

CS212

Java Practicum

Lecture 2

SaM

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Announcements

- <http://www.cs.cornell.edu/courses/cs212/>
- Part1 posted on website:
 - solo
 - hint: http://en.wikipedia.org/wiki/Bit_shift
- CMS sign-up (if you're not on already); e-mail patwell@cs.cornell.edu

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What is SaM? Why SaM?

- **From last lecture:**
 - computer stores data and instructions in memory
 - *fetch-and-decode cycle*:
 - JVM is _____ of computers
 - bytecodes are _____
- **SaM:**
 - a **simple stack machine** (with a **heap**)
 - see SaM on CS212 for full instruction set
 - gives us legible instruction set
 - you will write compiler to generate *Samcode*
 - BTW, what's a compiler? (last panel...)

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Samcode Instructions

- **Low-level instructions:**
 - push and pop values in memory
 - *mnemonics* for bit patterns
- **Structure:**
 - opcode*
 - opcode operand*
- **Example:**

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Using SaM

- **Areas:**
 - demo in lecture
 - download JAR file from SaM link on website
- **Simplest use:**
 - Create text file with Samcode
 - Open file
 - Run
- **Examples**

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Structure of Samcode File

- ASCII Text! (What's ASCII?)
- Write instructions on new lines
- *//* indicates single-line comments, which are ignored
- Program ends with _____
- Program must leave how many items on Stack?

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Focus on Stack

- **Call Stack** (and other names):
 - function calls function calls ...
 - when last function done, go back, then back, then ...
 - how to picture this structure?
- **Frame:**
 - each function's portion of Stack
 - variables, data, administrative info
- Cells and addresses
 - **start at 0!**
- Helpful picture?

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Useful Registers

- **Frame Based Register (FBR)**
 - administrative information
 - keeps track of current frame (and thus, function)
- **Stack Pointer (SP)**
 - use register
 - store location of next free cell in stack
- Helpful picture?

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Some Instructions

- ALU:
 - **ADD** (**SUB**, **TIMES**,)
 - **AND** (**OR**, **XOR**, **NOT**, ...) (**0** is false; all else is true)
 - **EQUAL** (**LESS**, ...)
 - Generally follows *below op top* (see documentation)
- Stack Manipulation:
 - **PUSHIMM** *c* (**PUSHIMMF** *f*, ...)
 - **DUP** (**SWAP**, ...)
 - **PUSHABS** *k*, **STOREABS** *k*
 - **PUSHOFF** *k*, **STOREOFF** *k*
- Register: **ADDSP** *n*
- Control: **STOP**
- Many others!
 - see on-line documentation
 - see **Chapter 1**

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Some Examples

- Notation:
 - Prefix: $(1 - 2) - 3$
 - Postfix: $1 2 - 3 -$
- Logical: $\sim(4 \leq 5)$
 - Samcode rem: *below op top*
- Samcode?

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Program Storage?

- Main memory model:
 - store programs as _____
 - so, instructions have patterns of _____
- Where are they in SaM?
 - Samcode read into an array
 - array stores instruction objects
- Want more? See documentation and source code
 - **SaM**→**Individual Files**→**Core**→**Instructions**
 - See next page for example
- How to load your own instructions?
 - recompile everything (a pain)
 - or...use SaM's *instruction loader*

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Example

```
package edu.cornell.cs.sam.core.instructions;
import edu.cornell.cs.sam.core.*;

public class SAM_ADD extends SamInstruction {
    public void exec() throws SystemException {
        int type1 = mem.getType(cpu.get(SP) - 2);
        int type2 = mem.getType(cpu.get(SP) - 1);
        mem.push(higherPrecedence(type1, type2), mem.pop() + mem.pop());
        cpu.inc(PC);
    }
}
```

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Variable Scope

- Example:
 - is the following legal?

```
int x(int x) { return x++; }
int y(int x) { return x(x); }
```
 - why? why not?
- Scope of variable:
 - region of code in which variable represents something
 - how does Java indicate?
- Local and global variables:
 - each function has its own local variables
 - global variables shared

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Variables and Frames

- A way to picture variables in frames...
 - variable gets cell
 - Aside: SaM shows type of cell
- Samcode program:
 - allocate cell
 - fill cell
 - later retrieve/change contents
 - finally deallocate cell (why?)

