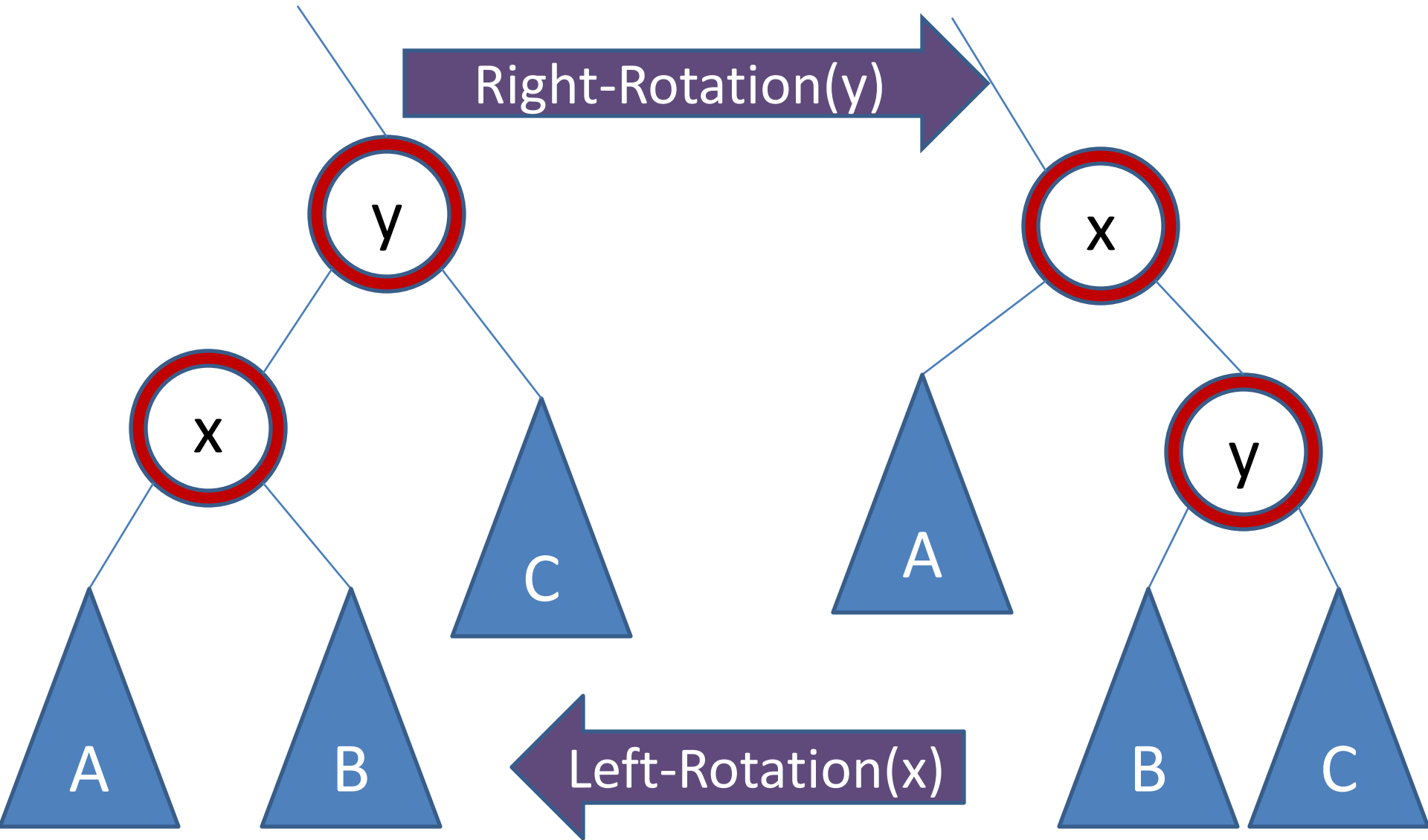


# 3110: Red-Black Trees

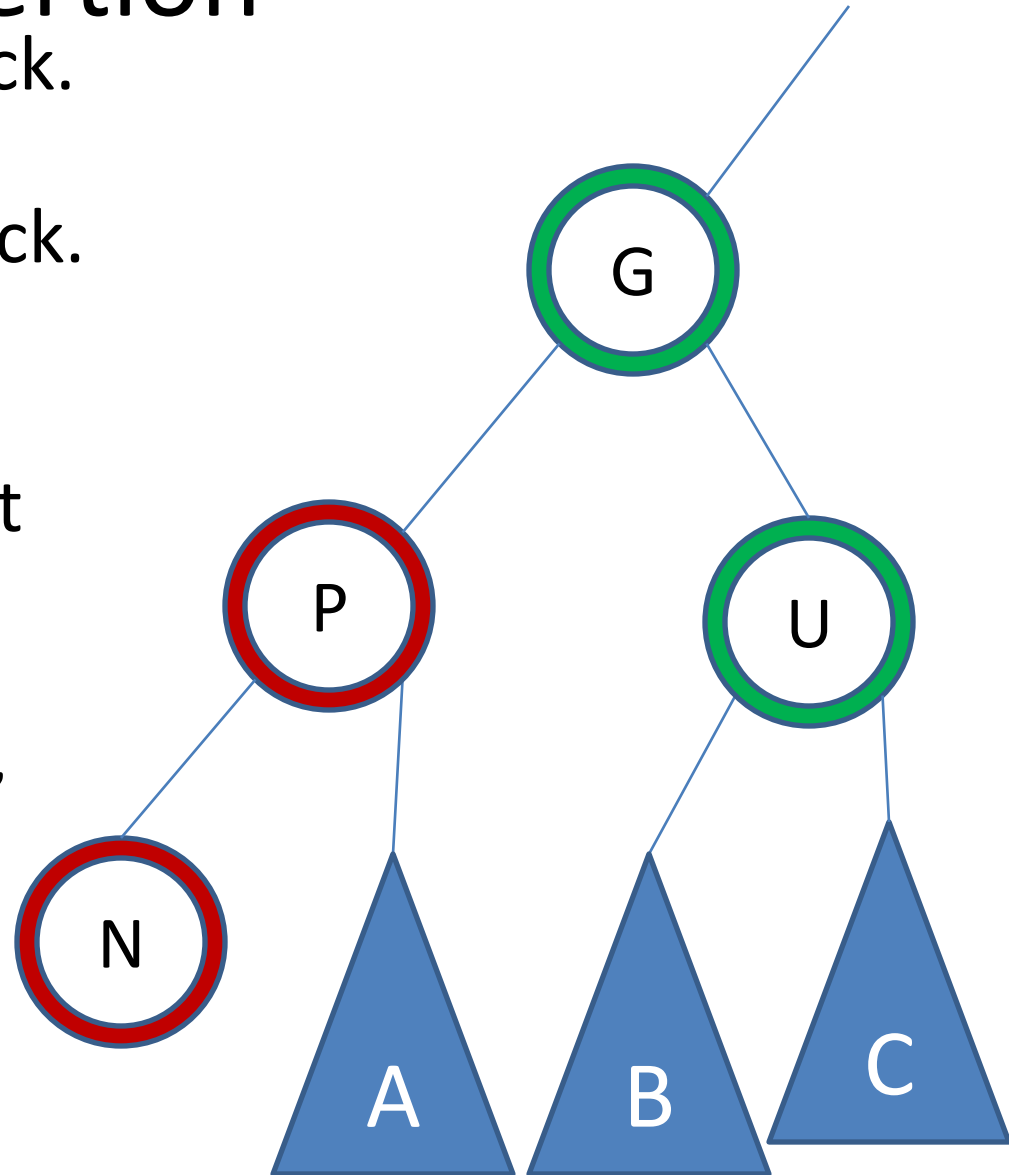
- Project?
- Homework – answer.
- Insertion
- An allude to Deletion

