17. Searching and Sorting

Topics:

Linear Search

Binary Search

Measuring Execution Time

The Divide and Conquer Framework

Merge Sort

Search

Examples:

Is this song in that playlist?

Is this number in that phone book?

Is this name in that phone book?

Is this fingerprint in that archive of fingerprints?

Is this photo in that yearbook?

More on Using Phone Books

The Manhatten phone book has 1,000,000+ entries.

How is it possible to locate a name by examining just a tiny, tiny fraction of those entries?



There must be a great search algorithm behind the scenes.

LinSearch: The Spec

```
def LinSearch(x,a):
    """ Returns an int k with the
    property that a[k]==x is True.
    If no such k exists, then
    k==-1.
```

PreC: a is a nonempty list of ints and x is an int.

** ** **

0 1 2 3 4 5 6 7 8 9 10 11 a->86 73 43 35 23 45 42 62 15 25 51 35

```
def LinSearch(x,a):
    for k in range(len(a)):
        if x == a[k]:
            return k
    return -1
```

0 1 2 3 4 5 6 7 8 9 10 11 a->86 73 43 35 23 45 42 62 15 25 51 35

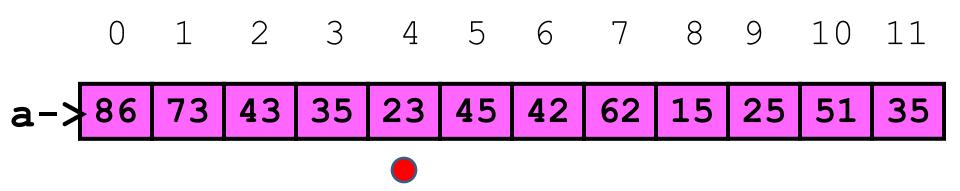
```
def LinSearch(x,a):
    for k in range(len(a)):
        if x == a[k]: Nope
        return k
    return -1
```

```
def LinSearch(x,a):
    for k in range(len(a)):
        if x == a[k]: Nope
        return k
    return -1
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```

0 1 2 3 4 5 6 7 8 9 10 11 a->86 73 43 35 23 45 42 62 15 25 51 35

```
def LinSearch(x,a):
    for k in range(len(a)):
        if x == a[k]: Nope
        return k
    return -1
```



```
def LinSearch(x,a):
    for k in range(len(a)):
        if x == a[k]: Yup
        return k
    return -1
```

0 1 2 3 4 5 6 7 8 9 10 11 a->86 73 43 35 23 45 42 62 15 25 51 35

```
def LinSearch(x,a):
    for k in range(len(a)):
        if x == a[k]:
            return k Alldone
    return -1
```

Linear Search: No Match Case

```
def LinSearch(x,a):
    for k in range(len(a)):
        if x == a[k]: Nope
        return k
    return -1
```

Linear Search: No Match Case

```
x-> 7
```

```
def LinSearch(x,a):
    for k in range(len(a)):
        if x == a[k]:
            return k
    return -1 Yup
```

Linear Search: While Implementation

```
def LinSearchW(x,a):
    k=0
    while k<len(a) and a[k]!=x:
        k+=1
    if k==len(a):
        return -1
    else:
        return k</pre>
```

Binary Search

Now we assume that the list to be searched is sorted from little to big.

