

Announcements

- **Extra office hours** today (instead of DIS sections); Zoom links on Canvas
- **P6 due tonight** at 11pm
- **Test 2B** feedback and grade estimation on website
- **Final exam:** Mon, 5/18, 9am. “2.5 hr” take-home, 48 hr submission window
- Optional **review session:** Sunday, 5/17, 2pm, Zoom (see Canvas)
- Please fill out **course evaluation**, *worth one BONUS point*, which can be used against any point lost on the final exam (150 points).
- Regular office/consulting hours end today. **Study period hours** are posted on Canvas and course website.

- Previous Lecture (and exercise):
 - Algorithms for sorting and searching
 - Insertion Sort
 - (Read about Bubble Sort in *Insight*)
 - Linear Search
 - Binary Search
 - Efficiency (complexity) analysis: analyze loops, count number of operations, use timing functions
 - Time efficiency vs. memory efficiency
- Today, Lecture 26:
 - Another “divide and conquer” strategy: **Merge Sort**
 - Review recursion
 - Semester wrap-up

Binary search is efficient, but we need to sort the vector in the first place so that we can use binary search

- Many different algorithms out there...
- We saw insertion sort (and read about bubble sort)
- Let's look at **merge sort**
- Another example of the “divide and conquer” approach (like binary search) but using recursion

Which task fundamentally requires less work: **sort a length 1000 array**, or **merge*** two length 500 sorted arrays into one?

A. Sort

B. Merge

C. The same

*Merge two sorted arrays so that the resultant array is sorted (not concatenate two arrays)

Comparison counting

How many comparisons (between elements) are required to run *insertion sort* on the following vector?

[9, 13, 24, 96, 12, 18, 56]

A. 6

C. 12

B. 7

D. 21

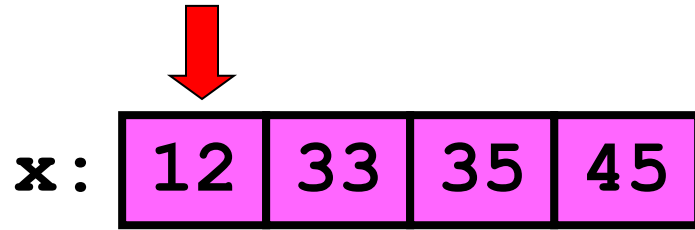
The central sub-problem is the **merging** of two sorted arrays into one single sorted array

12	33	35	45
----	----	----	----

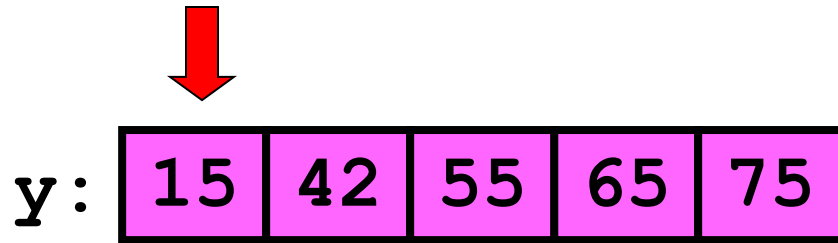
15	42	55	65	75
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12	15	33	35	42	45	55	65	75
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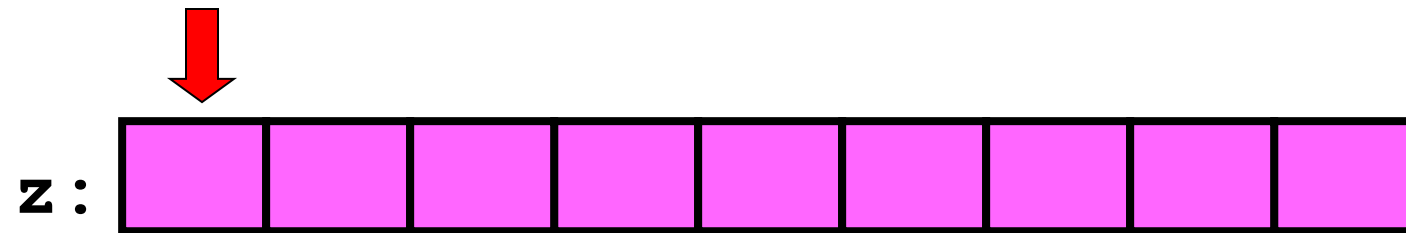
Merge



ix: [1]



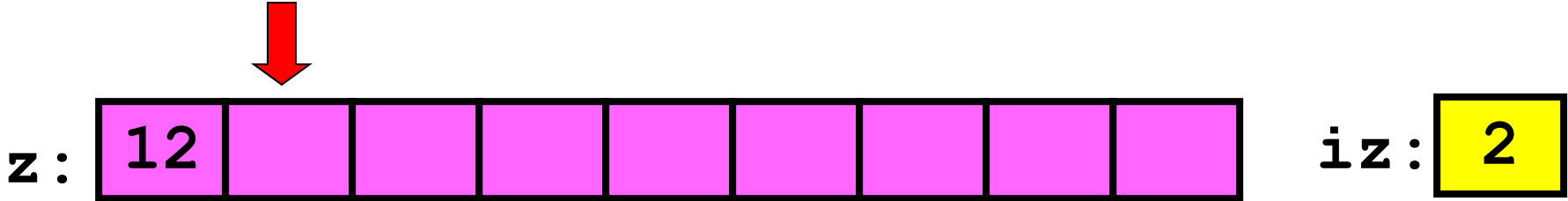
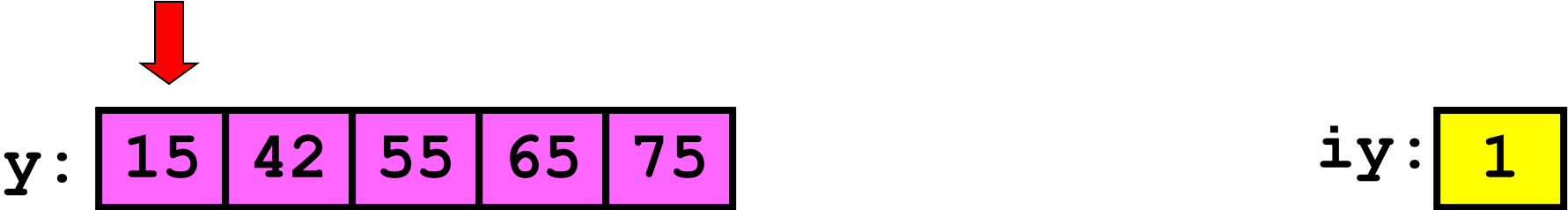
iy: [1]



iz: [1]

