

CS 2110: Priority Queues and Heaps

Prelim 2 scores released soon
median: 75

Priority Queue : extracts highest-priority element (minimum)

- shortest path (+ A*)
- simulation (next event)
- compression (Huffman coding)
- sorting

interface Priority Queue < E, P > {
 /** Add element to queue. */ // aka insert, add.
 void push (E element);
 /** Effect: remove highest-priority element */
 E pop(); // aka extract min, remove....
 /** Effect: increase priority of e.
 Requires: e is in queue. */
 void increasePriority (E elem, P priority);
}

← Java
Priority Queue
lacks this.

Implementation: Binary Heap

good asymptotic performance & simple
good constants

"binary heap" $\neq$ "memory heap"

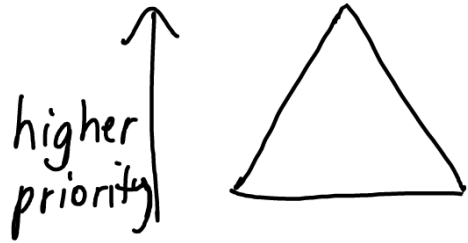
Heap: tree satisfying order invariant

Order

For all $n \neq \text{root}$,

$$n.\text{priority} \geq n.\text{parent.priority}$$

$\xrightarrow{\text{higher priority}}$



Binary heap: binary tree
also satisfying Shape invariant.

- All leaves at depth h or $h-1$
- If a leaf is at depth h :
 - all leaves to right are at depth $h-1$ or h
 - all possible depth- h leaves to left exist.

$$\Rightarrow h = O(\lg n)$$

Shape



Binary heaps?

Heap: tree satisfying order invariant Order
 For all $n \neq \text{root}$,
 $n.\text{priority} \geq n.\text{parent}.\text{priority}$

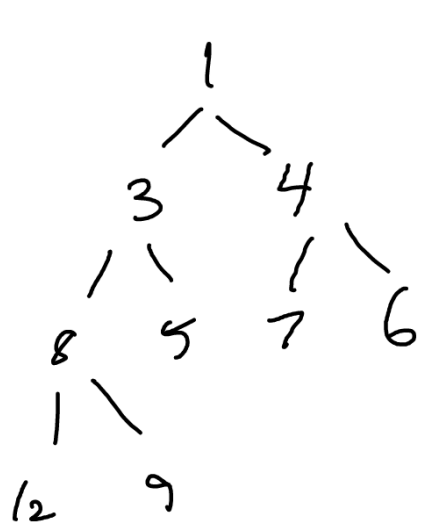
↑
higher priority



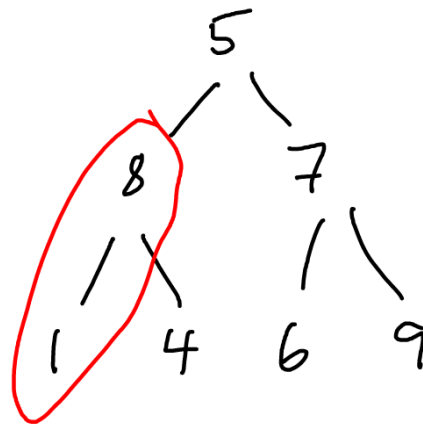
Binary heap: binary tree also satisfying Shape invariant.

Shape

- All leaves at depth h or $h-1$
 - All depth- h leaves are to the left of the $h-1$ leaves.
- $\Rightarrow h = O(\lg n)$

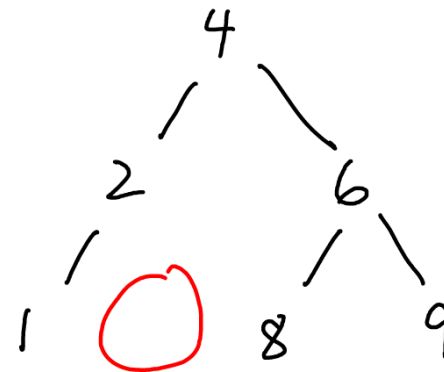


A



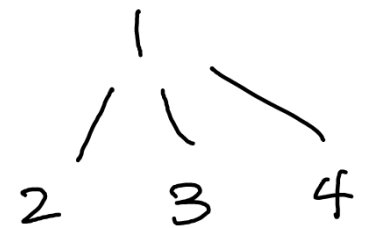
Order

B



Shape

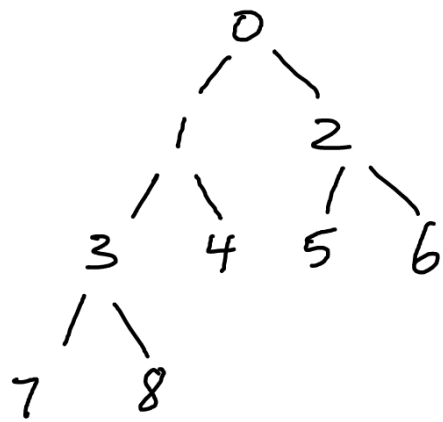
C



D

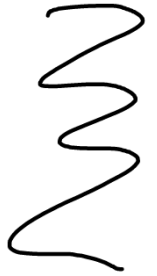
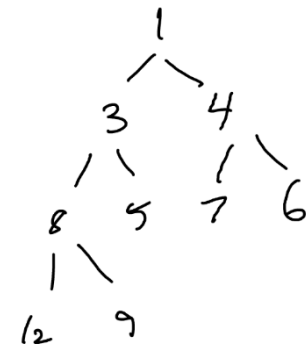
- not binary

Trick: store in resizable array. (no pointers)



array
indices

0	1	2	3	4	5	6	7	8
1	3	4	8	5	7	6	12	9



Shape: no gaps.

