CS 2800: Discrete Structures, Fall 2014

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#### Not to be confused with...



#### Arithmeticorum Lib. II.

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QVÆSTIO VIII.

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quadratos. Ponatur rurfus priconstat latus diuidendi. Esto i-

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#### Fermat's Last Theorem:

 $x^n + y^n = z^n$  has no integer solution for n > 2

### Recap: Modular Arithmetic

- Definition:  $a \equiv b \pmod{m}$  if and only if  $m \mid a b$
- Consequences:
  - $-a \equiv b \pmod{m} \text{ iff } a \bmod m = b \bmod m$   $(\text{congruence} \Leftrightarrow \text{Same remainder})$
  - If  $a \equiv b \pmod{m}$  and  $c \equiv d \pmod{m}$ , then
    - $a + c \equiv b + d \pmod{m}$
    - $ac \equiv bd \pmod{m}$

(congruences can sometimes be treated like equations)

• If p is a prime number, and a is any integer, then

$$a^p \equiv a \pmod{p}$$

• If p is a prime number, and a is any integer, then

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• If a is not divisible by p, then

$$a^{p-1} \equiv 1 \pmod{p}$$

#### Examples:

$$-21^7 \equiv 21 \pmod{7}$$
... but  $21^6 \not\equiv 1 \pmod{7}$ 

$$-111^{12} \equiv 1 \pmod{13}$$

$$-123,456,789^{2^{57,885,161}-2} \equiv 1 \pmod{2^{57,885,161}-1}$$

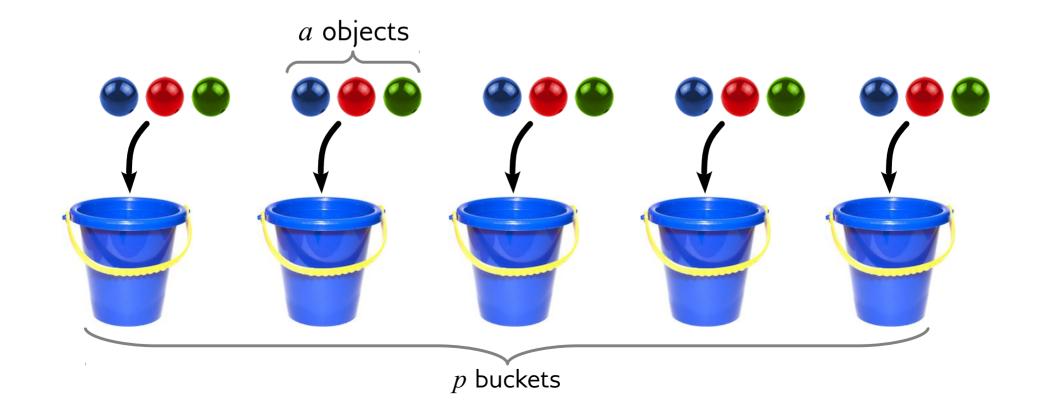
### Two proofs

- Combinatorial
  - ... counting things
- Algebraic
  - ... induction
- We'll consider only non-negative a
  - ... the result for non-negative  $\alpha$  can be extended to negative integers

(try it using what we know of congruences!)

# Counting necklaces

- Due to Solomon W. Golomb, 1956
- Basic idea:  $a^p$  suggests we see how to fill p buckets, where each is filled with one of a objects

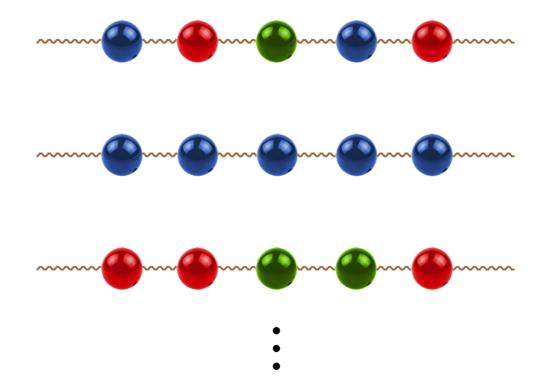


# Strings of beads

- Each way of filling the buckets gives a different sequence of p objects ("beads")
  - a<sup>p</sup> such sequences