```
a = [10,20,40,60,90]
a = ['brown','dog','fox','lazy','quick','the']
```

Back to Using Phone Books

The Ithaca phone book has 10,000+ entries.

The Manhatten phone book has 1 000 000+ entries

1,000,000+ entries. But it does not take $100 \times longer$ to look something up. Why?

wide at	SuperPages.com	195	Car C	
entiti A deliver 7	Cartage New England Inc	Carter F 24 Hillock Ros 02131	Carter Nella E	
17 566-1282	26 Allen Ln Ipswich 01938978 356-9960	Faye & Ricky 357 Columbus Av Bos 02116617 437-7331	333 Maschsts Av Bos 02115617 267-6483 Nicholas S F	
81 447-4101	Cartagema Lydia 18 Jewett Ros 02131	Francis S 134 Temple W Rox 02132 617 323-6781	115 Randolph Av Mil 02186	
01 447-4101	Cartagena Avith	Franklin & Anne	Nick 21 Fairfield Bos 02116	
00 257-9981	9 Bancroft Rox 02119617 442-9780	221 Mt Auburn Cam 02138617 354-0798	Nick & Debbi	
30 237 7701	B Hyd 02136	Fred 42 Haverford Jam 02130 617 524-3078	196 Herrick Rd Newton 02459617 527-0480	
17 566-1282	Jessica 50 Decatur Cha 02129617 241-0152	Fred 96 Hinckley Rd Mil 02186617 698-1343	Nicole617 698-0713	
17 364-5188	Lucilla 174 Harvard Cam 02139 617 491-5621	G & R 8 Verdun Dor 02124617 436-8906	Norman G	
101656 BB 2007	M 95 Rowe Ros 02131 617 323-9713	G T 27 Franklin Av Som 02145617 623-7121	38 Chickatawbut Dor 02122 617 822-1203	
361-0380	Melvin 501 Green Cam 02139617 576-1061	Gayle 25 Frontenac Dor 02124 617 825-0322	P 94 Crestwood Pk Rox 02121617 427-4754	
	Carte Nicholas	Geo S 115 Moss Hill Rd Jam 02130617 522-3215	P E 501 E Sixth S Bos 02127	
17 566-4548	18 Appleton Boston 02116	George 125 Nashua Bos 02114617 367-9548	P L 44 Hutchings Rox 02121	
	Cartegena O 4 Milford Bos 02118617 338-8219	Carter Halliday Associate	P R 91 Bynner Jam 02130617 983-8692	
17 628-8248	Carten Thos J Sr & Claire	107 S Street Bos 02111	Paul & Constance 114 Anawan Av W Rox 02132	
	1 Paradise Rd Mil 02186 617 698-6163	Carter Harry F 26 Runno Brk Rd W Rox 02132 617 325-5465	Paul E 501 F Sixth St S Bos 02127617 323-2030	
17 445-5116	Thomas & Kathleen	26 Runng Brk Rd W Rox 02132 617 325-5465 Carter Hide Co Inc	Paul M 27 Union Bri 02135617 787-2115	
N. C.	50 Thompson Ln Mil 02186		Carter Pile Driving Inc 17 Beaver Ct	
7 822-2982	Carter A Ros 02131	146 Summer Bos 02110	Framingham 01702 Wellesley TelNo-781 235-8488	
17 427-5712	A Roxbury	Horace	Carter Prudence	
7 569-2698	A 260 Putnam Av Cambridge 02139 617 492-4174	241 Walnut Av Roxbury 02119617 442-5307	46 Franklin Watertown 02172 617 393-3782	
17 667-5190	A M 255 Maschets Av Bos 02115 617 492-4174	Howard Jr 26 Notre Dme Rox 02119.617 445-5552	Prudence	
1,001-2140	A M 255 Mascrists AV Bos 02115 017 200-7155 Adams 361 Centre St Mil 02186 617 698-9074	J Cam	46 Franklin Watertown 02172 617 926-7063	
17 569-1417	Alice 108 Kilmarnock Bos 02215 617 425-0193	J 15 Chatham Bro 02446	Reginald	
itu Dr	Alice 45 Market Cambridge 02139 617 945-2711	J 518 Harvard Bro 02446617 730-9483	106 Brunswick Dorchester 02121617 541-2843	
17 338-9110	Andrew F 62 Vinal Av Som 02143 617 625-7623	J 775 Vfw Pkwy West Roxbury 02132 617 323-5574	Renee & Andrew	
7 825-9195	Carter Anne MD	Carter J Jacques MD	10 Walnut Bos 02108617 720-3765	
Miles register	1101 Beacon Bro 02446617 739-1022	1 Brookline Pl Bro 02446	Carter Rice Dowd	
17 296-1593	Carter Athens	Carter J M	Bulldey Dunton Publishing 163 Main Wilmington 01887 Toll Free-Dial '1' & Then	
	272 Newbury Boston 02116	1410 Columbia Rd S Bos 02127 617 464-1040	Cust Svc-Industrial Prod 613 Main Wilmington	
17 670-2078	B E 68 Gladeside Av Mat 02126 617 296-6911	Carter J M Ornamental Ironworks	Toll Free-Dial '1' & Then800 619-7447	
17 623-9001	Carter Barbara L MD	CallPembroke TelNo-617 436-5353	Cust Svc-Printing 613 Main Wilmington	
Sec. March 1999	Tufts-New England Medical Center Bos 02111 Call	Carter J Veal Co 48 Newmarket Sq Rox 02118617 442-1775	Toll Free-Dial '1' & Then800 648-7447	
17 296-4725	Carter Becky 80s 02114	Carter James	Headquarters 613 Main Wilmington 01887	
2 540 1501	Bernard J	1573 Cambridge St Cam 02138617 492-1214	Call	
17 542-1521	112 Gladstone E 8os 02128617 567-3430	James 182 Fisher Av Roxbury 02120617 739-2193	Ingalls Cronin 163 Main Wilmington 01887 Toll Free-Dial '1' & Then	
17 364-5232	Bithiah 25 Medway Dor 02124617 298-8713	James	Carter Richard	
17 541-5649	Blake 26 Mt Vernon Bos 02108617 367-9931	37 Gold Star Rd Cambridge 02140 617 876-8841	1079 Commwith Av Brighton 02215 617 987-0836	
17 341 3047	Carter Broadcasting Co	Jas L 14 Roseberry Rd Mat 02126617 361-0773	Richard A 97 Mt Vernon Bos 02108617 566-7293	
7 739-2662	20 Park Piz Bos 02116 617 423-0210	Jane 114 Adena Rd Newton 02465617 964-0435	Carter Richard A MD	
27 2302	Carter & Burgess Consultants Inc	Jeffrey 41 Warren Av Bos 02116617 426-5994	170 Commwith Av Bos 02116	
17 879-0030	23 East St Cam 02141617 225-0200	John 11 Mansfield Bri 02134 617 987-2163	Carter Richard K	
7 541-3948	Carter C 2000 Comnwith Av Bri 02135 617 782-2118	John 327 Summer Bos 02210617 423-4334	15 Mercer S Bos 02127617 268-0448	
17 436-1513	C 228 Faywood Av East Boston 02128617 569-1545	John 40 Westwind Rd Dor 02125 617 282-1235	Robert L 175 Richdale Av Cam 02140. 617 864-1535	
17 569-4119	C 359 Harvard Cam 02138	June O 329 A Summit Av Bri 02135 617 734-6109	Roger 150 St Botolph Bos 02115617 424-6148	
ton 02128	C 610 Walk Hill Mat 02126617 296-6392	K 38 Browning Av Dorchester 02124 617 265-8456	Roy 44 Concord Av Cam 02138617 491-6115	
00 569-8782	C & M 43 Burroughs Jam 02130 617 524-9558	K 17 Esmond Dorchester 02121617 282-1593	Royce 18 Seminary Cha 02129 617 241-0418	

Key Idea: Repeated Halving

To Derek Jeter's number...

```
B = phone book
while (B is longer than 1 page):
   1. P = middle page of B
   2. Let Q be the first name on P
   3. if 'Jeter" comes before Q:
         Rip away the 2<sup>nd</sup> half of B
      else:
         Rip away the 1st half of B.
Scan remaining page P line-by-line for 'Jeter'
```

What Happens to Phone Book Length?