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Allocation and Deallocation

- Pushing:
 - **PUSHIMM**... (see SaM website)
- Allocating:
 - Allocate **v** amount of vars: **ADDSP v**
 - Deallocate **v** amount of vars: **ADDSP -v**
- Example:

```
ADDSP 3
ADDSP -1
ADDSP -1
ADDSP -1
STOP
// error mesg (why?)
```

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How to access a variable?

- Addressing of variables:
 - absolute
 - relative
- **Absolute**:
 - don't worry about your current frame
 - figure out variable address on stack
 - eg) globals
- **Relative**:
 - do worry about your current frame
 - figure out variable address with respect to FBR value
 - eg) locals

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Absolute Address

- **Instructions:**
 - To **store** a value **v** at location **i**:
 - **PUSHIMM v**: $\text{Stack}[\text{SP}] \leftarrow v$; $\text{SP}++$
 - **STOREABS i**: $\text{Stack}[\text{i}] \leftarrow \text{Stack}[\text{SP}-1]$; $\text{SP}--$
 - To **retrieve** a value **v** from location **k**:
 - **PUSHABS k**: $\text{Stack}[\text{SP}] \leftarrow \text{Stack}[\text{k}]$; $\text{SP}++$
- **Example:**

```
int rv;          ADDSP 3
int x;          PUSHIMM 10
int y;          STOREABS 1
x = 10;         PUSHIMM 20
y = 20;         STOREABS 2
rv = x + y;     PUSHABS 1
return rv;     PUSHABS 2
               ADD
               STOREABS 0
               ADDSP -2
               STOP
```

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Relative Address

- **Instructions:**
 - To store a value **v** at location **i**:
 - **PUSHIMM v**: $\text{Stack}[\text{SP}] \leftarrow v$; $\text{SP}++$
 - **STOREOFF i**: $\text{Stack}[\text{i}+\text{FBR}] \leftarrow \text{Stack}[\text{SP}-1]$; $\text{SP}--$
 - To retrieve a value **v** from location **k**:
 - **PUSHOFF k**: $\text{Stack}[\text{SP}] \leftarrow \text{Stack}[\text{k}+\text{FBR}]$; $\text{SP}++$
- **Picture?**

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Example

```
ADDSP 1 // rv of program
JSR add // new frame (jump to "add")
STOREOFF 0 // store rv of "add"
STOP // done

add: // code for "add" function
LINK // store old FBR (0) and set new FBR (2)
ADDSP 3 // allocate space for x, y, rv of add
// rv of add is at relative address 1
PUSHIMM 10 // push value 10
STOREOFF 2 // store 10 in x's cell
PUSHIMM 20 // push value 20
STOREOFF 3 // store 20 in y's cell

PUSHOFF 2 // retrieve x
PUSHOFF 3 // retrieve y
ADD // x+y
STOREOFF 1 // store x+y as rv of add
ADDSP -2 // deallocate x, y

SWAP // exchange rv of add for old FBR
UNLINK // restore old FBR (0)
SWAP // exchange rv of add for return address
RST // return to Samcode just after "JSR add"
```

```
public int add()
{
    int x, y;
    x = 10;
    y = 20;
    return x+y;
}
```

NOTE: We will use a different frame structure later!

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Human Compiling

- **Compiling:**
 - translate **code** (like Java) to machine **code** (like Samcode)
 - compiler (like **javac**) does the work for you
- **Human Compiling** (Part 1 of CS212):
 - you identify simple expressions and statements
 - you convert them into Samcode
 - you test your Samcode problems in SaM
 - we grade your correctness and style

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