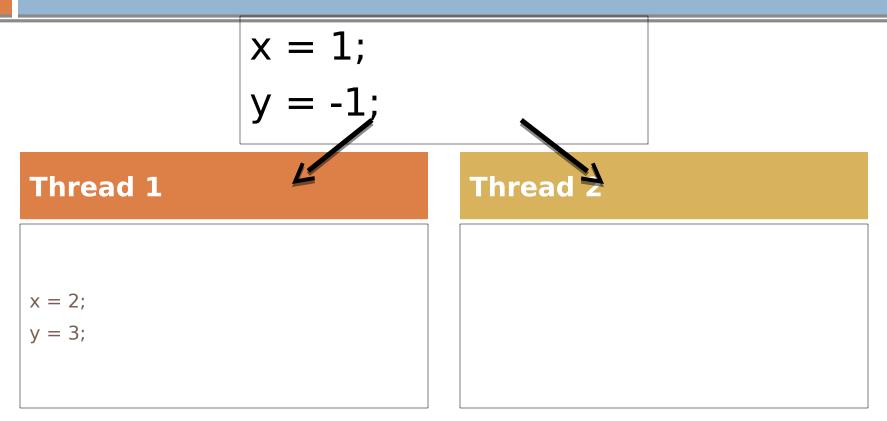
## **CONCURRENCY 3**

CS 2110 - Fall 2016

## Consistency



## What is printed?

0, 1, and 2 can be printed!

## Consistency

#### **Thread 1 on Core 1**

**Thread 2 on Core 2** 

Write 2 to x in local cache Write 3 to y in local cache 3 gets pushed to y in memory

2 gets pushed to x in memory

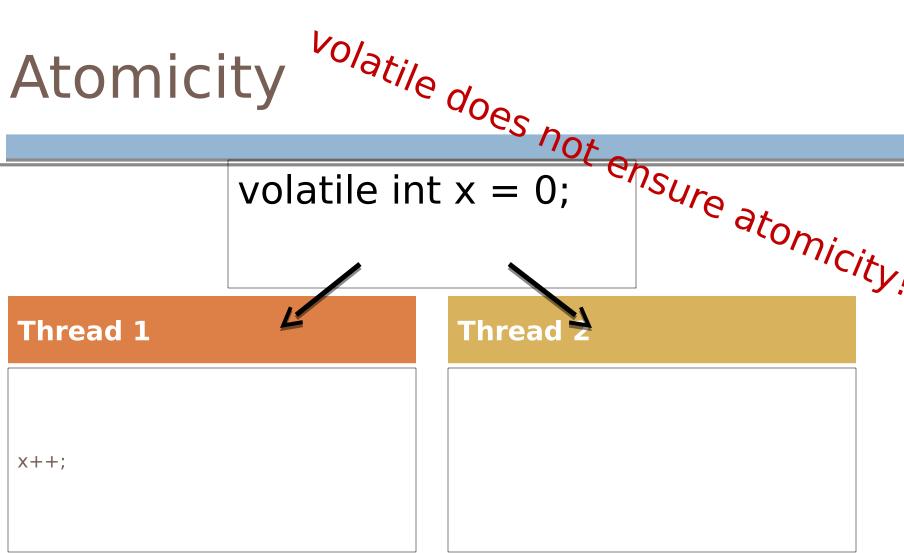
Not sequentially consistent!

## Harsh Reality

- Sequential Consistency
  - There is an interleaving of the parallel operations that explains the observations and events
  - Currently unknown how to implement efficiently
- Volatile keyword
  - Java fields can be declared volatile
  - Writing to a volatile variable ensures all local changes are made visible to other threads
  - x and v would have to be made volatile to

# Atomicity





What is the value of x?

Can be both 1 and

## java.util.concurrent.atomic

- class AtomicInteger,AtomicReference<T>, ...
  - Represents a value
- method set(newValue)
  - has the effect of writing to a volatile variable
- method get()
  - returns the current value
- effectively an extension of volatile
- but what about atomicity???

## Compare and Set (CAS)

- boolean compareAndSet(expectedValue, newValue)
  - If value doesn't equal expectedValue, return false
  - if equal, store newValue in value and return true
  - executes as a single atomic action!
  - supported by many processors
  - without requiring locks!

```
AtomicInteger n = new AtomicInteger(5);
n.compareAndSet(3, 6); // return false - no change
n.compareAndSet(5, 7); // returns true - now is 7
```

## Incrementing with CAS

```
/** Increment n by one. Other threads use
n too. */
public static void increment(AtomicInteger
n) {
  int i = n.get();
  while (n.compareAndSet(i, i+1))
     i = n.get();
```

// AtomicInteger has increment methods

### Lock-Free Data Structures

- Usable by many concurrent threads
- using only atomic actions no locks!
- compare and swap is god here
- but it only atomically updates one variable at a time!

Let's implement one!