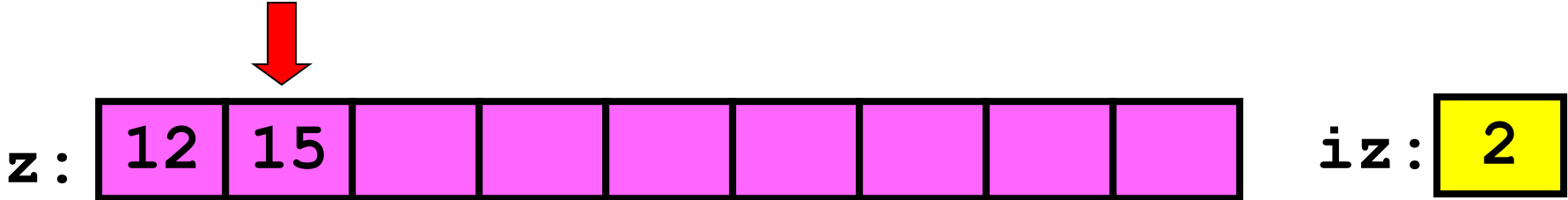
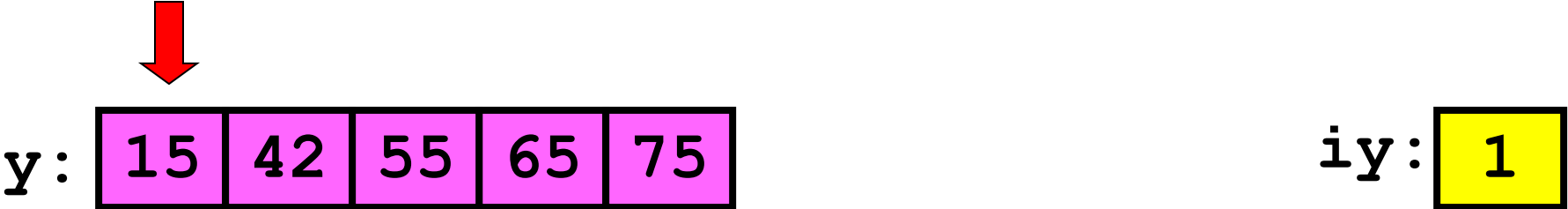
$ix \leq 4$ and $iy \leq 5$: $x(ix) \leq y(iy)$???

Merge



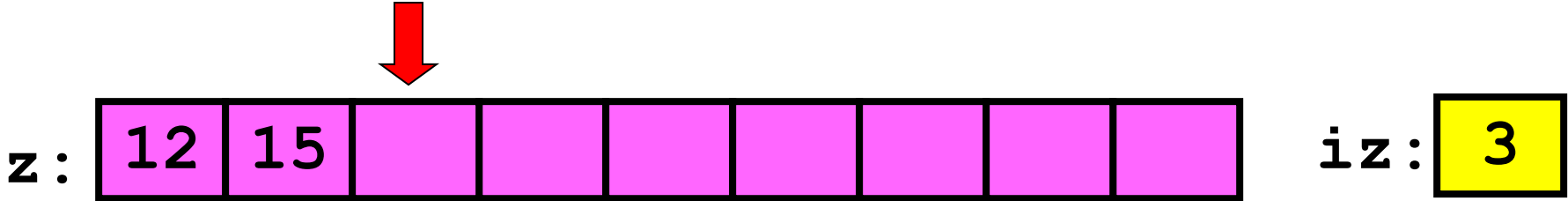
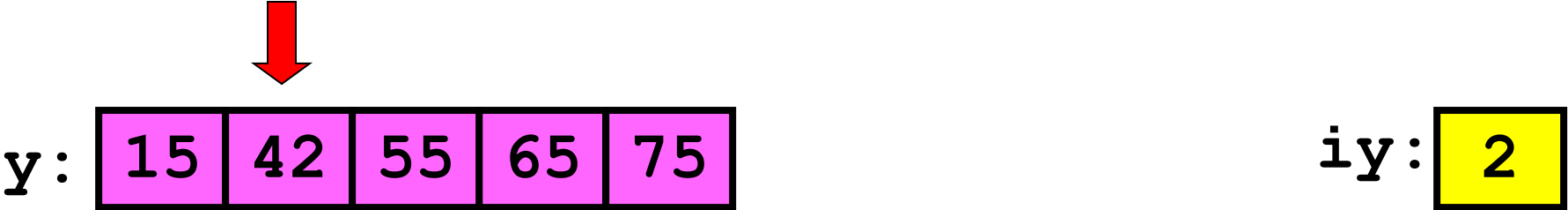
$ix \leq 4$ and $iy \leq 5$: $x[ix] \leq y[iy]$???