```
Original: 3000 pages
```

After 1 rip: 1500 pages

After 2 rips: 750 pages

After 3 rips: 375 pages

After 4 rips: 188 pages

After 5 rips: 94 pages

After 12 rips: 1 page

Binary Search

The idea of repeatedly halving the size of the "search space" is the main idea behind the method of binary search.

An item in a sorted array of length n can be located with approximately log_2 n comparisons.

What is $log_2(n)$?

n	ceil(log ₂ (n))	
10	4	
100	7	
1000	10	
10000	14	
100000	17	
1000000	20	

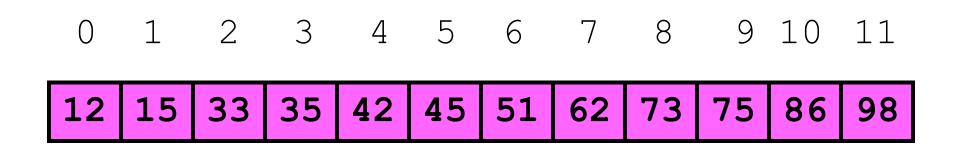
BinSearch: The Spec

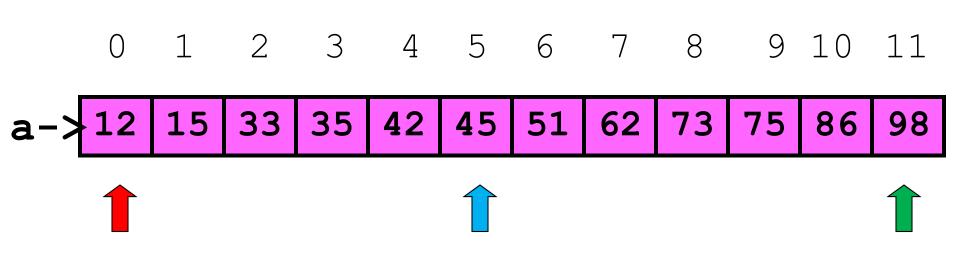
```
def BinSearch(x,a):
    """ Returns an int k with the
    property that a[k]==x is True.
    If no such k exists, then
    k==-1.
```

PreC: a is a nonempty list of ints that is sorted from smallest to largest. x is an int.

