

Week 9

Discussion of Assignment Part 3

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CS 212 – Spring 2004

Announcements

- No section tonight
 - Sections *will* be held next week (M & W)
- Do not alter the files we provide
 - Several groups altered the Part 2 files (e.g., some groups changed the packages)
- Make use of Office Hours!
- If your Part 2 did not compile or if it failed many tests
 - The graders are *not expected* to determine the exact nature of any problems with your code
 - If there is some small error, you can request a regrade
 - ◆ Describe the problem
 - ◆ Describe the fix
 - ◆ Provide working code

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Grammar for Bali (Part 3)

```
program -> function*
function -> functionHeader functionBody
functionHeader ->
  ( type | void ) name ( [ parameters ] )
functionBody ->
  { variableDeclaration* } { statement* }
type -> ( int | boolean ) [ [ ] ]
parameters -> type name ( , type name )*
variableDeclaration -> type name ( , name )* ;
```

- There must be a *main function*
- Functions can be *overloaded*
- Arrays are one-dimensional
- Valid types are *int*, *boolean*, *int array*, or *boolean array*

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More Grammar for Bali (Part 3)

```
statement -> return [ expression ] ;
statement -> { statement* }
statement -> if expression then statement
  [ else statement ]
statement -> while expression do statement
statement -> do statement while expression ;
statement -> expression ;
statement -> print expression ;
statement -> ;
statement -> target = expression ;
target -> name [ subscript ]
```

- The Part 3 sam-code for statements should be nearly the same as for Part 2
- The *expression statement* is executed for its side-effects (it might sort an array, for instance)
- To parse an *assignment statement*, pretend it's an *expression statement* until you reach the equal sign (=)

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Rest of Grammar for Bali (Part 3)

```
expression -> expPart [ binaryOp expPart ]
expPart -> unaryOp expPart
expPart -> literal
expPart -> ( expression )
expPart -> name [ functionArgs | subscript ]
functionArgs -> ( [ expressionList ] )
expressionList -> expression ( , expression )*
subscript -> [ expression ]
literal -> integer | true | false | null
binaryOp -> arithmeticOp | comparisonOp | booleanOp
arithmeticOp -> + | - | * | / | %
comparisonOp -> < | > | <= | >= | == | !=
booleanOp -> && | || | ^
unaryOp -> - | !
```

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The Major Tasks for Part 3

- The "hard stuff"
 - Implementing arrays
 - ◆ Use of the Heap
 - ◆ Use of **null**
 - Implementing functions
 - ◆ Stack frames
 - ◆ Overloading
 - Error handling
- Warning: finish this stuff before messing with any of the bonus work
- Bonus work
 - Multiple error reporting
 - Multidimensional arrays
 - ◆ Multiple subscripts
 - ◆ In declaration
 - ◆ In expression
 - ◆ In target for assignment statement
 - ◆ An additional kind of expression for array creation
 - Runtime error reporting

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Code Patterns for Arrays

Bali Code	Sam Code	Comment
myIntegers = int[4];	PUSHIMM 4 MALLOC	Push array's size onto Stack Create heap-block of size 5; push block's address onto Stack
	PUSHIMM 1 ADD STOREOFF 13	Address arithmetic We arbitrarily assume myIntegers is at offset 13 from the FBR
myIntegers[2] = 44;	PUSHOFF 13 PUSHIMM 2 ADD PUSHIMM 44 STOREIND	Push array's address onto Stack Subscript Address arithmetic Stores 44 into myIntegers[2]
x = myIntegers[2];	PUSHOFF 13 PUSHIMM 2 ADD PUSHIND STOREOFF 9	Push array's address onto Stack Subscript Address arithmetic Stored value (44) placed on Stack We arbitrarily assume x is at offset 9

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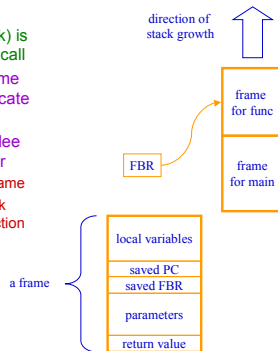
Use of `null` for Arrays

- Declaring an array
`int[] A;`
- Constructing an array
`A = int[6];`
- Initializing an array
`i = 0;`
`while i < 6 do {`
`A[i] = i;`
`i = i + 1;`
`}`
- When an array is declared but not yet constructed, the array variable has value `null`
- In the sam-code, an array variable (e.g., A) holds the address of the array
 - After array construction, this is an address in Heap
 - Before array construction, this should be the address 0 (an address clearly not within Heap)
- In other words, `null` in Bali-code corresponds to `0` in sam-code

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Recall: Stack Frames for Functions

- A new *frame* (on the stack) is created for each function call
 - We use the FBR (Frame Base Register) to indicate the current frame
 - The caller and the callee share responsibility for
 - ◊ creating the stack frame
 - ◊ cleaning up the stack frame when the function is done



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Recall: Signatures for Functions

- Functions in Bali can be *overloaded*
 - Functions can share same name as long as they differ in number or type of parameters
 - A function's *signature* determines which function to call
 - ◊ Signature encodes function's name as well as number and types of parameters
- Functions that share a name must *all* have same return type
- Bali does no automatic conversion of types
 - Thus function arguments and function parameters must match types exactly
- You can encode a function's signature in any way you want, but a Java String works fine

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Bonus: Multiple Error Reporting

- Error Handling
 - We will test your Part 3 compiler's response to errors in supplied Bali programs
 - Two kinds of errors
 - ◊ *Syntax errors*: code that violates the rules of the Bali grammar
 - ◊ *Semantic errors*: code that violates the rules of Bali semantics
- For bonus, use the `MultipleBaliException` class to accumulate and report multiple errors
- Which kind of error (*syntax error* or *semantic error*) is easier to deal with if we're trying to accumulate all errors?

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Bonus: Multidimensional Arrays

- Multiple subscripts
- An additional kind of expression for array creation

Changes in grammar:

```

type -> ( int | boolean ) ( [ ] )*
target -> name [ subscript* ]
expression -> [ [ expressionList ] ]
expPart -> name [ functionArgs | subscript* ]
  
```

Multiple brackets now allowed
Multiple subscripts now allowed
A new kind of expression
Multiple subscripts now allowed

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Bonus: Runtime Error Reporting

- Examples
 - An attempt to divide by zero
 - An array subscript out of bounds
 - Using an array before its construction
- In general, there is no program that can reliably detect such errors at compile time (see CS 381/481: *undecidable problems*)
- These errors can be detected at runtime, but...
 - You have to check for them and generate error messages using sam-code

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Recall: Expression Stmt vs. Assignment Stmt

- According to Bali's grammar
 - An expression statement and an assignment statement both start out looking like an expression
 - No way to tell that you are parsing an assignment statement until you get to the equal sign (=)
- Suggestion
 - Start parsing as if you are parsing an expression
 - Once the "expression" is complete, you check for the equal sign (=) to see if within an assignment statement
 - If in an assignment statement
 - ❖ You need to re-examine the AST you just built (for the expression) to see if it can be the target of an assignment statement
 - ❖ Your compiler should throw a `BaliSyntaxException` if the "expression" is inappropriate as a target

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