

Object-Oriented Programming

Object-Oriented Programming

- What do we mean by *object-oriented*?
- Why use it?
 - modularity (implementation hiding)
 - code reuse
 - type safety
 - inheritance (next time)
- Implementation
 - heap allocation of objects
 - references to objects

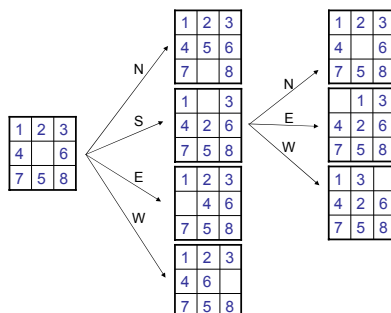
Some Context

- Programming “in the large”
 - big applications require many programmers
- General approach
 - break problem into smaller subproblems
 - assign responsibility for each subproblem to somebody
 - keep the interfaces small!
- Each subproblem must have a *specification*
 - *Functionality*: What services must code provide?
 - *Interface*: What input conditions does the code expect? What output conditions does it guarantee?
- Job of the programmer: provide an *implementation* (code) that meets the specification

The Message

- Separate the *specification* from the *implementation*
 - called *data abstraction* in the literature
 - more modular, easier to maintain
 - implementation is hidden from the client, can be changed without changing the interface
 - the client’s code does not break
- Object-oriented languages
 - encourage data abstraction
 - more modular code

The 8-Puzzle



Program Organization

- **class Puzzle**
 - an implementation of the game, written by you
 - functionality:
 - *init* – put puzzle in the initial state
 - *move* – move a tile N, S, E, or W to get a new state
 - *tile* – report which tile is in a given position
- **class TestPuzzle**
 - a client class, written by someone else
 - will communicate with **Puzzle** (your code) to play the game

Implementation

- Two subtasks
 - How do we represent a state (puzzle configuration)?
 - Given the representation, how do we implement `init`, `move`, and `tile`?
- Suppose no objects...

Representation of State

1	2	3
4	9	6
7	5	8

→ 123496758

- Model puzzle state as an integer between 123456789 and 987654321
 - 9 represents the empty square
- To convert integer `s` into a grid representation:
 - Remainder when `s` is divided by 10: tile in bottom right position
 - Java expression: `s % 10`
 - Quotient after dividing by 10 gives encoding of remaining tiles
 - Java expression: `s / 10`
 - Repeat remainder/quotient operations to extract remaining tiles
- This encoding may seem strange, but it arises many places in CS
 - Storing multidimensional arrays in memory

Implementing Operations

- `init`: put into initial configuration

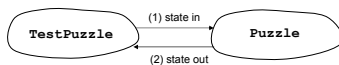

```
s = 123456879;
```
- `tile`: what tile is in position (`row, col`)?


```
return s / ((int)Math.pow(10, 8 - (3*row+col))) % 10;
```
- `move`: left to the reader

A Key Question

- Where do we keep the state?
 - method parameter/local variable
 - client keeps track of it
 - passed to `Puzzle` methods on each call
 - allocated on stack
 - class variable of `Puzzle` class
 - client does not see it
 - allocated in static area
- These implementation choices affect the interface of the `Puzzle` class

Interface L(ocal)



- State is implemented as local variable in class `TestPuzzle`
 - passed to/returned from methods in `Puzzle` class
- Interface of `Puzzle` class:

```
//return encoding of initial state
int init();
//return number of tile at grid (r,c)
int tile(int s, int r, int c);
//move to a new state, return new encoding
int move(int s, char d);
```

Implementation using L

```
class TestPuzzle {
    public static void main(String[] args) {
        int s = Puzzle.init();
        display(state);
        state = Puzzle.move(state, 'N');
        ...
    }

    public static void display(int s) {
        for (int r = 0; r < 3; r++) {
            for (int c = 0; c < 3; c++) {
                System.out.print(Puzzle.tile(state, r, c)
                    + " ");
            }
            System.out.println(); //newline after row
        }
    }
}

class Puzzle {
    public static int init() {
        return 123456879;
    }

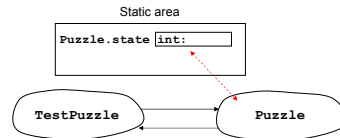
    public static int tile(int s, int r, int c) {
        return s / ((int)Math.pow(10, 8 - (3*r+c))) % 10;
    }

    public static int move(int s, char d) {
        ...
    }
}
```