Merge



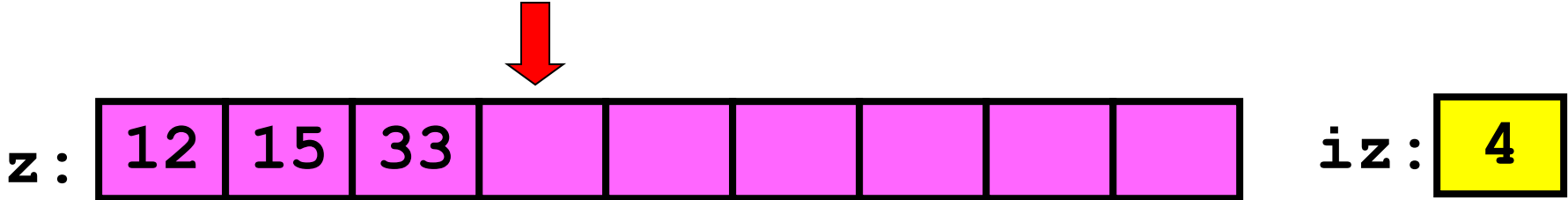
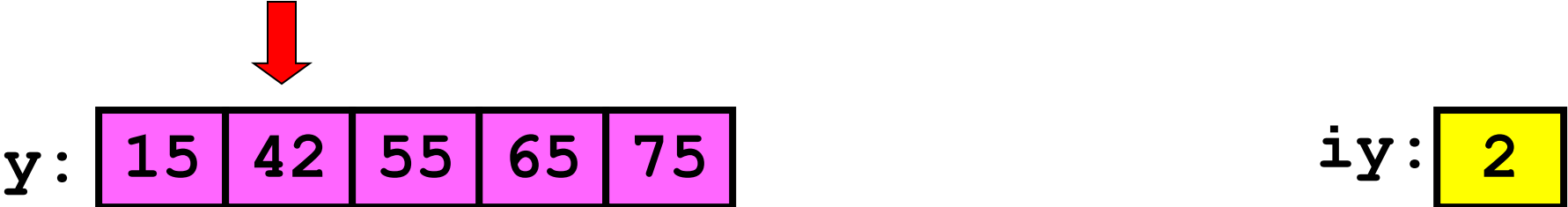
$ix \leq 4$ and $iy \leq 5$: $x(ix) \leq y(iy)$ **NO**

Merge



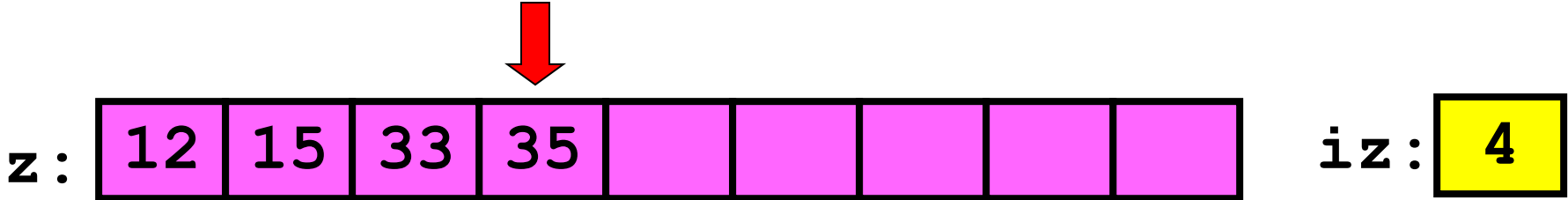
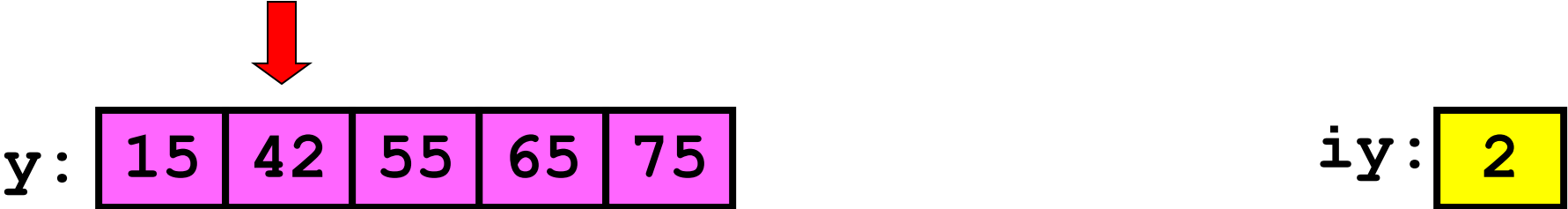
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Merge