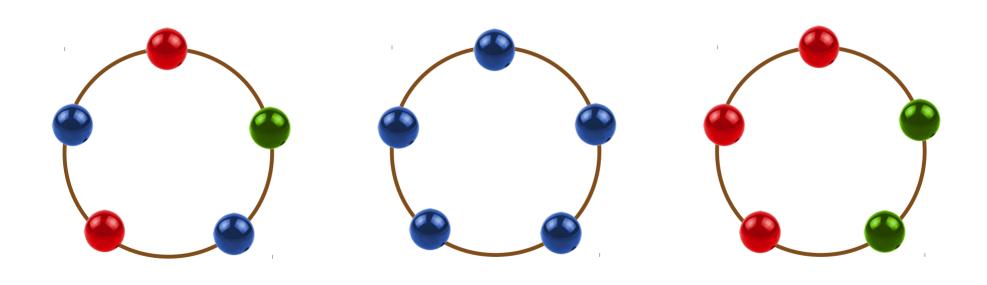
# Strings of beads

Now string the beads together...



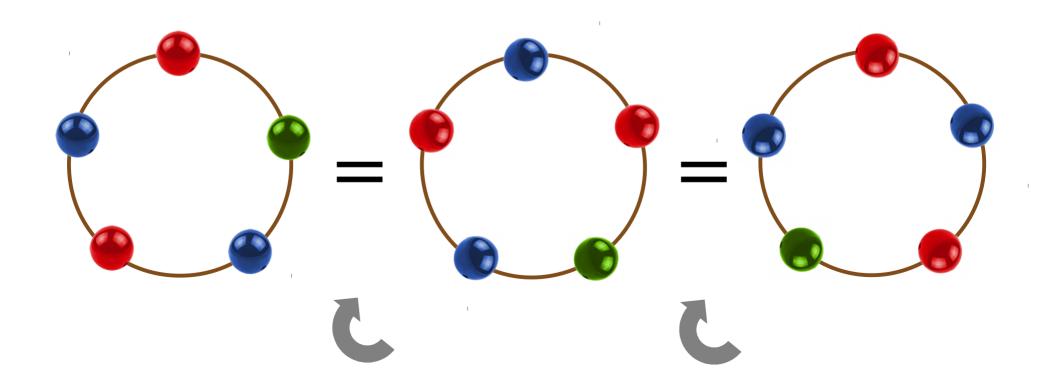
# Strings of beads

• ... and join the ends to form "necklaces"



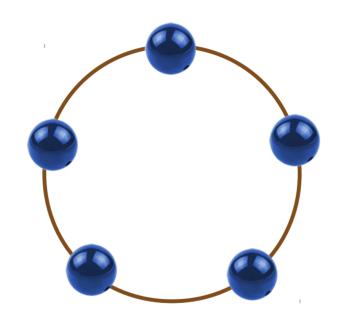
#### A necklace rotated...

- ... is the same necklace
  - Different strings can produce the same necklace when the ends are joined



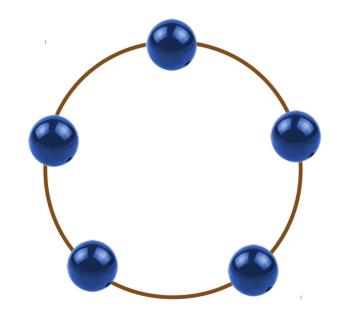
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Containing beads of a single color



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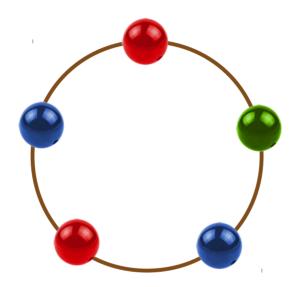