Critique of Interface L

- No data abstraction!
 - **Puzzle** class implementer chose to implement state as an **int**
 - This representation is exposed in the interface, so the client code is aware of it
 - Client's code may depend on this encoding
 - If **Puzzle** class implementer decides to change the implementation (say, to represent state as a **long**), **client code breaks**

Interface S(tatic)



- State is implemented as class variable in class **Puzzle**
 - state does not have to be passed back and forth
 - representation is hidden from client
- Interface of **Puzzle** class:


```
void init(); //initialize the state
int tile(int r, int c); //return tile in position (r,c)
void move(char d); //move in direction d
```

Implementation using S

```
class TestPuzzle {
    public static void main(String[] args) {
        Puzzle.init();
        display();
        Puzzle.move('N');
        ...
    }

    public static void display() {
        for (int r = 0; r < 3; r++) {
            for (int c = 0; c < 3; c++) {
                System.out.print(Puzzle.tile(r,c)
                    + " ");
            }
            System.out.println(); //newline
        }
    }
}

class Puzzle {
    private static int state;

    public static void init {
        state = 123456879;
    }

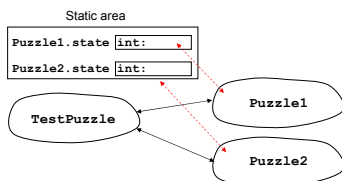
    public static int tile(int r, int c) {
        return state / ((int)Math.pow(10, 8 - (3*r+c))) % 10;
    }

    public static void move(char d) {
        ...
    }
}
```

Critique of Interface S

- Data abstraction: yes!
 - **Puzzle** class implementer chose to implement state as **int**
 - State representation is not visible outside of **Puzzle** class
 - If **Puzzle** class implementer decides to change implementation of state to **long**, client code does not have to change
- Problem: only one client and one puzzle at a time
 - state is a private class variable in class **Puzzle**
 - Mechanism we have used (class variable) gives right of puzzle creation to implementer of class rather than the client of the class

A Sneaky Solution



- Make copies of **Puzzle** class and rename them
- If client wants n puzzles, make n copies

Sneaky Implementation of S

```
class TestPuzzle {
    public static void main(String[] args) {
        Puzzle1.init();
        display1();
        Puzzle1.move('N');
        ...
        Puzzle2.init();
        display2();
        Puzzle2.move('N');
        ...
    }

    public static void display1() {
        ...
    }

    public static void display2() {
        ...
    }
}

class Puzzle1 {
    private static int state;

    public static void init {
        state = 123456879;
    }
    ...
}

class Puzzle2 {
    private static int state;

    public static void init {
        state = 123456879;
    }
    ...
}
```

Critique

- Data abstraction: yes
- Creation on demand: yes, but at cost of duplication of code
- Must know number of instances at compile time
- Naming issues

The Case for Objects

- Copying and renaming gives us
 - a unique name for each instance of the puzzle
 - a separate variable (**state**) to store the state of each instance
 - allows multiple simultaneous instances of the puzzle
- But all the instances are identical!
- Can we design language mechanisms to support the creation of separate instances?

Solution: Ask Gutenberg!

- Algorithm for making a copy of a book in the middle ages:
 - Hire a monk
 - Give monk paper and quill
 - Ask monk to copy text of book
- Algorithm for making n copies of a book
 - Hire a monk
 - Give monk lots of paper and quills
 - Ask monk to copy text of book n times
- Modern algorithm (Gutenberg, Strasbourg ca.1450 AD):
 - First make a template using movable type
 - Stamp out as many copies of book as needed
- Copying class code is like medieval approach to copying books!
- How do we exploit Gutenberg's insight in our context?
 - What is the template for puzzles?
 - How do we stamp out new puzzle instances from the template?
 - How do we name different puzzle instances?

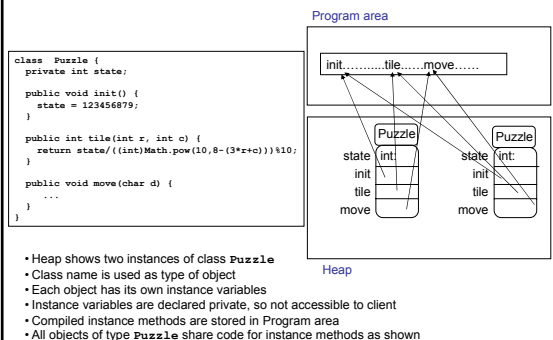


Gutenberg Bible
– The Huntington Collection

Object-Oriented Languages

- The **class definition** is the template
- Instances of the class are called **objects**
- Objects are stamped out (created) in an area of memory called the **heap**
- **instance variables**: when different instances are stamped out, they will each have their own copies of all instance variables (e.g. **state**)
- **instance methods**: code is shared among all instances of the same class, but references to instance variables in the code access those belonging to the correct object!
- **constructor**: a special method associated with a class invoked to create new instances of that class