# Rotations

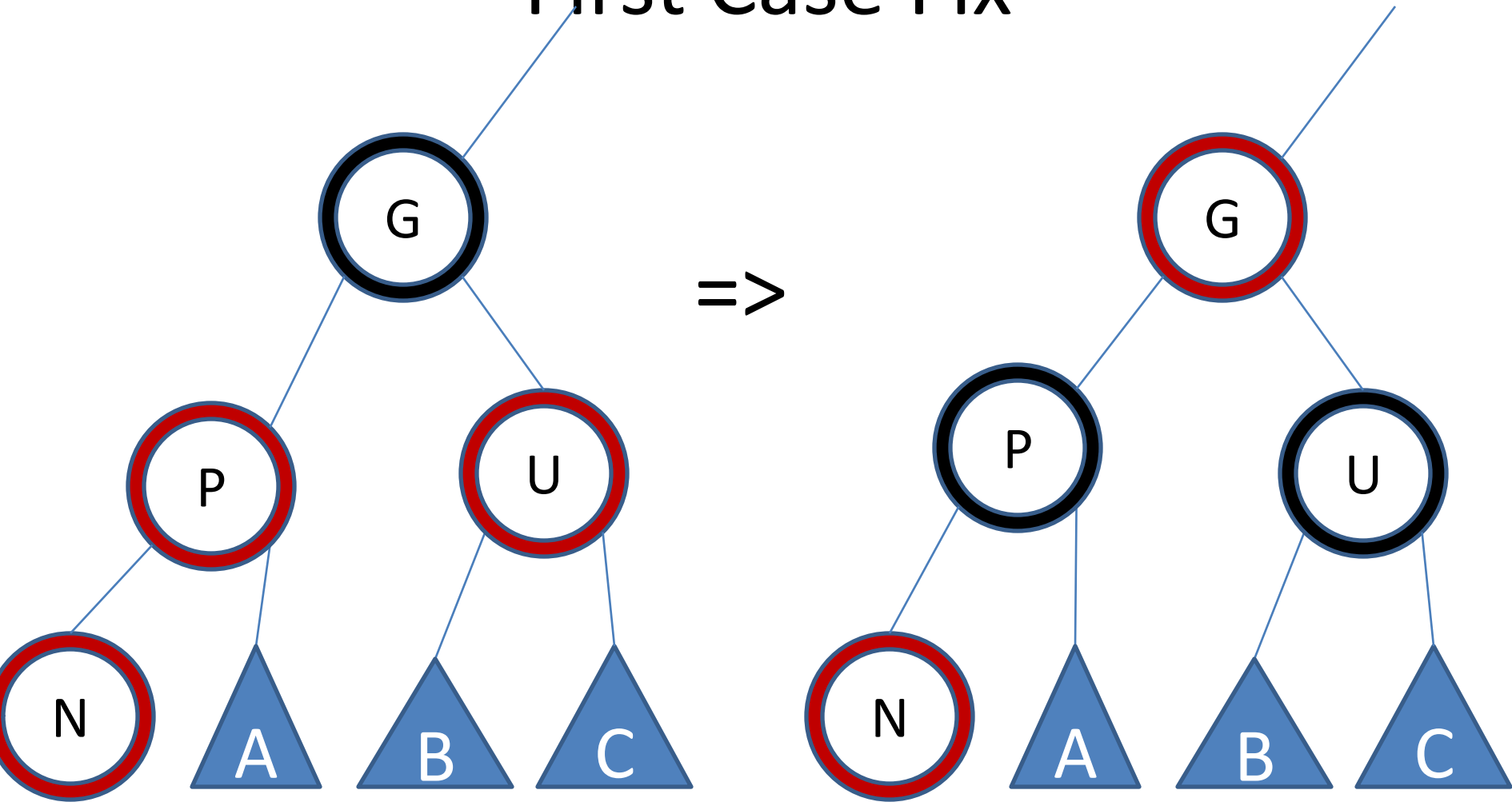


# Insertion

- If inserting root, color black.  
Done.
- If parent of inserted is black.  
Done.
- Otherwise have to break it down into cases.
- Consider the structure of new, parent, grandparent, and uncle:

