

LEC.15 09/29/99

## RECURSIVE ALGORITHMS

BY THE EUCLIDEAN ALGORITHM

$$\gcd(a, b) = \gcd(b, r_2) \quad \text{CIRCULAR PROCEDURE}$$

THIS GUARANTEES  
THE CONVERGENCE

NO!  
A SMALLER INPUT  
 $\min(a, b) \geq \min(b, r_2)$

def AN ALGORITHM IS CALLED RECURSIVE IF  
IT WORKS BY REDUCING TO ITS OWN VALUE  
ON SMALLER INPUTS (USUALLY PERFORMED BACKWARD)

EXAMPLE  $\gcd(a, b) = \gcd(b, r_2)$  WHERE  $a \geq b > 0$   
 $a = b \cdot q_1 + r_2 \quad 0 \leq r_2 < b$   

$$\begin{cases} \gcd(a, b) = \gcd(b, a \bmod b) \\ \gcd(d, 0) = d \text{ when } d > 0 \end{cases}$$

SPECIFIES A SEQUENCE OF REDUCTIONS THAT ENDS  
BY  $\gcd(d, 0) = d$

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def ITERATIVE PROCEDURE P:

$$\begin{cases} \text{INPUT } a \\ P_n(a) = P(P_{n-1}(a)) \end{cases} \quad \text{USUALLY PERFORMED STRAIGHT}$$

$a, P(a), P(P(a)), \dots, P(P(\dots(P(a) \dots))$   
 n TIMES

RECURSIVE

PROCEDURE factorial(n)  
IF n=1 THEN  
factorial(n) := 1  
ELSE  
factorial(n) :=  
n \* factorial(n-1)

ITERATIVE

PROCEDURE iterfactorial(n)  
x := 1  
FOR i:=1 TO n  
x := i \* x  
{x is n!}

PROCEDURE fibonacci(n)  
IF n=0 THEN fibonacci(n) := 0  
ELSE IF n=1 THEN  
fibonacci(n) := 1  
ELSE fibonacci(n) :=  
fibonacci(n-1) + fibonacci(n-2)

PROCEDURE iterfibonacci(n)  
IF n=0 THEN y := 0  
ELSE  
BEGIN x := 0 y := 1  
FOR i:=1 TO n-1  
BEGIN z := x+y, x := y, y := z  
END  
END {y is nth Fibonacci num}

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## RECURSIVE COMPUTING OF $a^n$

$$\begin{cases} a^0 = 1 \\ a^n = a \cdot a^{n-1} \end{cases} \quad \text{THE SAME PROCEDURE, SMALLER INPUT}$$

PROCEDURE POWER(a: NONZERO REAL, n: > 0 INTEGER)  
IF n=0 THEN POWER(a, n) = 1  
ELSE POWER(a, n) := a \* POWER(a, n-1)

## RECURSIVE SEQUENTIAL SEARCH ALGORITHM: FIND x

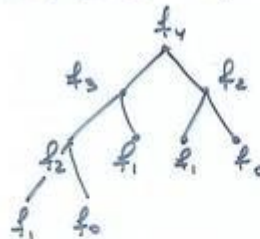
IN  $a_1, \dots, a_n$   
 PROCEDURE SEARCH(i, n, x)  
 IF  $a_i = x$  THEN  
     location := i  
 ELSE IF  $i = n$  THEN  
     location := 0  
 ELSE  
     SEARCH(i+1, n, x)

SEARCHES FOR x IN  $a_1, a_2, \dots, a_n$   
 DECREASING PARAMETER (n-i) IS ONE LESS AFTER EACH ITERATION  
 => CONVERGES

Th. IF  $f(n)$  IS DEFINED RECURSIVELY  

$$\begin{cases} f(0) = a \\ f(n) = \text{SOME FUNCTION } G \text{ of } f(0), \dots, f(n-1) \end{cases}$$
  
 AND G IS COMPUTABLE, THEN f IS COMPUTABLE.  
PROOF. COMPUTE  $f(0), f(1), f(2), \dots, f(n)$

## RECURSIVE FIBONACCI COMPUTATION



$f_{n+1}$  - 1 ADDITIONS TO  
FIND  $f_n$   
 (EXERCISE!)  
 EXPONENTIAL

## ITERATIVE FIBONACCI COMPUTATION

$$f_0, f_1, f_2 = f_1 + f_0, f_3 = f_2 + f_1, \dots, f_n = f_{n-1} + f_{n-2}$$

n-1 ADDITIONS / LINEAR

HW 3.4: 4, 14

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