

PROOF INDUCTION ON n .

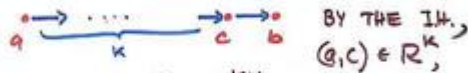
$n=1$ - BY DEFINITION

ASSUME THE THEOREM HOLDS FOR $n=k$ (I.H.)

THERE IS A PATH FROM a TO b OF LENGTH $k+1$

IFF FOR SOME c SUCH THAT $(c, b) \in R$

THERE IS A PATH FROM a TO c OF LENGTH k



THEREFORE $(a, b) \in R^{k+1}$

CONNECTIVITY RELATION $R^* = \{(a, b) \mid \text{THERE IS A PATH FROM } a \text{ TO } b \text{ IN } R\}$

$$R^* = \bigcup_{n=1}^{\infty} R^n$$

EXAMPLE. $R = \{(a, b) \mid \text{COMPUTER } a \text{ HAD AT LEAST ONE CONNECTION WITH A NET COMPUTER } b\}$

$$R^* = WWW$$

|| Th. THE TRANSITIVE CLOSURE OF R IS R^* .

PROOF. $TR.CL.(R) \subseteq R^*$ SINCE R^* IS TRANSITIVE

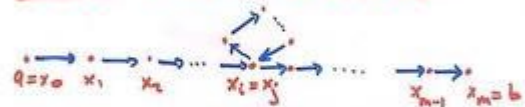
$R^* \subseteq TR.CL.(R)$ SINCE A CONNECTED PAIR (a, b) BELONGS TO ANY TRANSITIVE $S \supseteq R$

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Th. $R^* = R \cup R^2 \cup \dots \cup R^n$, WHERE $n = |A|$
 R IS A RELATION ON A .

PROOF. IMMEDIATELY FOLLOWS FROM AN OBSERVATION THAT IF a AND b FROM A ARE CONNECTED IN R , THEN THERE IS A PATH FROM a TO b WITH LENGTH NOT EXCEEDING n .

INDEED, BY THE PIGEONHOLE PRINCIPLE LONGER PATHS HAVE EQUAL VERTICES, THEREFORE CYCLES. CYCLES CAN BE DELETED!



$$\text{COR. } M_{R^*} = M_R \vee M_R^2 \vee \dots \vee M_R^n$$

$$\text{EXAMPLE. } M_R = \begin{bmatrix} 1 & 0 & 1 \\ 0 & 1 & 0 \\ 1 & 1 & 0 \end{bmatrix}; M_R^2 = \begin{bmatrix} 1 & 1 & 1 \\ 0 & 1 & 0 \\ 1 & 1 & 1 \end{bmatrix};$$

$$M_R^3 = M_R^2 \quad M_{R^*} = M_R \vee M_R^2 \vee M_R^3 = \begin{bmatrix} 1 & 1 & 1 \\ 0 & 1 & 0 \\ 1 & 1 & 1 \end{bmatrix}$$

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