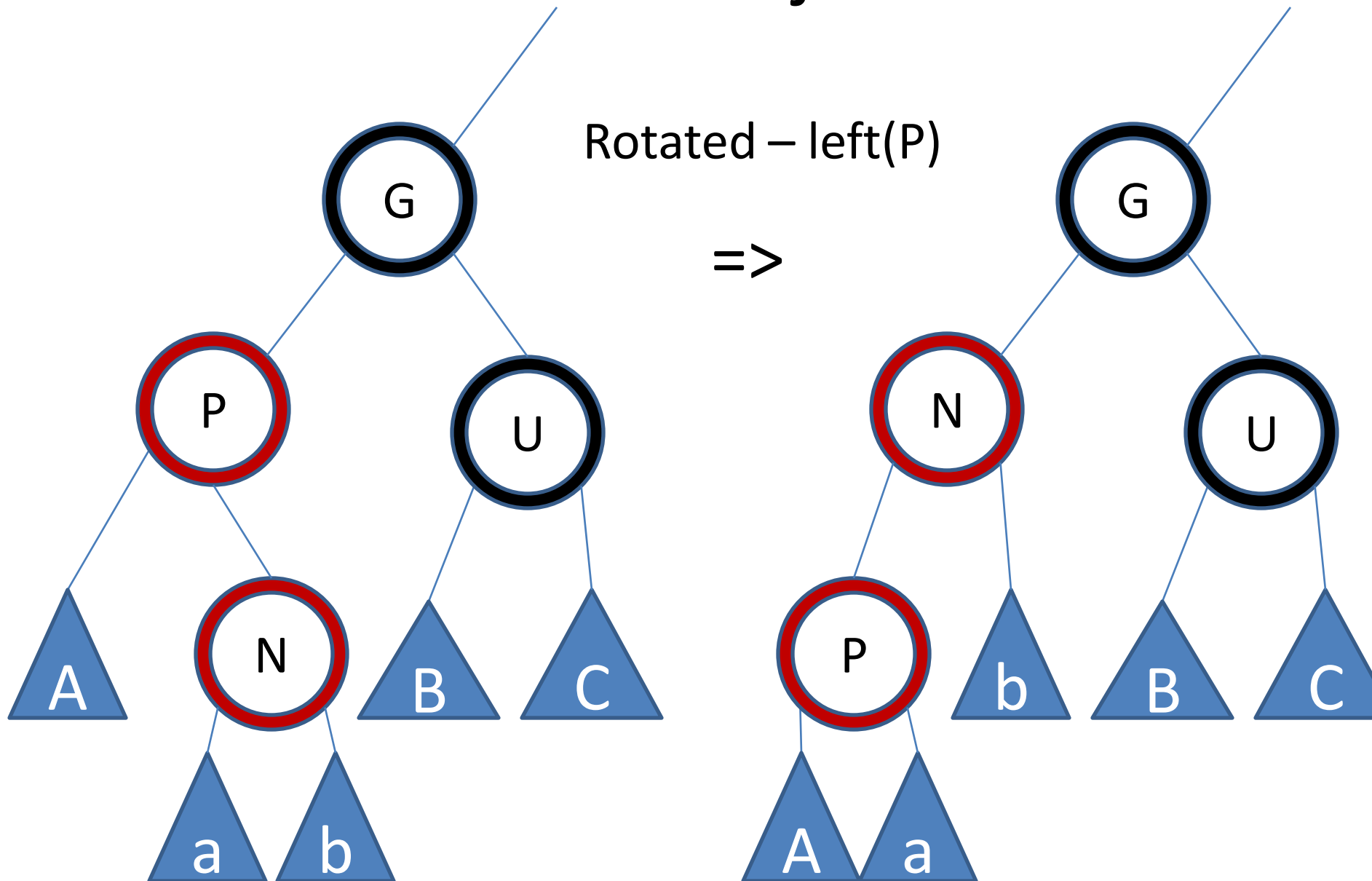


# First Case Fix



If parent of G is black, done. Otherwise consider G as the New node and recurse.