77 77 77

Example: Does this List have an Element With Value Equal to 70?





L: 0

 $a[Mid] \le x ????$

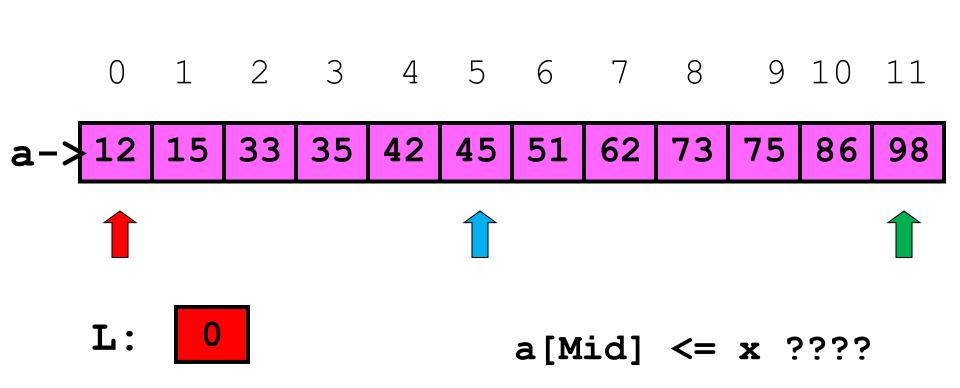
Mid: 5

R: 11 x: 70

Mid = (L+R)/2

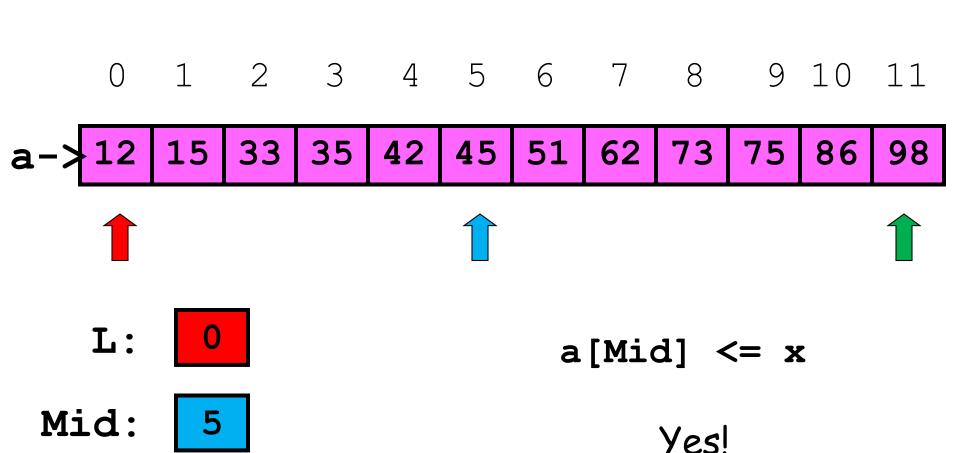
The Midpoint Computations

L	R	(L+R)/2
0	11	5
2	6	4
1	100	50



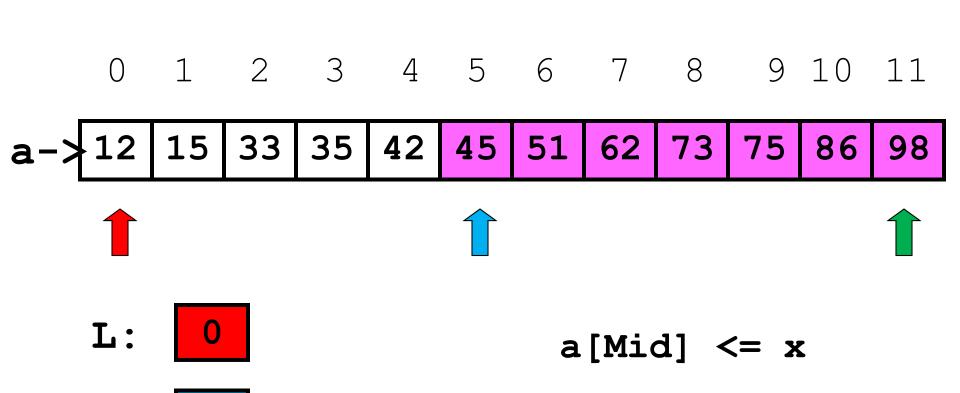
Mid: 5

R: 11 x: 70



R: 11 x: 70

So throw away
The "left half"

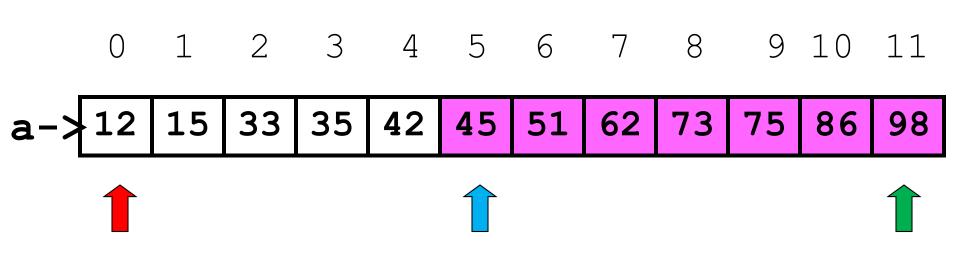


Mid: 5

R: 11

x: 70

Yes! So throw away The "left half"



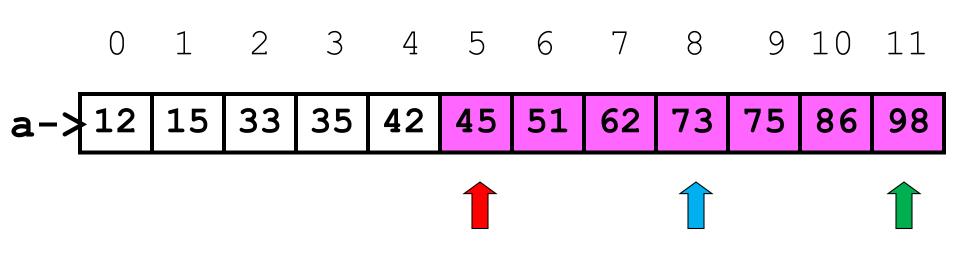
L: 0

 $a[Mid] \le x$

Mid: 5

Revise L and Mid

R: 11 x: 70

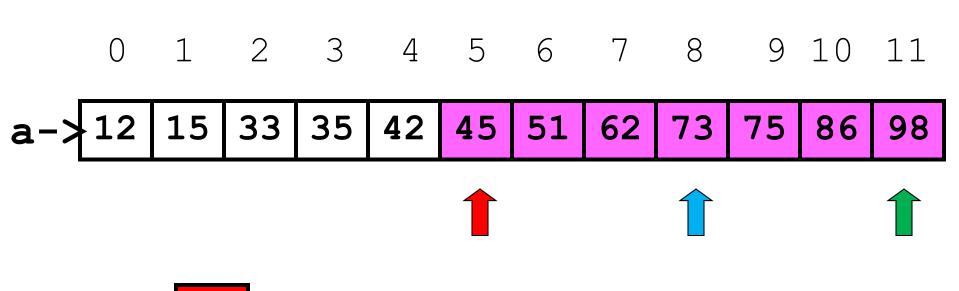


L: 5

 $a[Mid] \le x ???$

Mid: 8

R: 11 x: 70



L: 5

 $a[Mid] \le x$

Mid: 8

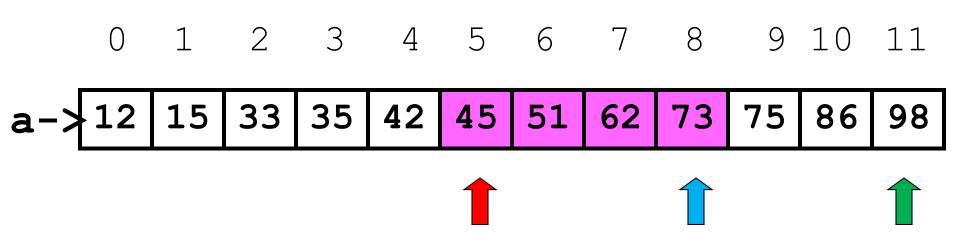