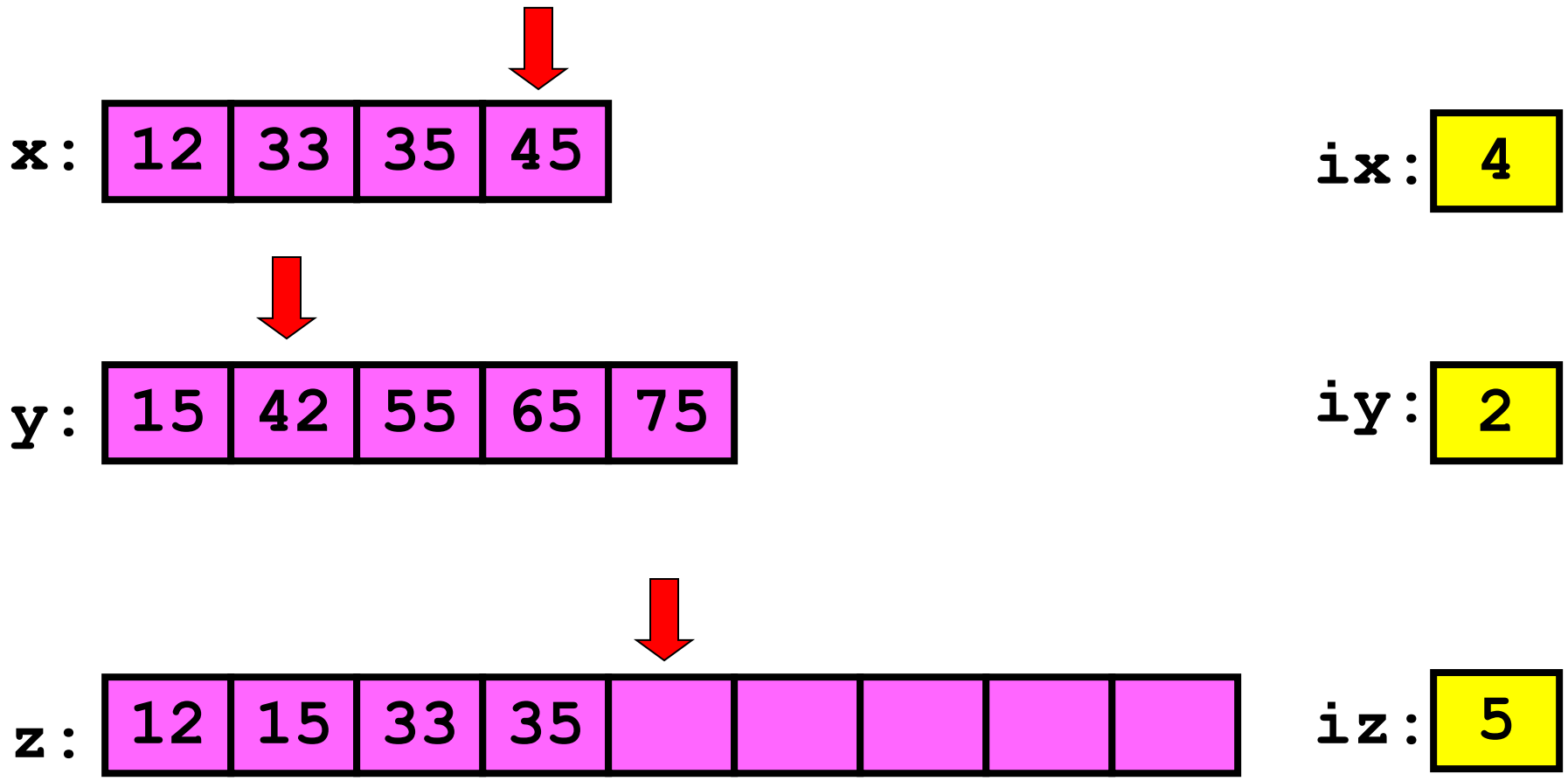
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Merge



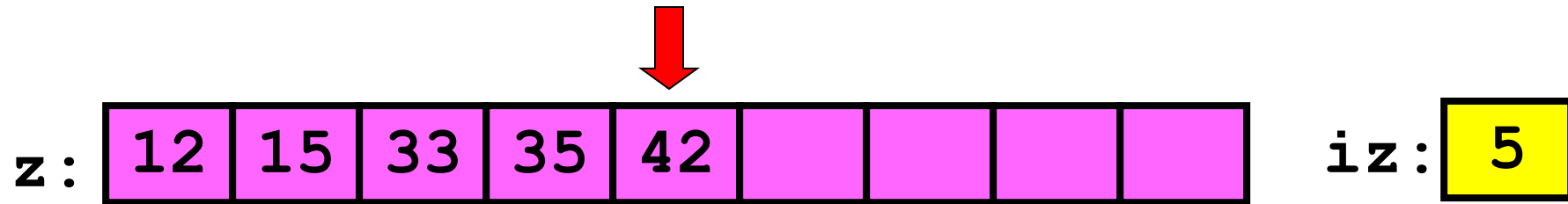
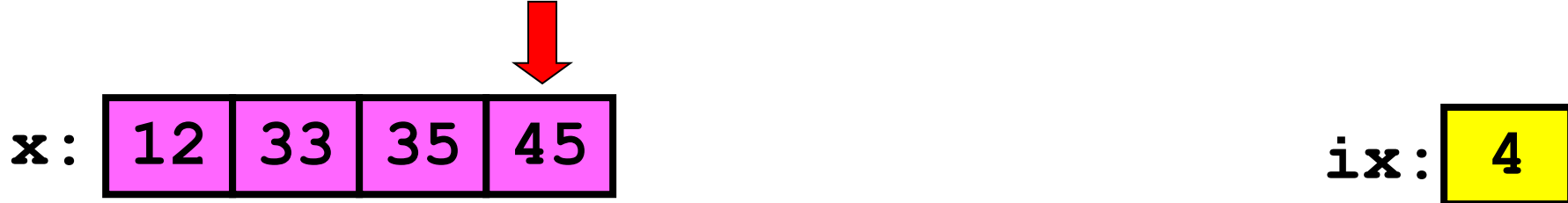
ix ≤ 4 and **iy** ≤ 5: **x(ix) ≤ y(iy)** **YES**

Merge



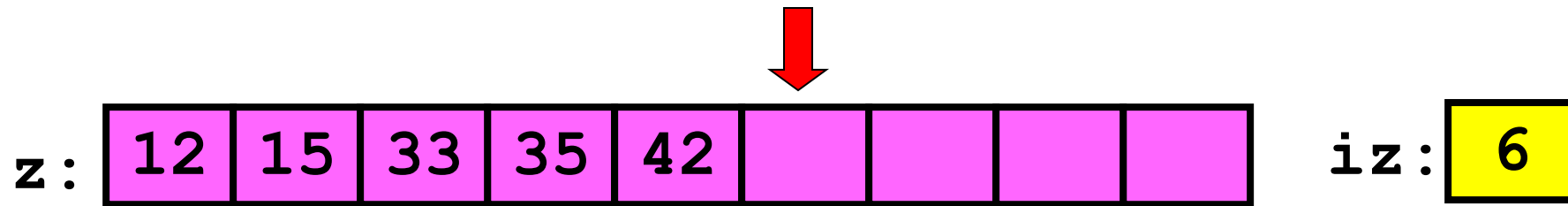
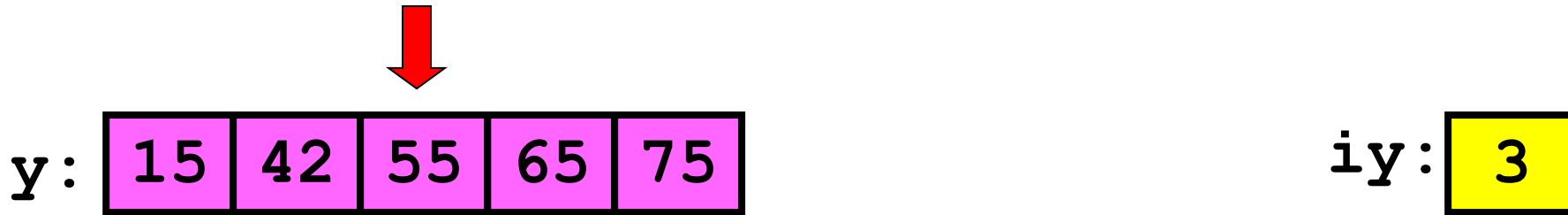
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Merge