Only one possible string

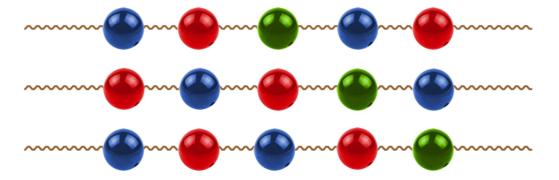


# Two types of necklaces

Containing beads of different colors

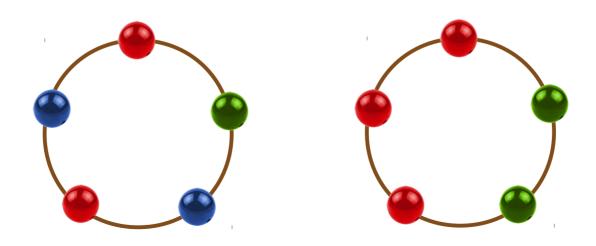


Many possible strings

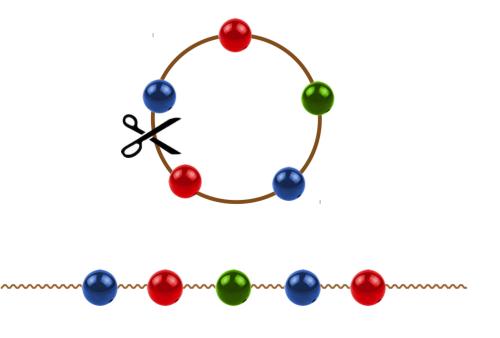


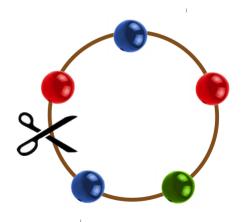
#### Lemma

- If p is a prime number and N is a necklace with at least two colors, every rotation of N corresponds to a different string
  - ... i.e. there are exactly p different strings that form the same necklace N

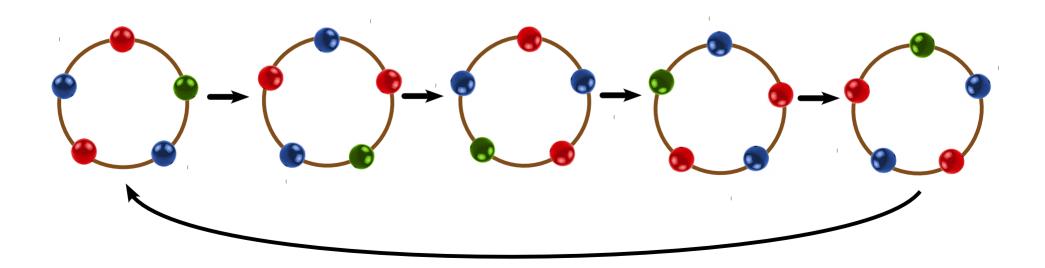


- First, note that each string corresponds to
  - a rotation of the necklace, and then...
  - ... cutting it at a fixed point





- No more than p strings can give the same necklace
  - There are only p (say clockwise) rotations of the necklace (that align the beads) before we loop back to the original orientation



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  - ... which is a contradiction, unless r = 0

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  - k = 1 (impossible if necklace has at least two colors) or
  - -k=p
- This proves the lemma

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    - $\Rightarrow a^p \equiv a \pmod{p}$  QED!

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- By the Binomial Theorem,

$$(a+1)^p = a^p + {p \choose 1}a^{p-1} + {p \choose 2}a^{p-2} + {p \choose 3}a^{p-3} + \dots + {p \choose p-1}a+1$$

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- By the Binomial Theorem, integer. P is prime, so it isn't canceled

Binomial coefficient  $\binom{P}{k}$  is

P!/k!(P-k)!, which is always an

out by terms in the denominator

$$(a+1)^p = a^p + \binom{p}{1}a^{p-1} + \binom{p}{2}a^{p-2} + \binom{p}{3}a^{p-3} + \dots + \binom{p}{p-1}a+1$$

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Hence proved by induction