No

R: 11

x: 70

So throw away the "right half"

Let's Look For x = 70



L: 5

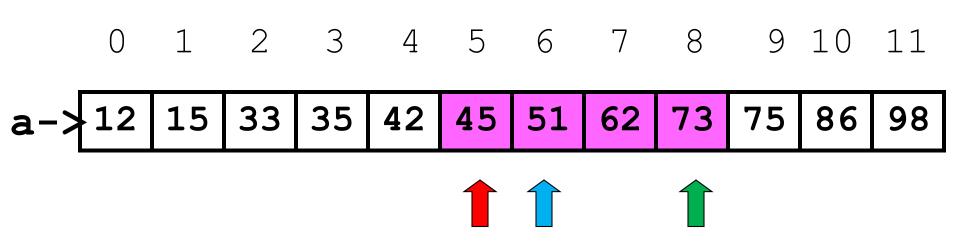
 $a[Mid] \le x$

Mid: 8

Revise R and Mid

R: 11 x: 70

Let's Look For x = 70



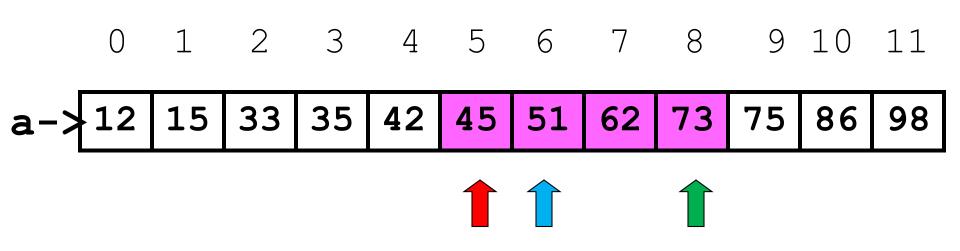
L: 5

 $a[Mid] \le x$

Mid: 6

Revise R and Mid

R: 8 x: 70

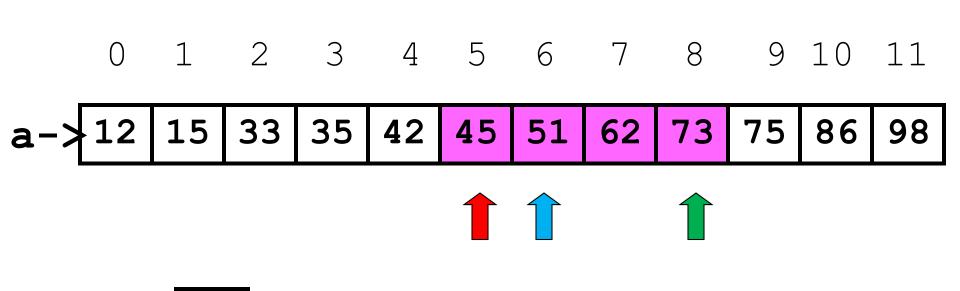


L: 5

 $a[Mid] \le x ????$

Mid: 6

R: 8 x: 70



L: 5

 $a[Mid] \le x$

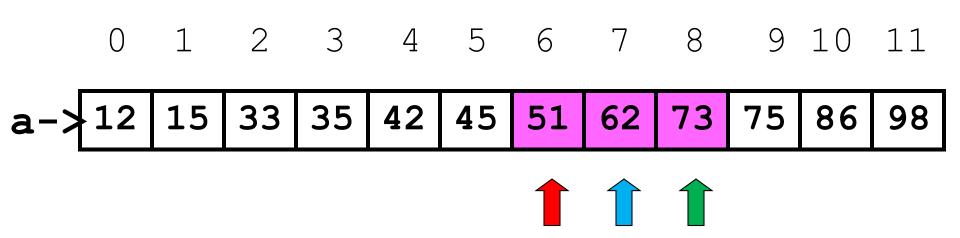
Mid: 6

Yes

R: 8

K: 70

Throw away the Left half



L: 6

 $a[Mid] \le x$

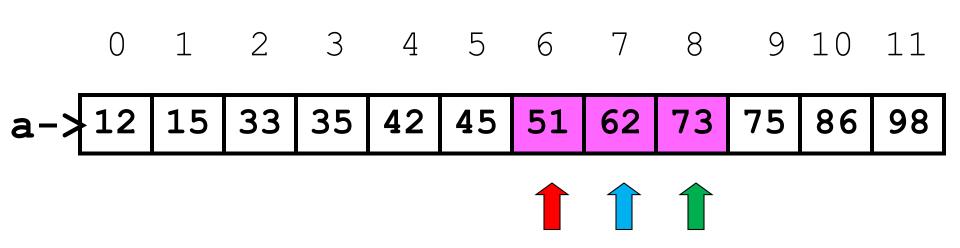
Mid: 7

Yes

R: 8

x: 70

Let's Look For x in a



L: 6

 $a[Mid] \le x$

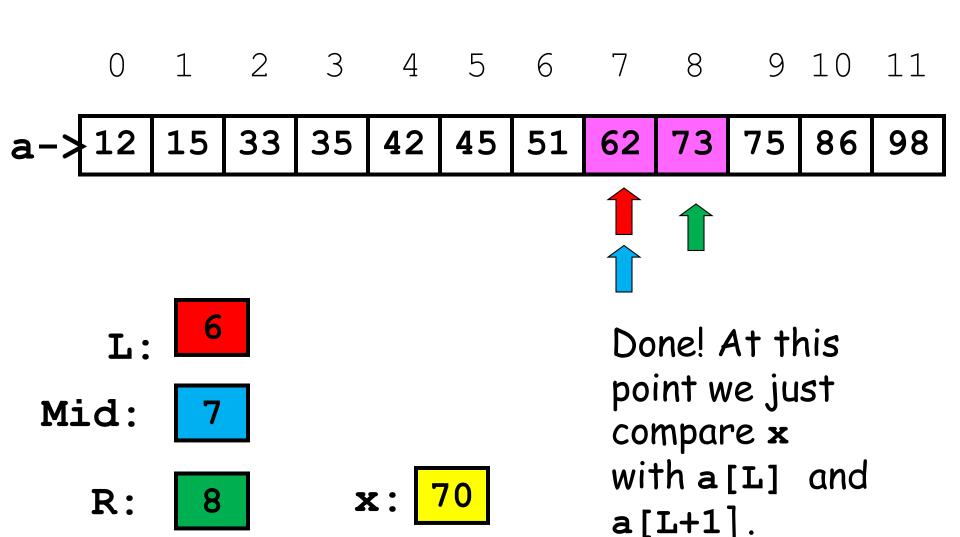
Mid: 7

Throw away the left half

R: 8

x: 70

Let's Look For x in a



What We Just Did