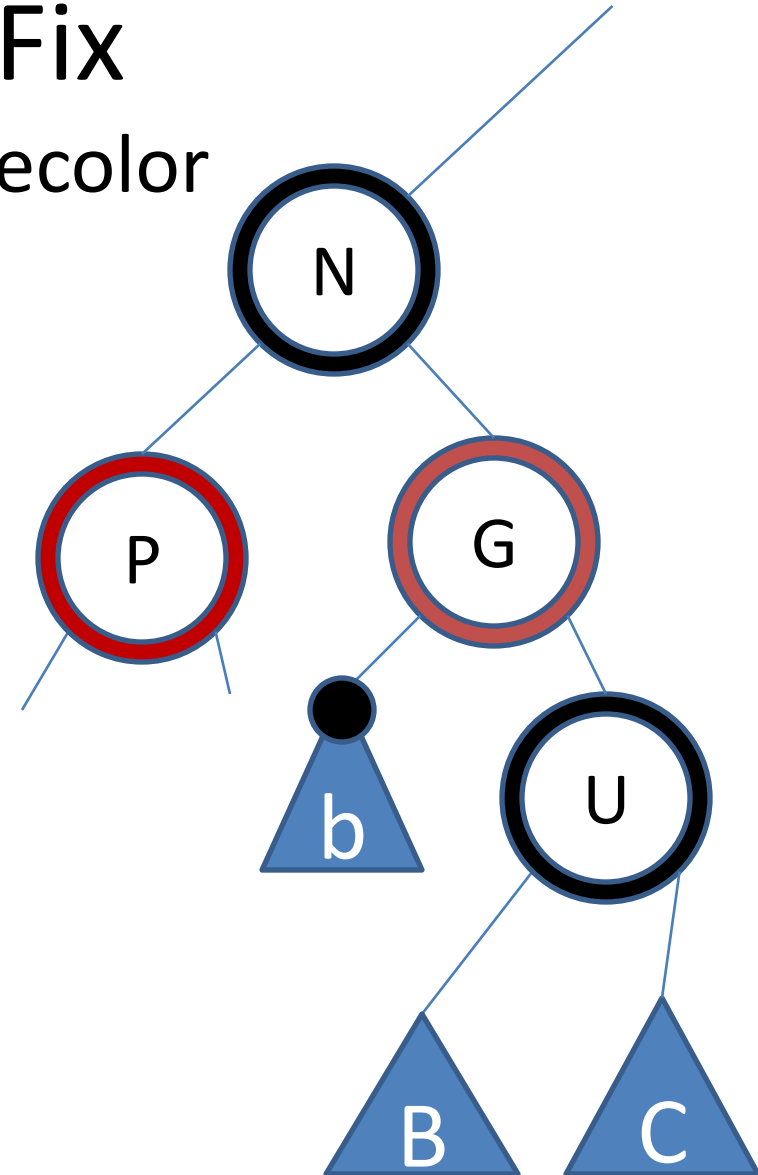
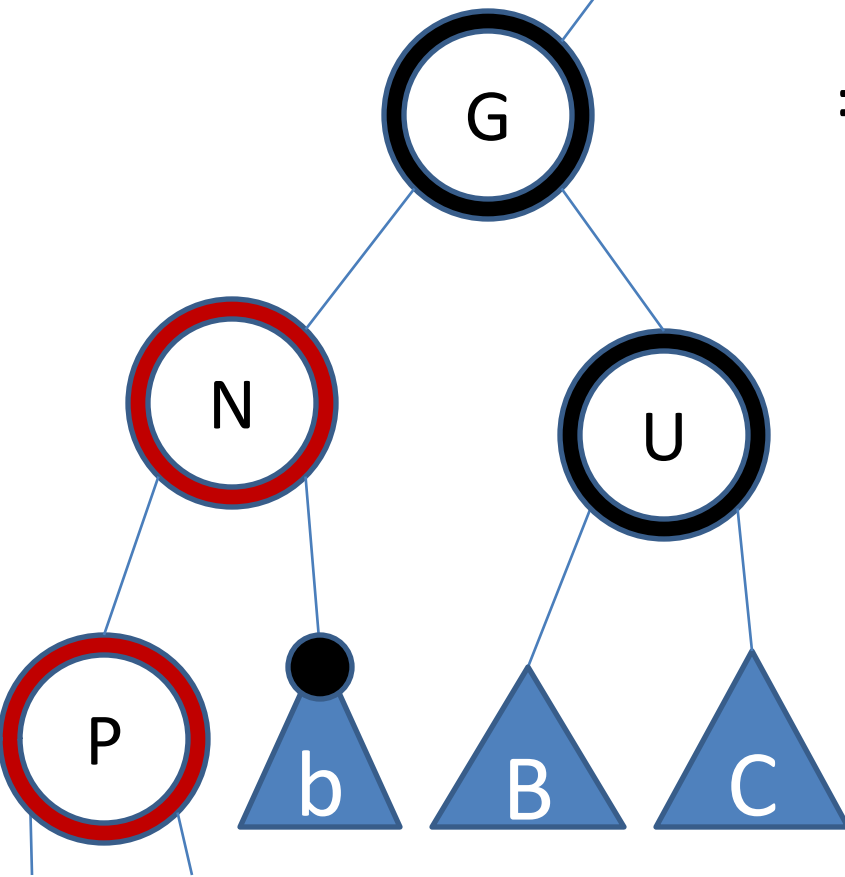
# Second Case Adjustment



