

## Finish Graphs

Lecture 25  
CS211 - Fall 2006

## Announcements

- Paul Chew's office hours for today (Tuesday, Nov 21) are cancelled
- No Section-meetings this week due to Thanksgiving Break

## Recall: Greedy Algorithms

- Dijkstra's Algorithm is an example of a **Greedy Algorithm**
- The Greedy Strategy is an algorithm design technique
  - Like Divide & Conquer
- The greedy algorithms are used to solve optimization problems
  - The goal is to find the *best* solution
- Works when the problem has the **greedy-choice property**
  - A global optimum can be reached by making locally optimum choices
- Problem: Given an amount of money, find the smallest number of coins to make that amount
- Solution: Use a Greedy Algorithm
  - Give as many large coins as you can
- This greedy strategy produces the optimum number of coins for the US coin system
- Different money system  $\Rightarrow$  greedy strategy may fail
  - Example: suppose the US introduces a 4¢ coin

## Minimum Spanning Trees

### Definition

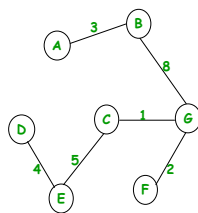
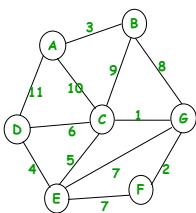
A *spanning tree* of an undirected graph  $G$  is a *tree* whose nodes are the vertices of  $G$  and whose edges are a subset of the edges of  $G$

### Definition

A *Minimum Spanning Tree (MST)* for a weighted graph  $G$  is the spanning tree of least cost (sum of edge-weights)

- Alternately, an MST can be defined as the least-cost set of edges so that all the vertices are connected
  - This has to be a tree... Why?
- A greedy strategy works for this problem
  - Add vertices one at a time
  - Always add the one that is closest to the current tree
  - This is called *Prim's Algorithm*

## An Example Graph and Its MST



## Prim's Algorithm

- $s$  is the start vertex
- $c(i,j)$  is the cost from  $i$  to  $j$
- Initially, vertices are unmarked
- $\text{dist}[v]$  is length of smallest tree-to- $v$  edge
- Initially,  $\text{dist}[v] = \infty$  for all  $v$

```

prim(s):
  dist[s] = 0;
  while (some vertices are unmarked){
    v = unmarked vertex with
      smallest dist;
    Mark v;
    for (each w adj to v){
      dist[w] = min[ dist[w], c(v,w) ];
    }
  }
  
```

- Runtime analysis
  - $O(n^2)$  for adj matrix
    - While-loop is executed  $n$  times
    - $O(n)$  time for for-loop
  - $O(m + n \log n)$  for adj list
    - **Total** time in for-loop is  $O(m)$
    - Use a PQ to find smallest dist
      - Regular PQ produces time  $O(m \log m)$
      - Can improve to  $O(m + n \log n)$  by using fancier heap

## Similar Code Structures

```
while (some vertices are unmarked) {
  v = best of unmarked vertices;
  Mark v;
  for (each w adj to v)
    Update w;
}
```

- **bfsDistance**
  - best: next in queue
  - update:  $\text{dist}[w] = \text{dist}[v] + 1$
- **dijkstra**
  - best: next in PQ
  - update:  $\text{dist}[w] = \min [\text{dist}[w], \text{dist}[v] + \text{cost}(v, w)]$
- **prim**
  - best: next in PQ
  - update:  $\text{dist}[w] = \min [\text{dist}[w], \text{cost}(v, w)]$

## Remembering Your Choices

- How can you remember which choices were made?

- Whenever  $\text{dist}[w]$  is updated we can remember the current  $v$  by using  $\text{parent}[w] = v$ ;

- Can use the parent info to construct the *bfs tree*, the *shortest path tree*, or the *minimum spanning tree*

```
while (some vertices are unmarked) {
  v = best of unmarked vertices;
  Mark v;
  for (each w adj to v)
    Update w;
    if (w changed) parent[w] = v;
}
```

## Two MST Algorithms (Both Greedy)

### Kruskal's Algorithm

- Choose the shortest edge  $e$  such that
  - $e$  is not yet processed
  - $e$  does not make a cycle

### Prim's Algorithm

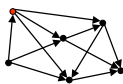
- Choose the shortest edge  $e$  such that
  - $e$  touches the tree
  - $e$  touches a vertex not in the tree

## Depth-First Search

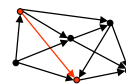
- Follow edges depth-first starting from an arbitrary vertex  $s$ , using a *Stack* to remember where you came from
- When you encounter a vertex previously visited, or there are no outgoing edges, retreat and try another path
- Eventually visit all vertices reachable from  $s$
- If there are still unvisited vertices, repeat

Easy to see this takes  $O(m)$  time

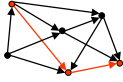
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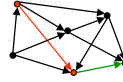
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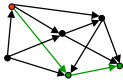
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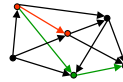
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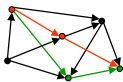
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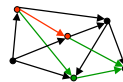
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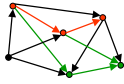
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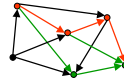
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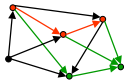
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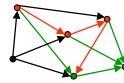
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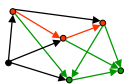
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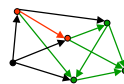
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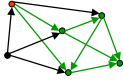
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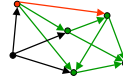
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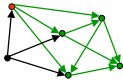
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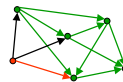
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### Depth-First Search



## DFS Notes

- Same as BFS, except we use a Stack instead of a Queue to determine which edge to explore next
- Can also implement DFS recursively
  - The Stack is represented *implicitly* in the Stack Frames created by the recursive calls

## Topological Sort using Recursive DFS

- Recall topological sort: find an ordering for the nodes of a dag so that all edges are "forward" edges
- Can use recursive DFS to do topological sort
  - Call DFS\_visit on each *unmarked vertex*

DFS\_visit (Vertex v):  
 Mark v as active  
 for each successor w of v:  
   if w not yet marked:  
     DFS\_visit(w)  
   if w is active: Exception  
 Mark v as done  
 Add v to *head* of topSort list



## Graph Overview

- Graph Definitions
  - Directed graph (digraph)
  - Undirected graph
  - Directed acyclic graph (dag)
  - Paths & cycles
- Graph Properties
  - Graph coloring
  - Planarity
  - Bipartite graphs
- Graph Implementations
  - Adjacency matrix
  - Adjacency lists
- Graph Searching
  - Breadth First Search (BFS)
  - Depth First Search (DFS)
- Other Graph Algorithms
  - Topological Sort
    - Using in-degree-0 nodes
    - Using DFS
  - Single-source shortest paths
    - Dijkstra's Algorithm
  - Minimum spanning tree (MST)
    - Prim's Algorithm
    - Kruskal's Algorithm