```
R = len(a)-1
while R-L > 1:
                           A Loop
     \# a[L] <= x <= a[R]
                           Invariant
      Mid = (L+R)/2
      if x \le a[mid]:
           R = Mid
      else:
           L = Mid
```

What We Just Did

```
R = len(a)-1
while R-L > 1:
    \# a[L] <= x <= a[R]
     Mid = (L+R)/2
     if x \le a[mid]:
          R = Mid
     else:
          L = Mid
```

What We Just Did

```
L = 0
R = len(a)-1
while R-L > 1:
    \# a[L] <= x <= a[R]
     Mid = (L+R)/2
     if x \le a[mid]:
          R = Mid
     else:
          L = Mid
```

After the Loop Ends



This is True: $a[L] \le x \le a[L+1]$

After the Loop Ends

```
a[L]
                       a[L+1]
if x==a[L]:
   return L
elif x==a[L+1]:
   return L+1
else:
   return -1
```

Measuring Execution Time

We now have two ways to search a list:

LinSearch(x,a) BinSearch(x,a)

Intuition: BinSearch much faster.

Can we quantify this with a "stop watch"?

The timeit Module

This module can be used to time how long it takes to execute a chunk of code.

Typical chunk = some function of interest.

This is called benchmarking.

Benchmarking

Let's benchmark LinSearch (x,a) and BinSearch (x,a).

Compare how long it takes when len(a) equals 1000, 10000, 100000, and 1000000.

Our intuition tells us that as len(a) increases, BinSearch will be dramatically faster.

BinSearch vs LinSearch

n	tBin	tLin	tLinW
1000	0.0007	0.0064	0.0119
10000	0.0009	0.0668	0.1203
100000	0.0011	0.8296	1.2082
1000000	0.0015	17.7388	13.9341

```
tBin = time for BinSearch
tLin = time for LinSearch (for loop version)
tLinW = time for LinSearch (while-loop version)
```

BinSearch vs LinSearch

n	tLin/tBin	
1000	9	
10000	74	
100000	754	
1000000	7095	

Reporting ratios is more illuminating since we do not really care about the time units in this informal comparison

Using the timeit Module

We show how this module was use to get the results on the previous slides.

Our LinSearch vs BinSearch example is very typical: is one function faster than another?

```
from timeit import *
S = """
   Set-up code
                             Yes, these are doc
                             strings.
// // //
B = """
   Code to Benchmark
// // //
p = 10; m = 100
t = min(Timer(B, setup=S).repeat(p, m))
```

The Set-Up and Bench Codes

```
from random import randint as randi
from ShowSearch import BinSearch
n = 10000
s = [randi(0,10*n) for i in range(n)]
s.sort()
x = s[n/2]
```

```
k=BinSearch(x,s)
```

The set-up code is run once. It is not timed.

It just sets up the code to be timed.

```
from timeit import *
     11 11 11
   Set-up code
                                m times.
// // //
     11 11 11
   Code to Benchmark
// // //
p = 10; m = 100
t = min(Timer(B, setup=S).repeat(p, m))
```

An "experiment" consists of running the blue code

The stopwatch will time how long it takes to do one experiment

Larger values necessary if the blue code executes very quickly

```
from timeit import *
      11 11 11
    Set-up code
                                   Timer returns
// // //
                                   a length-p
      11 11 11
                                   list. Each
                                   element is
    Code to Benchmark
                                   the stopwatch
                                   time for 1
// // //
                                   experiment
p = 10; m = 100
t = min(Timer(B, setup=S).repeat(p, m))
```

This helps control for other stuff that may be running on your computer.