# Third Case Fix

Rotate-right(G) and recolor

=>



# Case Progressions

First Case

First Case

Second Case

Third Case

Second Case

$\Rightarrow$

Third Case

Third Case

Done

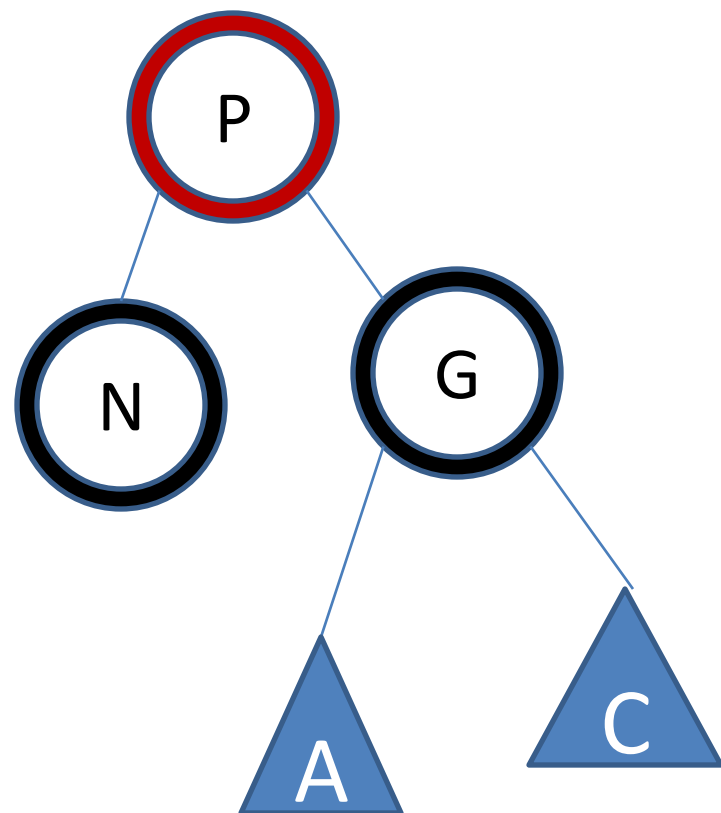
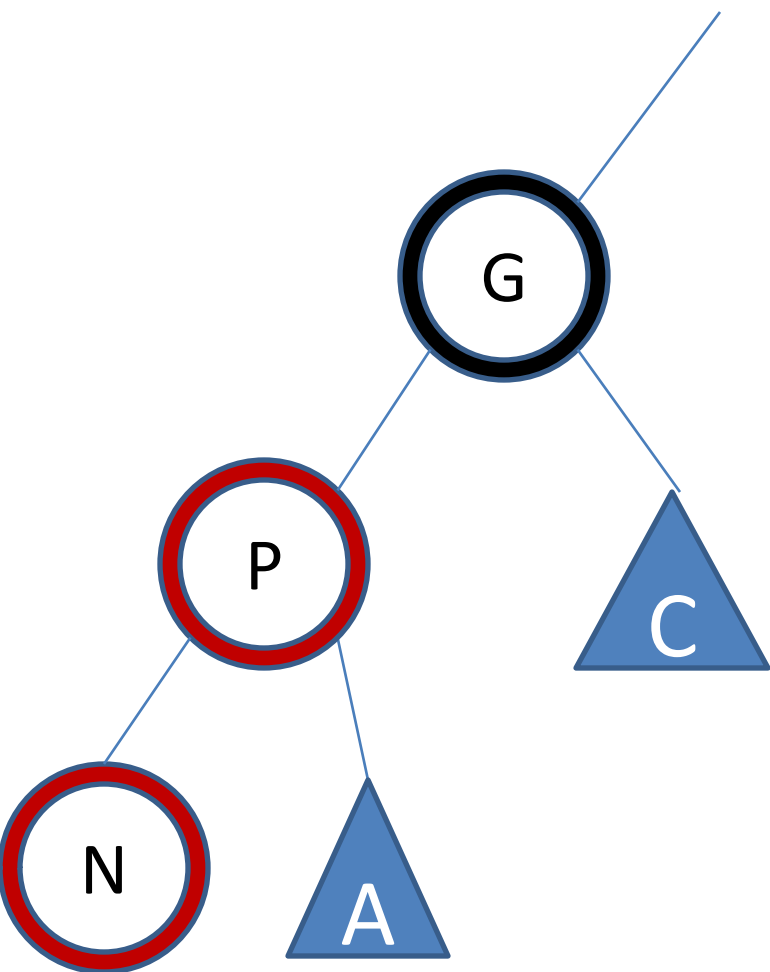


# Number of Operations

- First Case operations always move you up the tree and require a constant number of color flips,  $O(\log(n))$
- At most 2 rotations going through Second Adjustment and Third Case. Constant.
- $O(\log(n))$

# Alternative Method

- More rotations, but only 1 case, easier to code.
- <http://portal.acm.org/citation.cfm?id=968583>