Heap Allocation

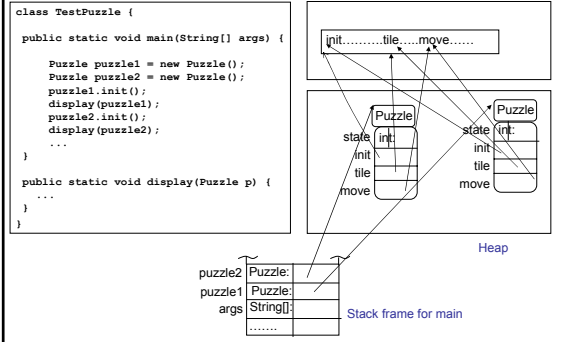


Naming Instances

- **Reference:** a variable that is a name for objects of some class
 - contains either a pointer to some object or `null`
- **Type of reference = class name**
`Puzzle p1; //declare a reference variable`
- Creation of an object using a constructor and assignment to a reference:

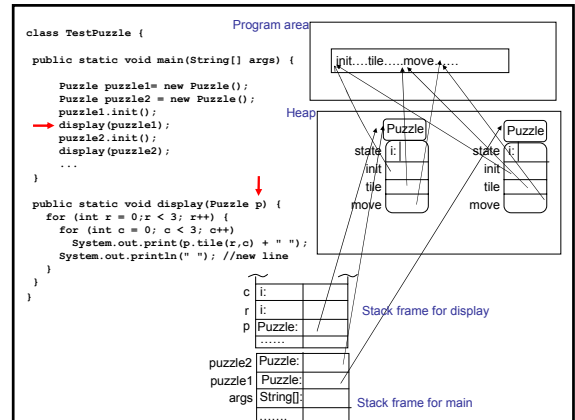

```
p1 = new Puzzle(); //create a new object, call it p1
Puzzle p2 = new Puzzle(); //can do both at once
```
- **Invoking instance method**
`p1.init();`
- **Implementation:**
 - examine object pointed to by `p1`
 - look inside object for starting address of method named `init`
 - invoke that method

Client Code

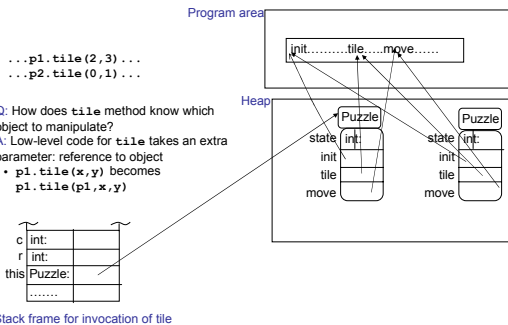


Method Invocation

- **References can be passed as parameters**
 - formal parameter becomes name for object in callee
 - callee can manipulate object using that name
 - on method return, caller sees any changes made to object by callee
- **Example: display method**
 - no need to have different code for each puzzle instance

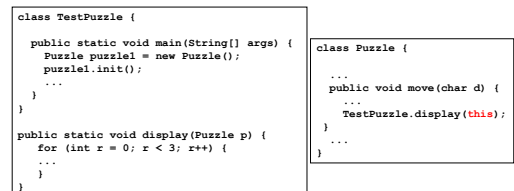


Accessing Instance Variables



Keyword this

- In instance method, **this** is a reference to object in which the method exists



Critique

- Data abstraction: yes
- Creation on demand: yes
- Duplicate class code: no
- Duplicate client code: no

Garbage Collection

- Intuitively, an object is *live* at time t if that object is still in use and can be accessed by the program after time t
- Formally (recursive definition), an object O is *live* if:
 - The runtime stack contains a reference to O
 - There is a live object O' that contains a reference to O
- Everything else is *garbage*
- Periodically, system detects garbage and reclaims it
- Start with the stack, trace all references, mark all objects seen – anything not marked is garbage
- **C, C++:**
 - Pointer arithmetic makes it hard to determine what is a reference
 - Storage reclamation must be done explicitly by programmer (`malloc`, `free`)
 - Highly error-prone

Conclusion

- Object-oriented languages support data abstraction and code reuse
- Objects (instances of a class) can be created on demand by client without breaking abstraction
- Client can hold a reference to an object, but implementation is hidden from it
- User-defined types: class names are used as types of objects and references