



# Lecture 23: Shortest Paths

CS 2110

April 16, 2026

# Today's Learning Outcomes

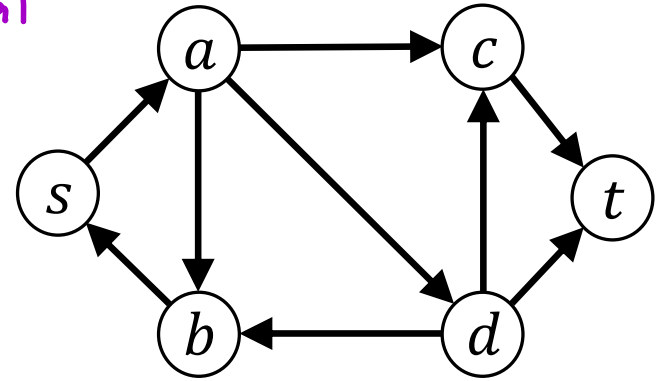
- 89. Identify substructures in graphs such as neighborhoods, paths, and cycles both visually and programmatically.
- 92. Visualize the state of a graph traversal at a given point of its execution, describing each node as either undiscovered, discovered, visited, and/or settled.
- 93. Describe Dijkstra's shortest path invariant and use it to establish properties on the unvisited portions of a graph at a given point in the execution of Dijkstra's algorithm.
- 94. Implement Dijkstra's shortest path algorithm and analyze its time/space complexities.

# Recall: Traversal Levels

The level of vertex  $v$ ,  $l(v)$ , in a traversal starting at  $s$  is the length of the shortest  $s \rightsquigarrow v$  path.

BFS visits vertices in level order.

Think in phases:





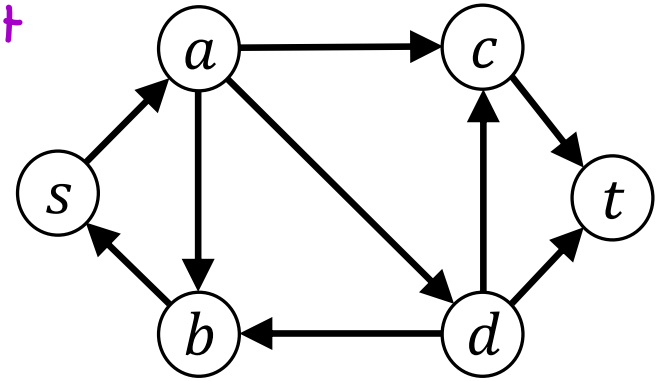
# Coding Demo: Level-Augmented BFS



# Shortest Paths in Unweighted Graphs

For a vertex  $v$  at level  $l$ , its shortest  $s \rightsquigarrow v$  path will include  $l$  edges

$$s \xrightarrow{1} u_1 \xrightarrow{2} u_2 \dots \xrightarrow{l-1} u_{l-1} \xrightarrow{l} v$$





# Coding Demo: BFS with PathInfo



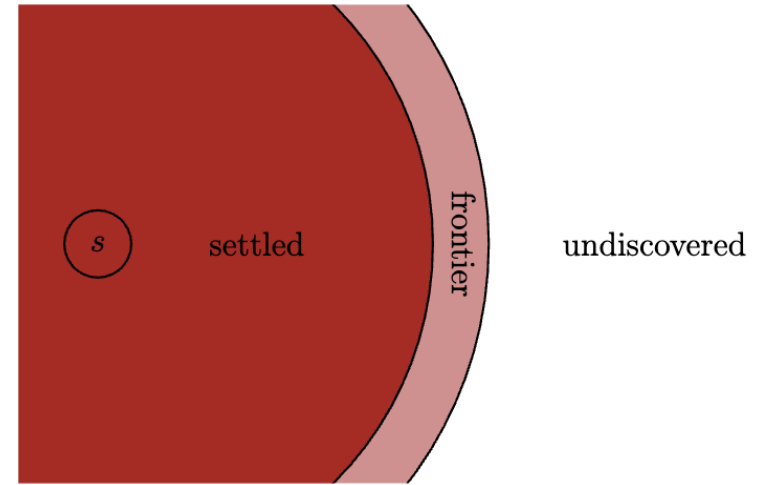


# Key Ideas of Unweighted Shortest Paths

1. At any point in the algorithm, our discovered map records the shortest (known) path distance to every node that we've discovered.

2. The next vertex to be removed from the frontier queue is always the *unvisited* vertex with the lowest level.

3. As soon as a vertex is visited/settled, we are guaranteed that we have located the shortest path to it from the source vertex, and this path contains only vertices that were previously settled.

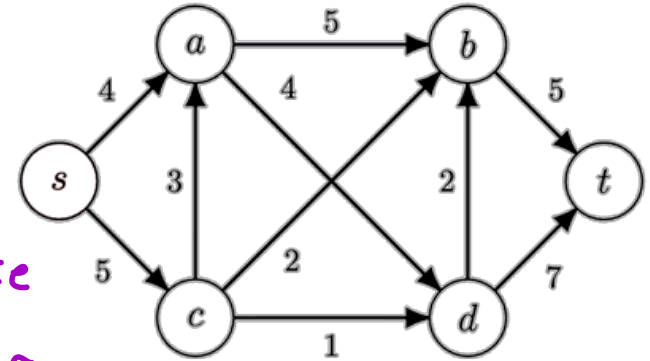


# Adapting to Weighted Graphs

(non-negative weights)

Some similar ideas to BFS:

- track discovered, frontier vertices
- visit vertices in distance order from source
- during visit, use outgoing edges to discover new vertices