```
from timeit import *
      11 11 11
    Set-up code
                                    In general, it is
// // //
                                    best to take
      11 11 11
                                    the mininum as
                                    the most reliable.
    Code to Benchmark
                                     The benchmark
                                    time is assigned
// // //
                                    to t
p = 10; m =
t = min(Timer(B, setup=S).repeat(p, m))
```

This helps control for other stuff that may be running on your computer.

Why Benchmarking is Important

Confirms/refutes what our intuition might say about efficiency.

Makes us sensitive to the various issues that affect efficiency.

Steers us away from simplistic comparisons of different methods that can be used on the same problem.

MergeSort

Binary Search is an example of a "divide and conquer" approach top probkem solving.

A method for sorting a list that features this strategy is MergeSort

Motivation

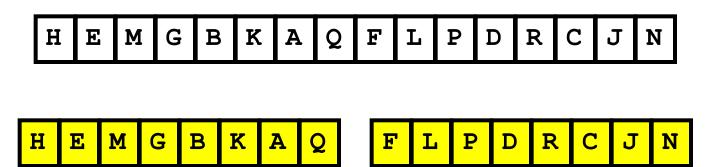
You are asked to sort a list but you have two "helpers": H1 and H2.

Idea:

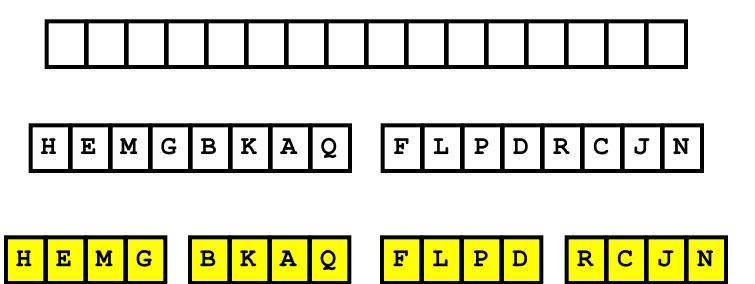
- 1. Split the list in half and have each helper sort one of the halves.
- 2. Then merge the two sorted lists into a single larger list.

This idea can be repeated if H1 has two helpers and H2 has two helpers.

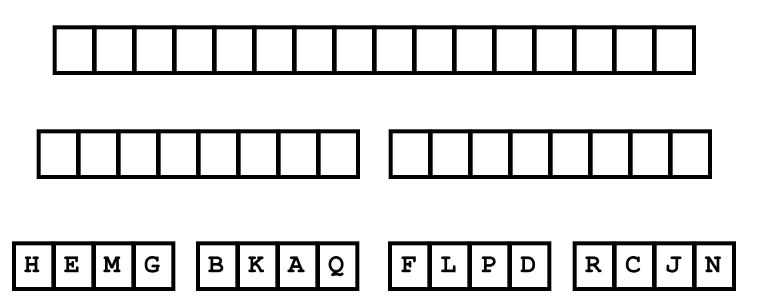
Subdivide the Sorting Task



Subdivide Again

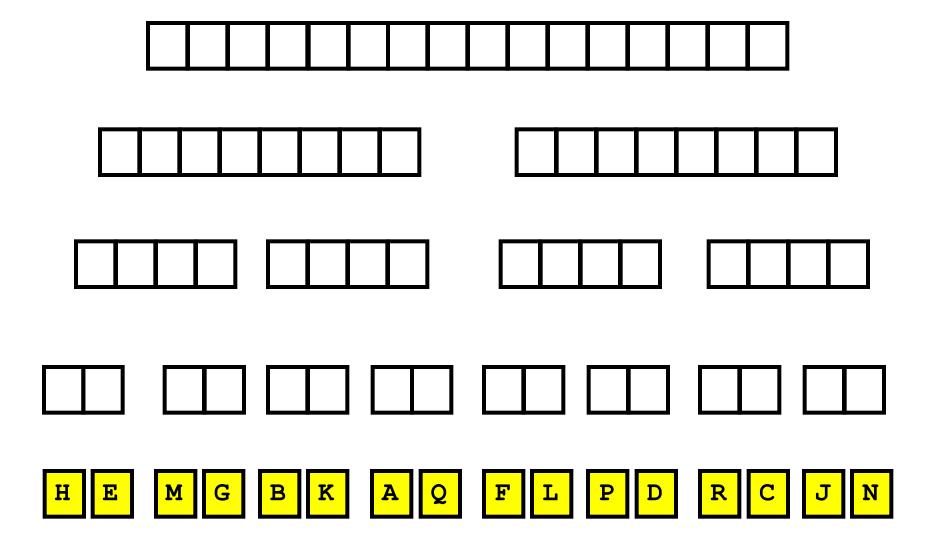


And Again

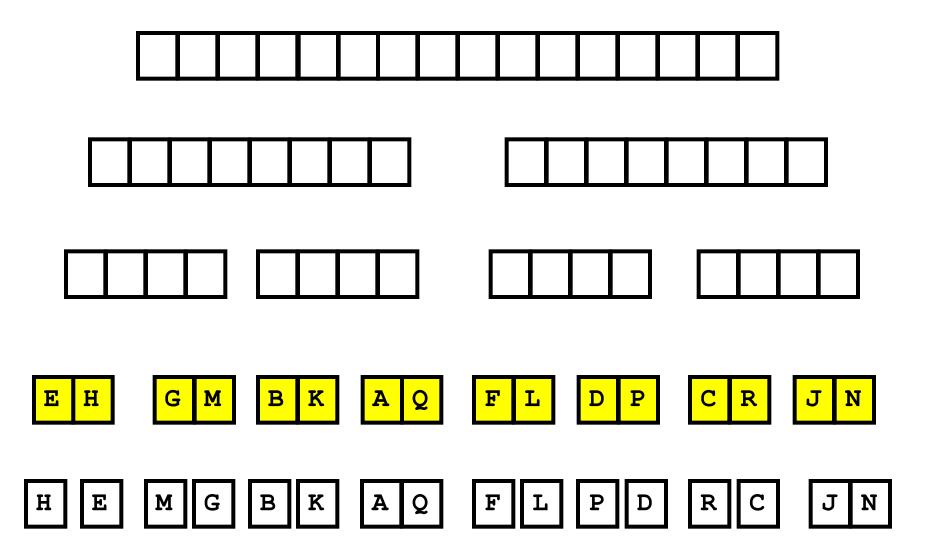


H E M G B K A Q F L P D R C J N

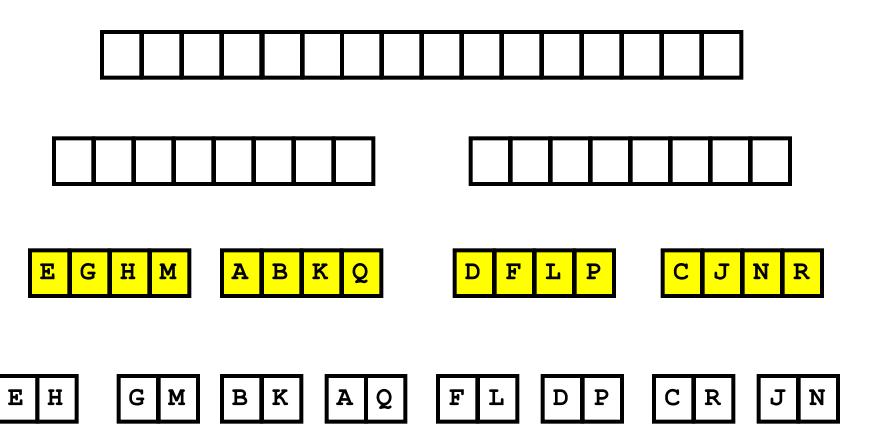
And One Last Time



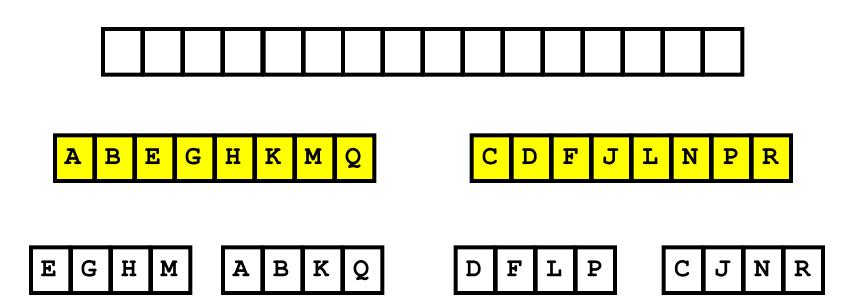
Now Merge



And Merge Again



And Again



And One Last Time







Done!

A B C D E F G H J K L M N P Q R

Done!

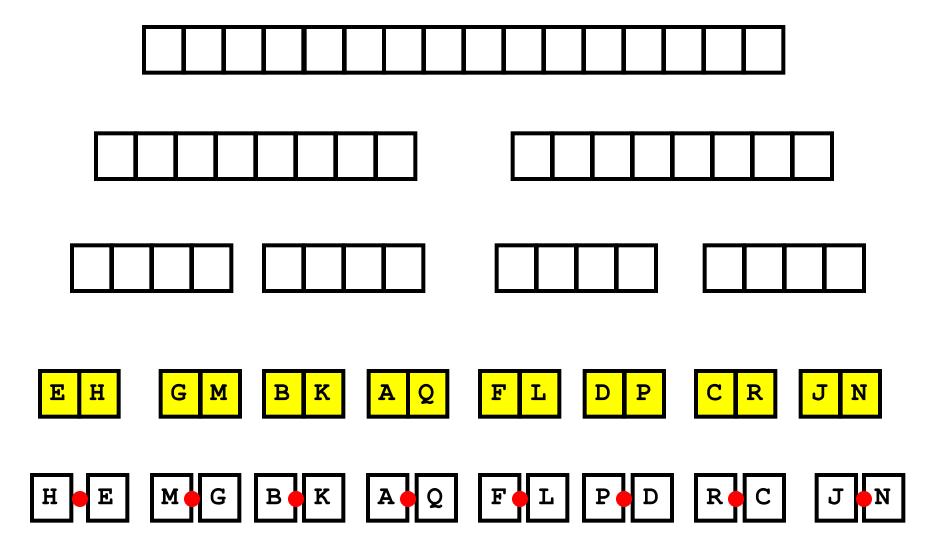


Let's write a function to do this making use of