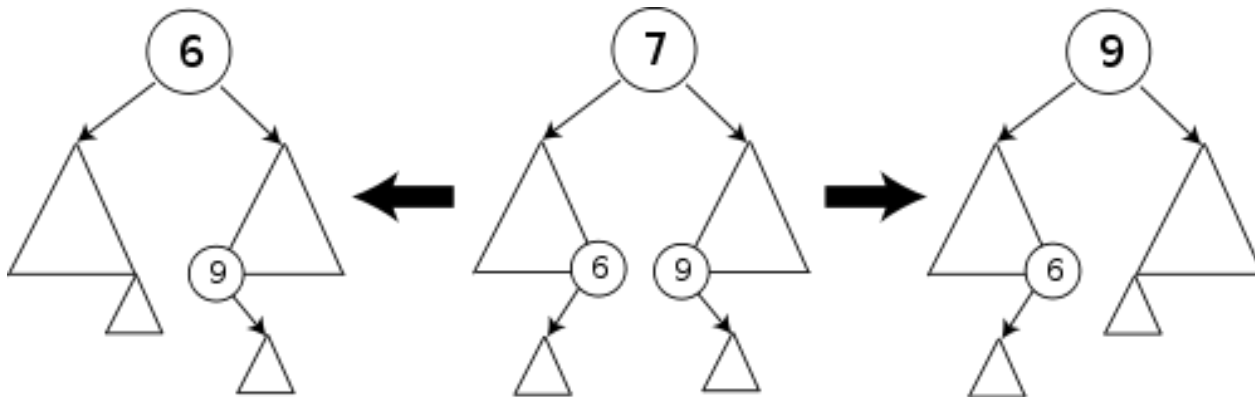
# Youtube Example

- <http://www.youtube.com/watch?v=vDHFF4wjWYU&feature=related>

# Deletion

First Step: find replacement value and copy. Use either max of left subtree or min of right subtree. Doing so maintains key order property.

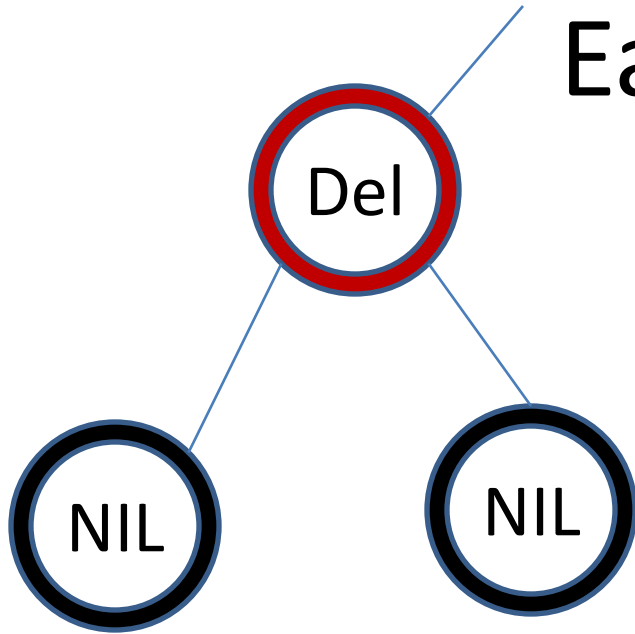
Second Step: delete the node of the copied value. Note, this node will have at most 1 child and can moved up in its place.