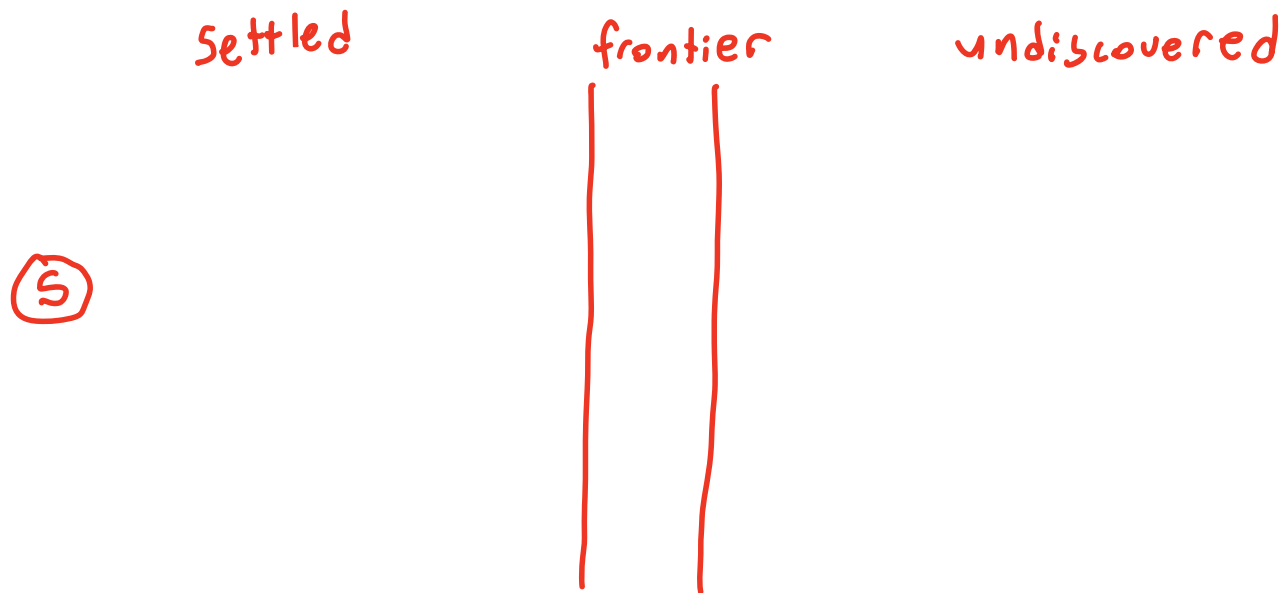
Some modifications needed:

# Dijkstra's Algorithm (High Level)

```
Initialize discovered  $\leftarrow$  {source} frontier  $\leftarrow$  {source}
while (frontier is not empty) {
  v  $\leftarrow$  frontier vertex with min known distance from s
  for each outgoing edge (v,w) {
    if w is undiscovered, discover it and add to frontier
      with best known distance  $d(s,w) = d(s,v) + \text{weight of edge (v,w)}$ 
    if w is discovered but we've found a shorter
      path to it, update its best known distance
  }
}
```

# Dijkstra's Invariant

At the start of each main loop iteration:



# Properties of Dijkstra's Algorithm

Vertices are visited /settled in increasing order of distance (just like BFS).

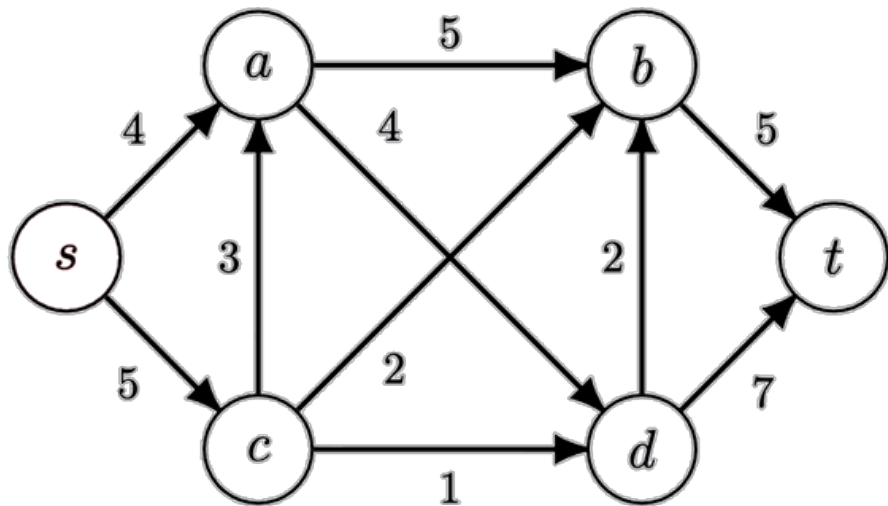
Best known distance to vertex only decreases as algorithm runs (when we find better path)



# Coding Demo: Dijkstra's Algorithm



# Dijkstra's Algorithm Walkthrough



	discovered					
vertex	<i>s</i>	<i>a</i>	<i>b</i>	<i>c</i>	<i>d</i>	<i>t</i>
distance						
prev						

frontier

→

→

# Dijkstra's Algorithm Complexity

```
while(!frontier.isEmpty()) {  
    V v = frontier.remove();  
    for (WeightedEdge<V> edge : v.outgoingEdges()) {  
        V neighbor = edge.head();  
        double dist = discovered.get(v.label()).distance() + edge.weight();  
        if (!discovered.containsKey(neighbor.label())) {  
            discovered.put(neighbor.label(), new PathInfo(dist, v.label()));  
            frontier.add(neighbor, dist);  
        } else if (discovered.get(neighbor.label()).distance > dist) {  
            discovered.put(neighbor.label(), new PathInfo(dist, v.label()));  
            frontier.updatePriority(neighbor, dist);  
        }  
    }  
}
```