

- Previous Lecture:
 - Vectorized code
 - 2-d array—matrix
- Today's Lecture:
 - More examples on matrices
 - Optional reading: contour plot (7.2, 7.3 in *Insight*)
- Announcements:
 - Fall Break Mon & Tues: no lec, dis, office/consulting hrs. Attendance at 10/16 (W) dis is optional, but the exercise is required. Attend any of the 10/16 dis sections if you wish. Location is Upton B7 lab.
 - Optional review sessions: W 5-6:30p, W 7-8:30p. Locations TBA

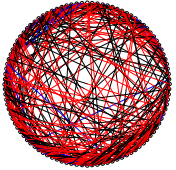
Storing and using data in *tables*

A company has 3 factories that make 5 products with these costs:

10	36	22	15	62
12	35	20	12	66
13	37	21	16	59

C

What is the best way to fill a given purchase order?



Connections between webpages

0	0	1	0	1	0	0
1	0	0	1	1	1	0
0	1	0	1	1	1	1
1	0	1	1	0	1	0
0	0	1	1	0	1	1
0	0	1	0	1	0	1
0	1	1	0	1	1	0

Lecture 13 5

Pattern for traversing a matrix M

```

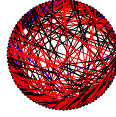
[nr, nc] = size(M)
for r= 1:nr
    % At row r
    for c= 1:nc
        % At column c (in row r)
        %
        % Do something with M(r,c) ...
    end
end
end
```

Lecture 14 6

Matrix example: Random Web

- N web pages can be represented by an N-by-N Link Array *A*.
- *A*(*i*,*j*) is 1 if there is a link on webpage *j* to webpage *i*
- Generate a random link array and display the connectivity:
 - There is no link from a page to itself
 - If *i*≠*j* then *A*(*i*,*j*) = 1 with probability $\frac{1}{1+|i-j|}$

⇒ There is more likely to be a link if *i* is close to *j*



Lecture 13 7

```

function A = RandomLinks(n)
% A is n-by-n matrix of 1s and 0s
% representing n webpages

A = zeros(n,n);
for i=1:n
    for j=1:n
        r = rand(1);
        if i~j && r<= 1/(1 + abs(i-j));
            A(i,j) = 1;
        end
    end
end
end
```

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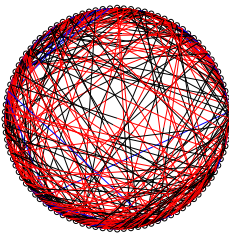
Random web
N = 20

```

0 1 1 1 0 0 0 0 0 1 0 0 1 0 0 0 0 0 0 0
1 0 0 0 1 0 0 0 1 1 1 0 0 0 0 0 0 0 1 0 0
0 1 0 1 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
0 0 1 0 1 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
0 0 0 1 0 0 0 0 0 0 1 1 0 0 0 0 0 0 0 0
0 0 0 0 0 0 0 0 0 0 0 0 0 0 1 0 1 0 0 0
0 1 1 1 1 1 0 0 0 1 0 1 1 0 0 0 0 0 0 0
0 0 0 0 0 0 1 0 0 0 0 1 0 0 0 0 0 0 0 1 1
0 1 0 0 0 0 0 0 0 0 1 0 0 0 1 0 0 0 1 0 0
0 0 0 0 0 0 0 1 1 0 1 0 0 0 0 0 0 0 0 0 1
0 0 0 0 0 0 1 0 0 0 0 0 1 1 0 0 0 0 0 0 0
0 0 0 0 0 0 1 0 0 1 0 0 1 1 0 0 0 0 0 0 0
0 0 0 0 0 0 1 0 0 0 0 0 0 0 0 1 1 0 0 0 0
0 0 0 0 0 1 0 1 0 0 0 0 0 1 0 0 0 0 0 0 1
0 0 0 0 0 0 1 0 0 0 1 0 0 0 0 0 0 1 0 1 0
0 1 0 0 0 0 0 0 1 0 0 0 0 0 1 0 1 0 1 1 0
0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 1 1 0 0 1
0 0 0 0 0 0 1 0 0 0 0 0 0 0 0 0 0 0 0 0 0
0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 1 0 1 0
```

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Represent the web pages graphically...

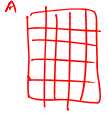


Bidirectional links are blue. Unidirectional link is black as it leaves page j , red when it arrives at page i .