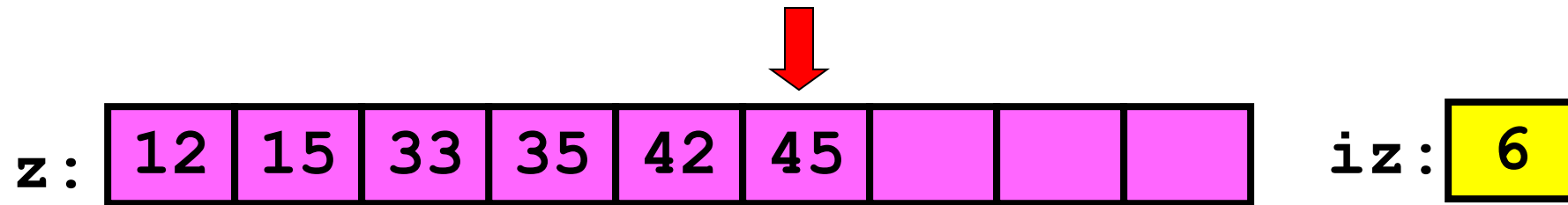
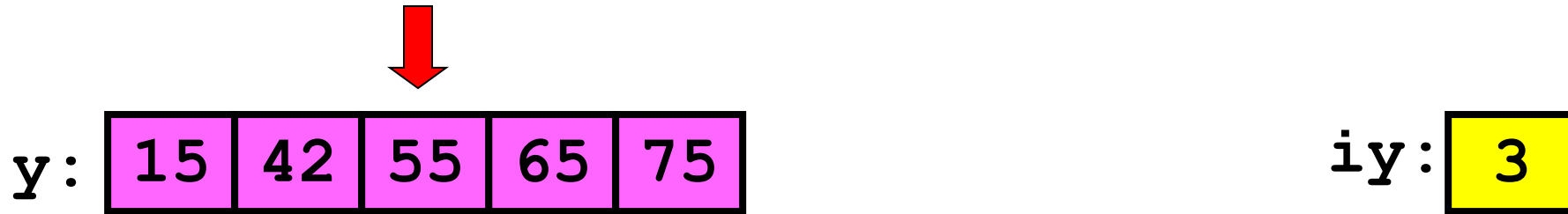
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Merge



$ix \leq 4$ and $iy \leq 5$: $x(ix) \leq y(iy)$???

Merge



$ix \leq 4$ and $iy \leq 5$: $x(ix) \leq y(iy)$ **YES**

Merge



x:

12	33	35	45
----	----	----	----

ix:

5



y:

15	42	55	65	75
----	----	----	----	----

iy:

3



z:

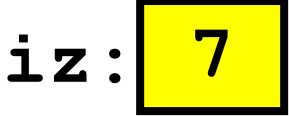
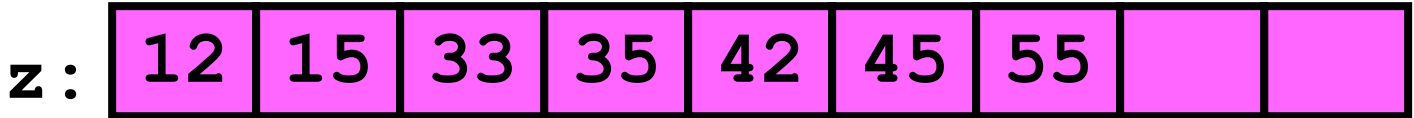
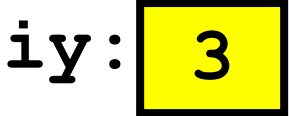
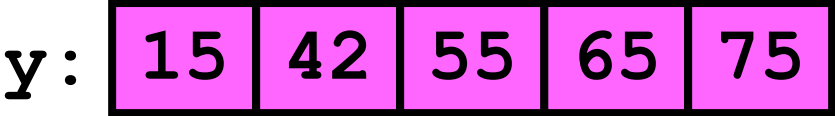
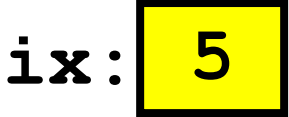
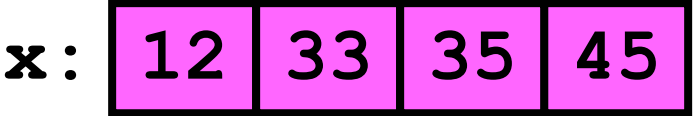
12	15	33	35	42	45			
----	----	----	----	----	----	--	--	--

iz:

7

ix > 4

Merge



`ix > 4: take y(iy)`

Merge

x:

12	33	35	45
----	----	----	----

ix:

5

y:

15	42	55	65	75
----	----	----	----	----

iy:

4

z:

12	15	33	35	42	45	55		
----	----	----	----	----	----	----	--	--

iz:

8

$iy \leq 5$

Merge



x:

12	33	35	45
----	----	----	----

ix:

5



y:

15	42	55	65	75
----	----	----	----	----

iy:

4



z:

12	15	33	35	42	45	55	65	
----	----	----	----	----	----	----	----	--

iz:

8

$iy \leq 5$

Merge

x:

12	33	35	45
----	----	----	----

ix:

5

y:

15	42	55	65	75
----	----	----	----	----

iy:

5

z:

12	15	33	35	42	45	55	65	
----	----	----	----	----	----	----	----	--

iz:

9

$iy \leq 5$

Merge

x:

12	33	35	45
----	----	----	----

ix:

5

y:

15	42	55	65	75
----	----	----	----	----

iy:

5

z:

12	15	33	35	42	45	55	65	75
----	----	----	----	----	----	----	----	----

iz:

9

$iy \leq 5$


```
function z = merge(x,y)
nx = length(x); ny = length(y);
z = zeros(1, nx+ny);
ix = 1; iy = 1; iz = 1;
```

```
function z = merge(x,y)
nx = length(x); ny = length(y);
z = zeros(1, nx+ny);
ix = 1; iy = 1; iz = 1;
while ix<=nx && iy<=ny

end
% Deal with remaining values in x or y
```

```
function z = merge(x,y)
nx = length(x); ny = length(y);
z = zeros(1, nx+ny);
ix = 1; iy = 1; iz = 1;
while ix<=nx && iy<=ny
    if x(ix) <= y(iy)
        z(iz) = x(ix); ix=ix+1; iz=iz+1;
    else
        z(iz) = y(iy); iy=iy+1; iz=iz+1;
    end
end
end
% Deal with remaining values in x or y
```

```
function z = merge(x,y)
nx = length(x); ny = length(y);
z = zeros(1, nx+ny);
ix = 1; iy = 1; iz = 1;
while ix<=nx && iy<=ny
    if x(ix) <= y(iy)
        z(iz) = x(ix); ix=ix+1; iz=iz+1;
    else
        z(iz) = y(iy); iy=iy+1; iz=iz+1;
    end
end
while ix<=nx % copy remaining x-values
    z(iz) = x(ix); ix=ix+1; iz=iz+1;
end
while iy<=ny % copy remaining y-values
    z(iz) = y(iy); iy=iy+1; iz=iz+1;
end
```

Merge sort: Motivation

If I have two helpers, I'd...

- Give each helper half the array to sort
- Then I get back the sorted subarrays and **merge** them.

Cost of dividing work

Suppose each comparison we make costs \$1

Given a vector with N elements,

- Insertion sort costs $\$N(N-1)/2$
- Merge costs $\$(N-1)$

(worst case)

Consider a vector with 8 elements

- Sorting by ourselves: \$26
- Sorting by delegating work:
 - Left delegate (4 elements): \$6
 - Right delegate (4 elements): \$6
 - Merge (8 elements): \$7
 - **Profit: \$7!**

Merge sort: Motivation

If I have two helpers, I'd...

- Give each helper half the array to sort
- Then I get back the sorted subarrays and **merge** them.

What if those two helpers each had two sub-helpers?

And the sub-helpers each had two sub-sub-helpers? And...

Subdivide the sorting task

H	E	M	G	B	K	A	Q	F	L	P	D	R	C	J	N
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

H	E	M	G	B	K	A	Q
---	---	---	---	---	---	---	---

F	L	P	D	R	C	J	N
---	---	---	---	---	---	---	---

Subdivide again



And again



H E M G

B K A Q

F L P D

R C J N

H E

M G

B K

A Q

F L

P D

R C

J N

And one last time



H E

M G

B K

A Q

F L

P D

R C

J N

Now merge



E H

G M

B K

A Q

F L

D P

C R

J N

H E

M G

B K

A Q

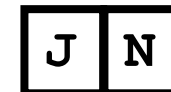
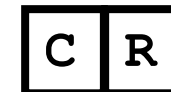
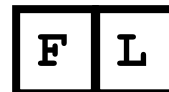
F L

P D

R C

J N

And merge again



And again



And one last time

A	B	C	D	E	F	G	H	J	K	L	M	N	P	Q	R
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

A	B	E	G	H	K	M	Q
---	---	---	---	---	---	---	---

C	D	F	J	L	N	P	R
---	---	---	---	---	---	---	---

Done!

A	B	C	D	E	F	G	H	J	K	L	M	N	P	Q	R
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---


```
function y = mergeSort(x)
% x is a vector.  y is a vector
% consisting of the values in x
% sorted from smallest to largest.

n = length(x);
if (task is trivial)
    % Base case
else
    % Divide work
    % Delegate subproblems

    % Merge results
end
```

```
function y = mergeSort(x)
% x is a vector.  y is a vector
% consisting of the values in x
% sorted from smallest to largest.

n = length(x);
if n==1
    y = x;
else
    % Divide work
    % Delegate subproblems

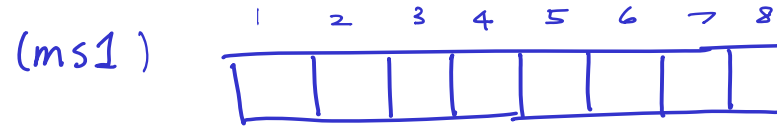
    % Merge results
end
```

```
function y = mergeSort(x)
% x is a vector.  y is a vector
% consisting of the values in x
% sorted from smallest to largest.
```

```
n = length(x);
if n==1
    y = x;
else
    m = floor(n/2);
    yL = mergeSort(x(1:m));
    yR = mergeSort(x(m+1:n));
    y = merge(yL, yR);
end
```