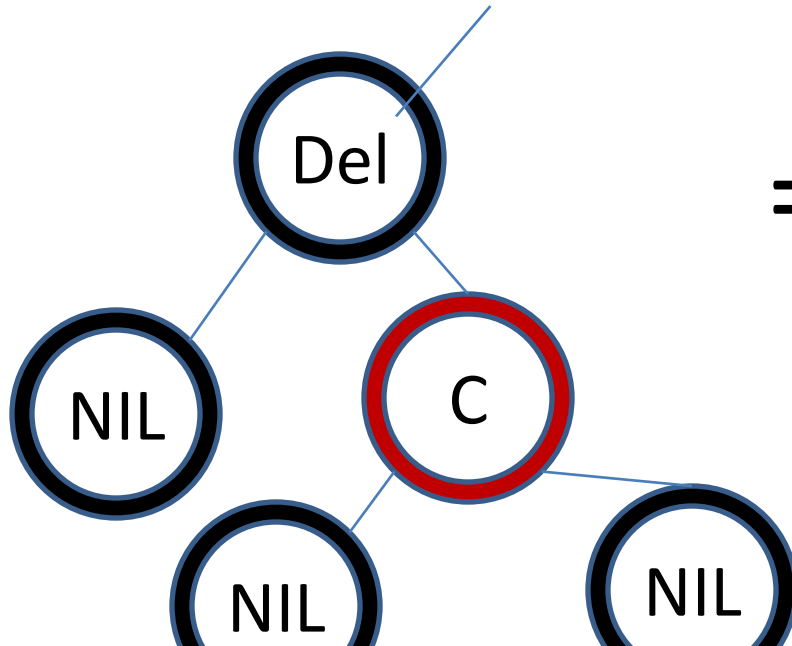
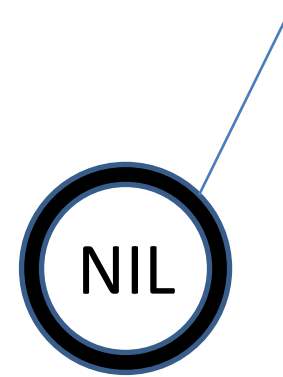
# Correcting

- Note, the value copy does not destroy the properties of the red-black tree. Only the deletion of 'foot' node may.
- If 'foot' node is red – no problem from deletion.
- Otherwise:

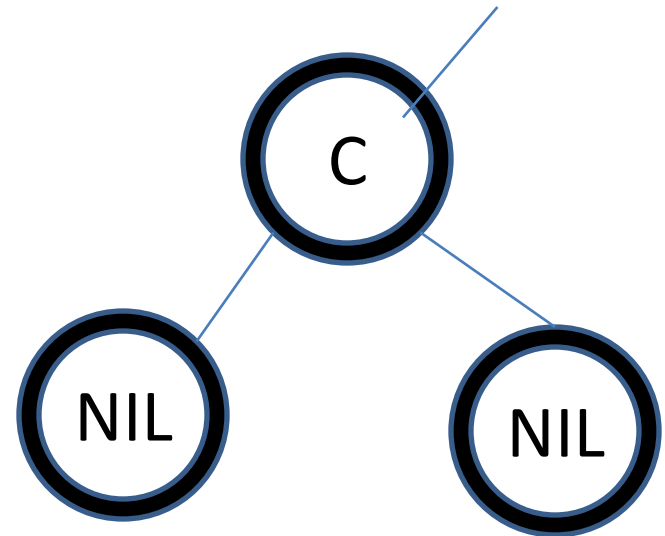
# Easy Cases



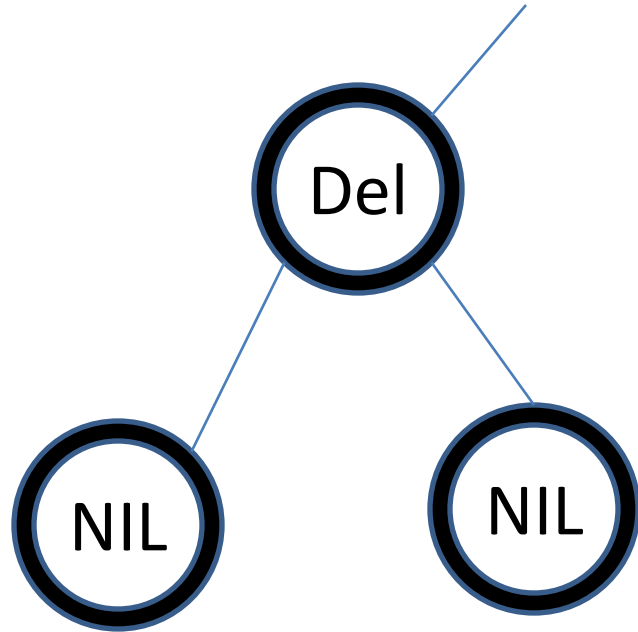
$\Rightarrow$



$\Rightarrow$



# Hard Case



Perform a series of rotations until we come across a red node and can convert it.

See CLRS or wikipedia for a list of all cases.



# Schemes

- Insertion algorithm destroys the coloring condition and seeks to rotate preserving black-height, pushing the bad coloring up the tree until it can recolor to satisfy.
- Deletion algorithm destroys the black-height and seeks to rotate preserving the coloring condition, pushing the unbalanced black height up the tree until it can color a red to black to satisfy.

