

List Overview

- Arrays
 - Random access :)
 - Fixed size: cannot grow on demand after creation : (
- · Characteristics of some applications
 - Do not need random access
 - Require a data structure that can grow and shrink dynamically to accommodate different amounts of data

 Lists satisfy these requirements

- Common operations
 List creation
 - Accessing elements in a list
 - Inserting elements into a list
 - Deleting elements from a list

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List Operations

- An ADT (Abstract Data Type):
 - Specifies public functionality
 - Hides implementation detail from users
 - Allows us to improve/replace implementation
 - Forces us to think about fundamental operations (i.e., the interface) separately from the implementation
- ype): List Operations:
 - Create
 - Insert object
 - Delete object
 - Find object
 - Replace object
 - Size? Empty?
 - Usually sequential access (not random access)
 - A Java interface corresponds nicely to an ADT

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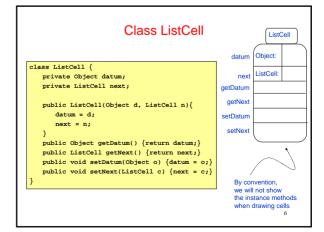
A Simple List Interface

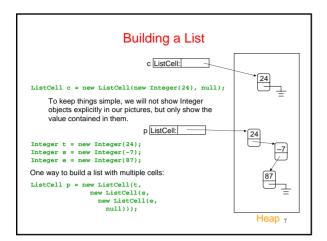
```
public interface List {
   public void insert(Object element);
   public void delete(Object element);
   public boolean contains(Object element);
   public int size();
}
```

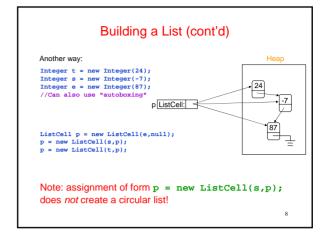
• Methods are specified, but no implementation

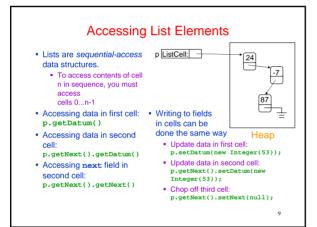
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List Data Structures · Can use an array · Can use a sequence of linked cells Need to specify array size · We'll focus on this kind of Inserts & Deletes require implementation moving elements Must copy array (to a larger We define a class ListCell from which we build lists array) when it gets full 24 87 24 -7 87 78 78 empty









```
Access Example: Linear Search

//Scan list looking for object x, return true if found
public static boolean search(Object x, ListCell c) {
    for (ListCell lc = c; lc != null; lc = lc.getNext()) {
        if (lc.getDatum().equals(x)) return true;
    }
    return false;
}

//Here is another version. Why does this work?
public static boolean search(Object x, ListCell c) {
    for (; c != null; c = c.getNext()) {
        if (c.getDatum().equals(x)) return true;
    }
    return false;
}

return false;
}
```

```
public static boolean search(Object x, ListCell c) {
   if (c == null) return false;
   if (c.getDatum().equals(x)) return true;
   return search(x, c.getNext());
}

public static boolean search(Object x, ListCell c) {
   return c != null &&
        (c.getDatum().equals(x) || search(x, c.getNext()));
}
```

Recursion on Lists

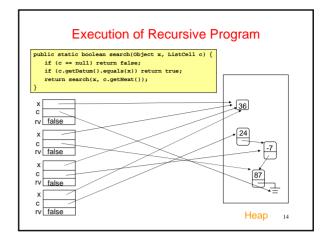
- Recursion can be done on lists
 - Similar to recursion on integers
- Almost always
 - Base case: empty list
 - Recursive case: Assume you can solve problem on (smaller) list obtained by eliminating first cell...
- Many list operations can be implemented very simply by using this idea
 - Although some operations are easier to implement using iteration

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Recursive Search

- · Base case: empty list
 - return false
- Recursive case: non-empty list
 - if data in first cell equals object x, return true
 - else return result of doing linear search on rest of list

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List with Header

public void setHead(ListCell c)

Iteration is Sometimes Better

- Given a list, create a new list with elements in reverse order
- Intuition: think of reversing a pile of coins

```
public static ListCell reverse(ListCell c) {
   ListCell rev = null;
   for (; c != null; c = c.getNext()) {
      rev = new ListCell(c.getDatum(), rev);
   }
   return rev;
}
```

• It is not obvious how to write this simply using a recursive style

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Heap

Variations on List with Header Header can also keep other info Reference to last -7 cell of list head Number of elements in list Search/insertion/ List deletion as head tail instance methods size 3 Heap 17

Special Cases to Worry About

- Empty list
 - add
 - find
 - delete
- Front of list
 - insert
- · End of list
 - find
 - delete

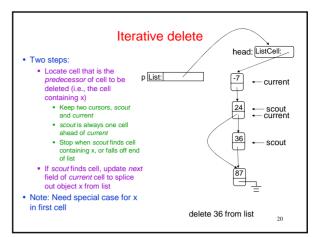
· Lists with just one element

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Example: Delete from a List

- Delete first occurrence of x from list c
 - Recursive delete
 - Iterative delete
- · Intuitive idea of recursive code
 - If list is empty, return null
 - If first element of c is x, return rest of list c
 - Otherwise, return list consisting of
 - · First element of c. and
 - List that results from deleting x from rest of list c

```
//recursive delete
public static ListCell delete(Object x, ListCell c) {
 if (c == null) return null;
 if (c.getDatum().equals(x)) return c.getNext();
 c.setNext(delete(x, c.getNext()));
 return c:
```



Iterative Code for Delete

```
public void delete (Object x) {
   if (head == null) return;
   if (head.getDatum().equals(x)) { //x in first cell?
       head = head.getNext();
       return:
   ListCell current = head;
   ListCell scout = head.getNext();
   while ((scout != null) && !scout.getDatum().equals(x)) {
       current = scout;
       scout = scout.getNext();
   if (scout != null) current.setNext(scout.getNext());
```

Doubly-Linked Lists

• In some applications, it is convenient to have a ListCell that has references to both its predecessor and its successor in the list.

```
class DLLCell {
  private Object datum;
   private DLLCell next;
   private DLLCell prev;
                                      45
                                next
```

Doubly linked vs. Singly linked

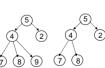
- · Advantages of doubly-linked lists over singly-linked lists
 - some things are easier e.g., reversing a doubly-linked list can be done simply by swapping the previous and next fields of each cell
 - don't need the scout to delete
- Disadvantages
 - doubly-linked lists require more heap space than singly-

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insert and delete take more time

Tree Overview • Tree: recursive data structure (similar to list)

- Each cell may have two or more successors (or children)
- Each cell has at most one predecessor (or parent) . Distinguished cell called root has
- no parent All cells reachable from root
- Binary tree: tree in which each cell can have at most two children: a left child and a right child



General tree

Binary tree





List-like tree

Not a tree