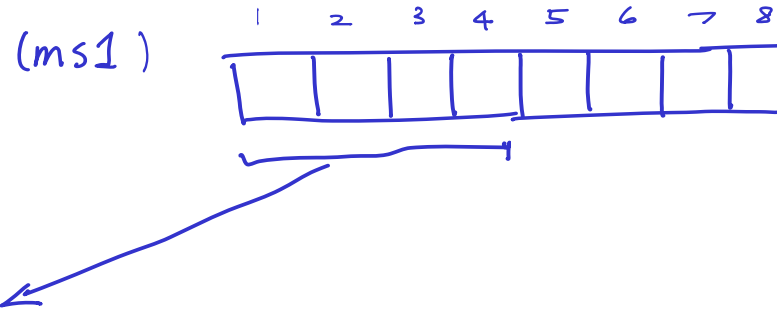
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mergeSort — 1st call



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mergeSort — 1st call

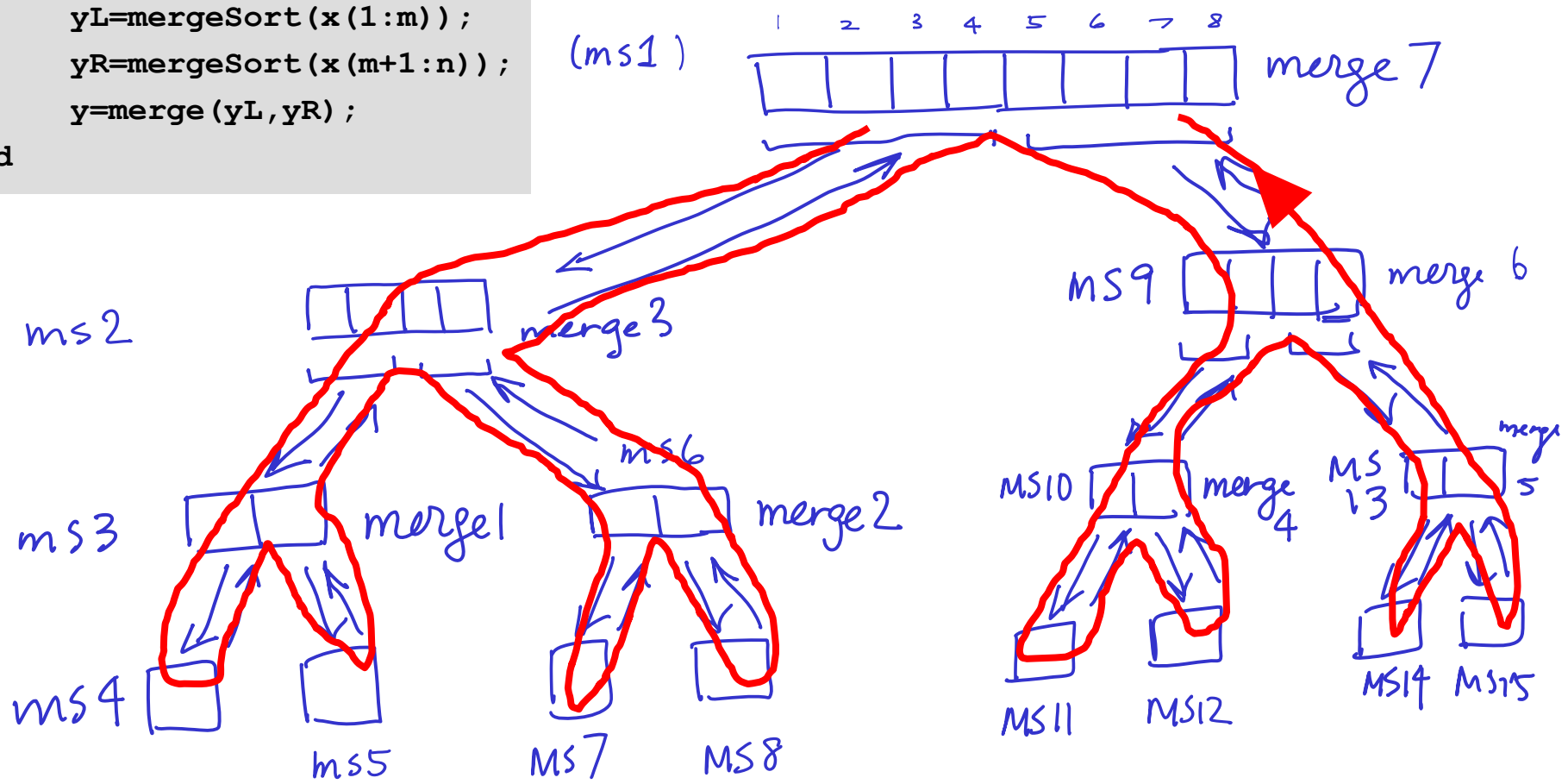


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    y=merge(yL,yR);
end

```

mergeSort - 1st call



How do merge sort and insertion sort compare?

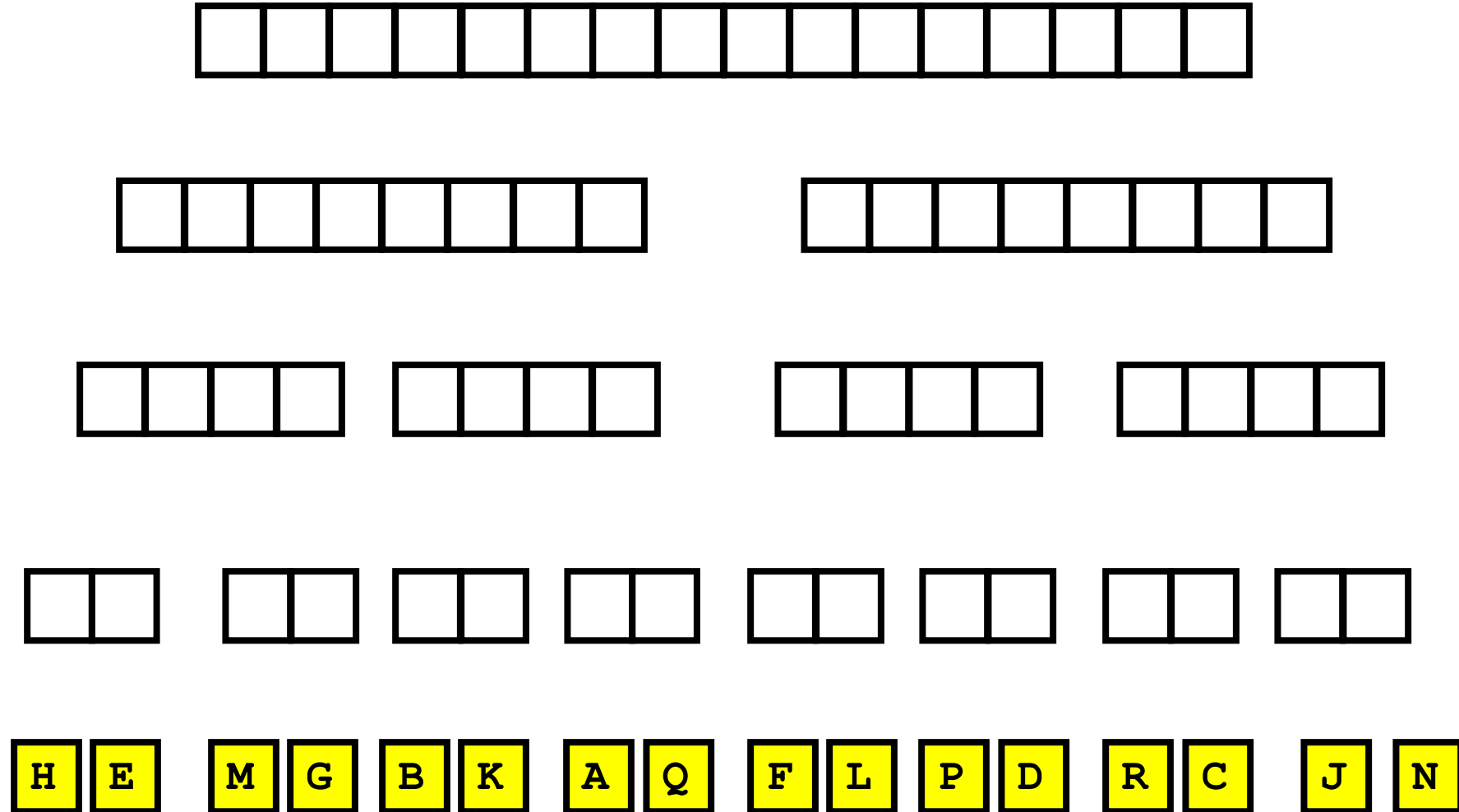
- Insertion sort: (worst case) makes k comparisons to insert an element in a sorted array of k elements. For an array of length N :
 $1+2+\dots+(N-1) = N(N-1)/2$, say N^2 for big N
- Merge sort:

```
function y = mergeSort(x)
% x is a vector.  y is a vector
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```

```
n = length(x);
if n==1
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    m = floor(n/2);
    yL = mergeSort(x(1:m));
    yR = mergeSort(x(m+1:n));
    y = merge(yL, yR);
end
```

All the comparisons between
vector values are done in merge

Merge sort: about $\log_2(N)$ “levels”;
about N comparisons each level



How do merge sort and insertion sort compare?

- Insertion sort: (worst case) makes i comparisons to insert an element in a sorted array of i elements. For an array of length N :
 $1+2+\dots+(N-1) = N(N-1)/2$, say N^2 for big N
- Merge sort: $N \cdot \log_2(N)$
- Insertion sort is done *in-place*; merge sort (recursion) requires extra memory (call frames plus merge area)

See `compareInsertMerge.m`

How to choose??

- Depends on application
- Merge sort is especially good for sorting **large data sets**
 - Easily adapted to work with files if data is too big for memory
- Sort “**stability**” matters for object handles (elements may compare equal, but are actually distinct)
 - Insertion, Merge are intrinsically stable. QuickSort is not, but MATLAB’s `sort()` does extra work to stabilize
- Insertion sort is “order N^2 ” at **worst case**, but what about an **average case**?
 - Insertion good for “fixing” a mostly sorted array, or adding just a few new elements