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```

for i = 1:n
  for j = 1:n
    end
  end
end
    
```



Lecture 13 12

A Cost/Inventory Problem

- A company has 3 factories that make 5 different products
- The cost of making a product varies from factory to factory
- The inventory/capacity varies from factory to factory

Lecture 14 14

Problems

A customer submits a purchase order that is to be filled by a single factory.

1. How much would it cost a factory to fill the order?
2. Does a factory have enough inventory/capacity to fill the order?
3. Among the factories that can fill the order, who can do it most cheaply?

Lecture 14 15

Cost Array

10	36	22	15	62
12	35	20	12	66
13	37	21	16	59

C

The value of $C(i, j)$ is what it costs factory i to make product j .

Lecture 14 16

Inventory (or Capacity) Array

38	5	99	34	42
82	19	83	12	42
51	29	21	56	87

Inv

The value of $Inv(i, j)$ is the inventory in factory i of product j .

Lecture 14 17

Purchase Order

PO

1	0	12	29	5
---	---	----	----	---

The value of $PO(j)$ is the number of product j 's that the customer wants

Lecture 14 18

10	36	22	15	62
12	35	20	12	66
13	37	21	16	59

C

PO

1	0	12	29	5
---	---	----	----	---

Cost for factory i:

```
s = 0; %Sum of cost
for j=1:5
    s = s + C(i,j)*PO(j)
end
```

Lecture 14 22

Encapsulate...

```
function TheBill = iCost(i,C,PO)
% The cost when factory i fills the
% purchase order

nProd = length(PO);
TheBill = 0;
for j=1:nProd
    TheBill = TheBill + C(i,j)*PO(j);
end
```

Lecture 14 23

Finding the Cheapest

```
iBest = 0; minBill = inf;
for i=1:nFact
    iBill = iCost(i,C,PO);
    if iBill < minBill
        % Found an Improvement
        iBest = i; minBill = iBill;
    end
end
```

Lecture 14 24

inf – a special value that can be regarded as positive infinity

$x = 10/0$ assigns **inf** to x

$y = 1+x$ assigns **inf** to y

$z = 1/x$ assigns 0 to z

$w < \text{inf}$ is always true if w is numeric

Lecture 14 25

Inventory/Capacity Considerations

What if a factory lacks the inventory/capacity to fill the purchase order?

Such a factory should be excluded from the find-the-cheapest computation.

Lecture 14 26

Who Can Fill the Order?

	38	5	99	34	42	Yes
Inv	82	19	83	12	42	No
	51	29	21	56	87	Yes
PO	1	0	12	29	5	

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Wanted: A True/False Function

```

    graph LR
      i --> iCanDo
      Inv --> iCanDo
      PO --> iCanDo
      iCanDo --> DO
  
```

DO is "true" if factory *i* can fill the order.
DO is "false" if factory *i* cannot fill the order.

Lecture 14 28

Example: Check inventory of factory 2

	38	5	99	34	42
Inv	82	19	83	12	42
	51	29	21	56	87
PO	1	0	12	29	5

Method 1: check the inventory for every product

Lecture 14 29

Initialization

	38	5	99	34	42	
Inv	82	19	83	12	42	DO 1
	51	29	21	56	87	
PO	1	0	12	29	5	

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Still True...

	38	5	99	34	42	
Inv	82	19	83	12	42	DO 1
	51	29	21	56	87	
PO	1	0	12	29	5	

DO = DO && (Inv(2,1) >= PO(1))

Lecture 14 31

Encapsulate...

```

function DO = iCanDo(i,Inv,PO)
% DO is true if factory i can fill
% the purchase order. Otherwise, false

nProd = length(PO);
DO = 1;
for j = 1:nProd
    DO = DO && ( Inv(i,j) >= PO(j) );
end
  
```

Lecture 14 36

Encapsulate...

```
function DO = iCanDo(i,Inv,PO)
% DO is true if factory i can fill
% the purchase order. Otherwise, false
nProd = length(PO);
j = 1;
while j<=nProd && Inv(i,j)>=PO(j)
    j = j+1;
end
DO = _____;
```

DO should be true when...

A j < nProd

B j == nProd

C j > nProd

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Back To Finding the Cheapest

```
iBest = 0; minBill = inf;
for i=1:nFact
    iBill = iCost(i,C,PO);
    if iBill < minBill
        % Found an Improvement
        iBest = i; minBill = iBill;
    end
end
```

Don't bother with this unless there is sufficient inventory.

Lecture 14 40

Back To Finding the Cheapest

```
iBest = 0; minBill = inf;
for i=1:nFact
    if iCanDo(i,Inv,PO)
        iBill = iCost(i,C,PO);
        if iBill < minBill
            % Found an Improvement
            iBest = i; minBill = iBill;
        end
    end
end
```

See Cheapest.m
for alternative implementation

Lecture 14 41

Finding the Cheapest

	10	36	22	15	62	1019	Yes
C	12	35	20	12	66	930	No
C	13	37	21	16	59	1040	Yes
PO	1	0	12	29	5	↑	↑

As computed by iCost

As computed by iCanDo

Lecture 14 43

```
% Given an nr-by-nc matrix M.
% What is A?
for r= 1: nr
    for c= 1: nc
        A(c,r)= M(r,c);
    end
end
```

- A A is M with the columns in reverse order
- B A is M with the rows in reverse order
- C A is the transpose of M
- D A and M are the same

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