Tree w/o pointers.

Children of node at index i ?

left: $2i+1$
right: $2i+2$

Parent of node at index i ?

$$\left\lfloor \frac{i-1}{2} \right\rfloor = (i-1) \gg 1$$

push(e)

- Add at end of array to preserve Shape.
- "Bubble up" to restore Order.
 - swap with parent until parent has higher priority



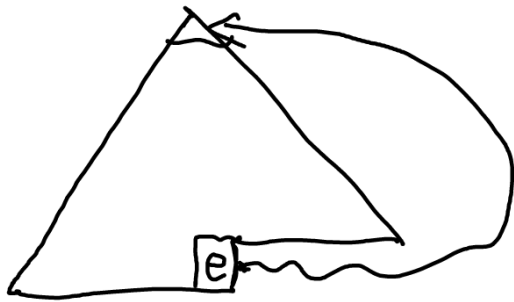
push(z)

Time: $O(h) = O(\lg n)$

pop(e) / Extract Min(e)

- Result = root value

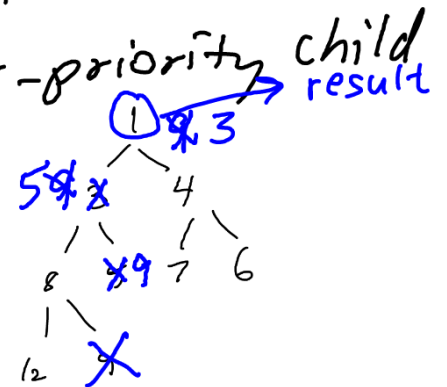
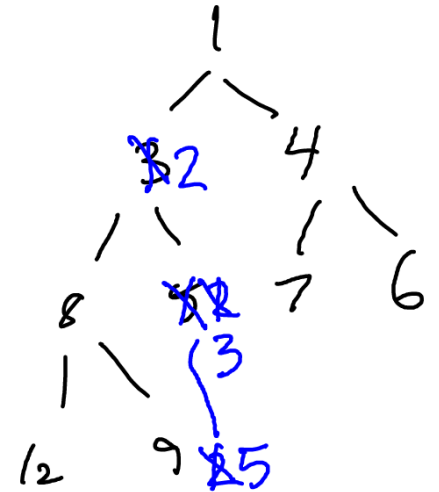
(remove)



fix
Shape

Time: $O(h) = O(\lg n)$

- Bubble down root
- Swap with higher-priority child (if necessary)



Increase Priority (e, p)

- Set new priority
- Bubble up to restore Order
- Time: $O(h) = O(\lg n)$

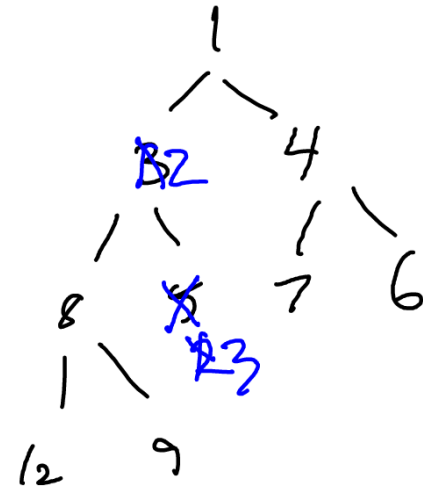
How to find element in the array?

Need location.

⇒ store index in element

or in hash table (element → index)

⇒ bubbling changes location



Heapsort

To sort an array:

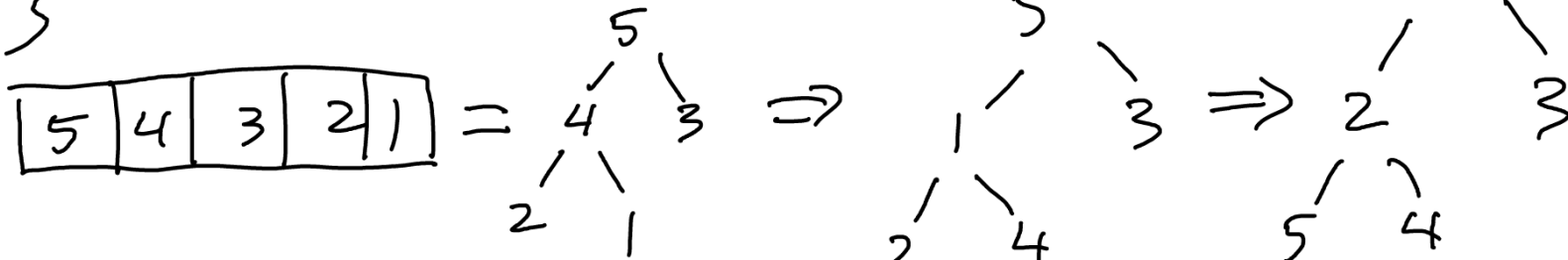
1. heapify
2. pop() until empty
— pops elements in priority order.

Heapify: $O(n)$

— $O(n \lg n)$

for $(i = \lfloor \frac{n}{2} \rfloor - 1; i \geq 0; i--) \{$
 bubbleDown(i);
}

$\lfloor \frac{n}{2} \rfloor \dots n-1$
are
leaves



loop invariant: subtrees of nodes $i+1, \dots, n-1$ are heaps
subtree of i is a heap exc. for i
bubble down $\Rightarrow$ fix subtree at i