```
def Merge(x,y):
    """ Returns a float list that is the
    merge of sorted lists x and y.

PreC: x and y are lists of floats
    that are sorted from small to big.
"""
```

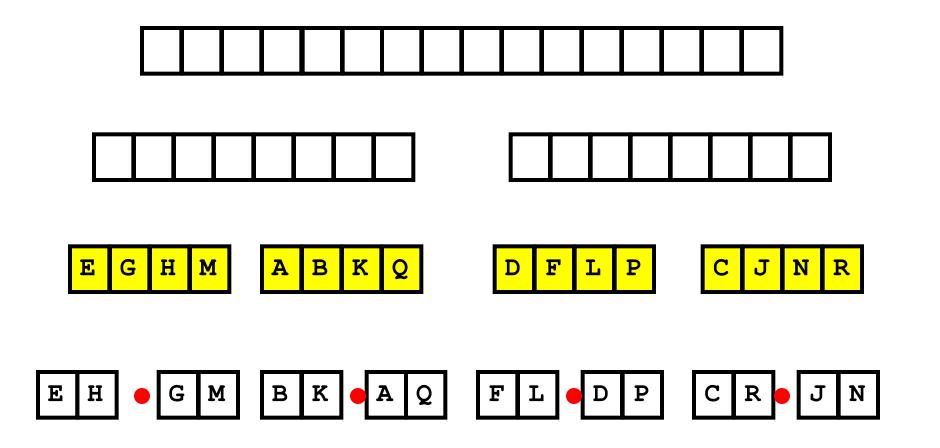
8 Merges Producing length-2 lists



Handcoding the n = 16 case

```
A0 = Merge(a[0],a[1])
A1 = Merge(a[2],a[3])
A2 = Merge(a[4], a[5])
A3 = Merge(a[6], a[7])
A4 = Merge(a[8], a[9])
A5 = Merge(a[10], a[11])
A6 = Merge(a[12], a[13])
A7 = Merge(a[14], a[15])
```

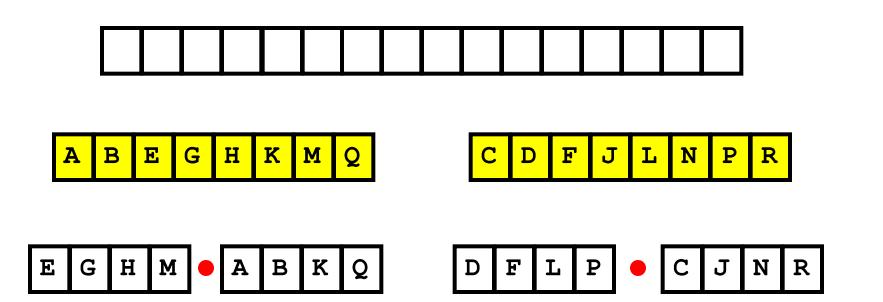
4 Merges Producing Length-4 lists



Handcoding the n = 16 case

```
B0 = Merge(A0,A1)
B1 = Merge(A2,A3)
B2 = Merge(A4,A5)
B3 = Merge(A6,A7)
```

2 Merges Producing Length-8 Lists



Handcoding the n = 16 case

```
C0 = Merge(B0,B1)
C1 = Merge(B2,B3)
```

1 Merge Producing a Length-16 List





All Done!

$$D0 = Merge(C0,C1)$$

For general n, it can be handled using recursion.

Recursive Merge Sort

```
def MergeSort(a):
     n = length(a)
     if n==1:
                               A function
          return a
                               can call
     else:
                               Itself!
         m = n/2
         u0 = list(a[:m])
         u1 = list(a[m:])
         y0 = MergeSort(u0)
         y1 = MergeSort(u1)
          return Merge (y0,y1)
```