What we learned...

- Develop/implement **algorithms** for problems
- Develop programming skills
 - Design, implement, document, test, and debug
- Programming “tool bag”
 - Functions for reducing redundancy
 - Control flow (if-else; loops)
 - Recursion
 - Data structures, type
 - Graphics
 - File handling

What we learned... (cont'd)

- Applications and concepts
 - Image processing
 - Object-oriented programming—custom type
 - Sorting and searching (you should know the algorithms covered)
 - Approximation and error
 - Simulation, sensitivity analysis
 - Computational effort and efficiency

Where to go from here?

- [Mathworks.com](https://www.mathworks.com) – Many free tutorials available on specific topics, e.g., signal processing, Simulink, ..., etc.
- More detailed intro to scientific and engineering uses: “*Getting Started with MATLAB*” by Rudra Pratap. Excellent for independent, non-course-based learning
- Just play, i.e., experiment, with MATLAB programs! Many programs available in MATLAB “[Community](#)” forum “[File Exchange](#)”

Some courses for future consideration

- ENGRD/CS 2110 Object-oriented programming and data structure
 - CS 2111 Programming practicum
- CS 2800 Discrete Math (logic, proof, probability theory)
- CS 3220 Computational Mathematics for Computer Science
- Short language courses (e.g., Python, C++)

*Highly recommended
companion to CS2110*

Computing gives us *insight* into a problem

- Computing is not about getting one answer!
- We build models and write programs so that we can “play” with the models and programs, learning—gaining insights—as we vary the parameters and assumptions
- Good models require domain-specific knowledge (and experience)
- Good programs ...
 - are correct and have been thoroughly tested
 - are modular and cleanly organized
 - are well-documented
 - use appropriate data structures and algorithms
 - are reasonably efficient in time and memory

Gro

2!

Best wishes
and
good luck with all your exams!

