CS1110 5 November 2009

Developing array algorithms. Reading: 8.3..8.5

Important point: how we create the invariant, as a picture

Haikus (5-7-5) seen on Japanese computer monitors

Yesterday it worked.
Today it is not working.
Windows is like that.

Serious error.
All shortcuts have disappeared.
Screen. Mind. Both are blank.

A crash reduces
Your expensive computer
To a simple stone.

The Web site you seek
Cannot be located, but
Countless more exist.

Three things are certain:
Death, taxes, and lost data.
Guess which has occurred?

Chaos reigns within.
Reflect, repent, and reboot.
Order shall return.

1

P2 review session in Hollister B14, Sunday 1-3

4. for loops. We may give you a problem that requires you to write a loop (with initialization) that processes a range of integers. You should be able to write a postcondition, write the loop header "for (int k ...)", write a loop invariant, and finally develop the various parts of the loops and initialization.
5. while loops. THIS ITEM HAS BEEN REMOVED. YOU ARE NOT RESPONSIBLE FOR IT.
6. Arrays. Everything on Sects 8.1 and 8.2 of the text — these pages discuss the technical details for using arrays in Java and for reasoning about arrays, including the range notation h..k (where h..h-1 is allowed, and indicates the empty range). The prelim will not deal with two-dimensional arrays.

2

Developing algorithms on arrays

We develop several important algorithms on arrays.

With each, we specify the algorithm by giving its precondition and postcondition as pictures.

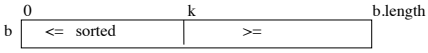
Then, draw the invariant by drawing another picture that "generalizes" the precondition and postcondition, since the invariant is true at the beginning and at the end.

Four loopy questions —memorize them:

1. How does loop start (how to make the invariant true)?
2. When does it stop (when is the postcondition true)?
3. How does repetend make progress toward termination?
4. How does repetend keep the invariant true?

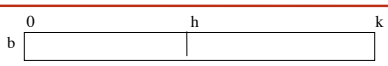
2

Horizontal notation for arrays, strings, Vectors



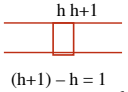
Example of an assertion about an array b. It asserts that:

1. b[0..k–1] is sorted (i.e. its values are in ascending order)
2. Everything in b[0..k–1] is ≤ everything in b[k..b.length–1]



Given the index h of the First element of a segment and the index k of the element that Follows the segment, the number of values in the segment is k – h.

b[h .. k – 1] has k – h elements in it.

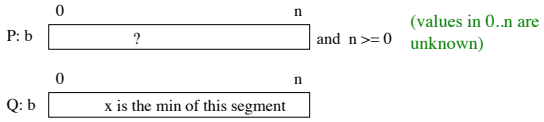


(h+1) – h = 1

4 3

Invariant as picture: Combining pre- and post-condition

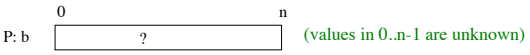
Finding the minimum of an array. Given array b satisfying precondition P, store a value in x to truthify postcondition Q:



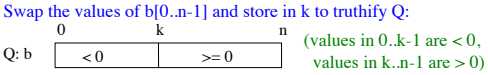
3

The invariant as picture: Combining pre- and post-condition

Put negative values before nonnegative ones. given precondition P:



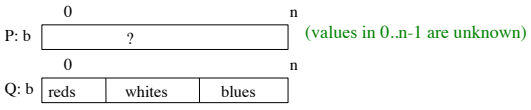
Swap the values of b[0..n-1] and store in k to truthify Q:



4

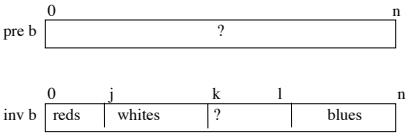
The invariant as picture: Combining pre- and post-condition

Dutch national flag. Swap values of 0..n-1 to put the reds first, then the whites, then the blues. That is, given precondition P, swap value of b [0.n] to truthify postcondition Q:



5

How to make invariant look like initial condition

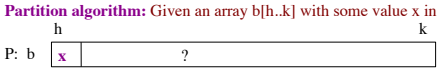


1. Make red, white, blue section empty: use formulas for no. of values in these sections, set j, k , l so that they have 0 elements.

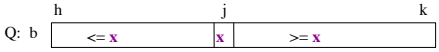
2. Compare precondition with invariant. E.g. in precondition, 0 marks first unknown. In invariant, k marks first unknown. Therefore, k and 0 must be the same.

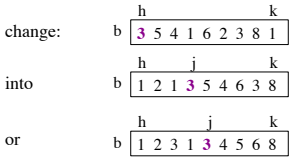
4

Partition algorithm: Given an array b[h..k] with some value x in b[h]:



Swap elements of b[h..k] and store in j to truthify P:





x is called the pivot value.

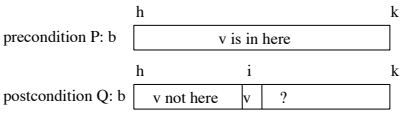
x is not a program variable; x just denotes the value initially in b[h].

6

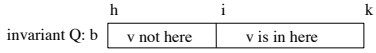
Linear search (for value known to be in array)

Vague spec.: Find first occurrence of v in b[h..k-1], which is known to be there.
Better spec.: Store an integer in i to truthify postcondition Q:

Q: 1. v is not in b[h..i-1]
2. v = b[i]



Can't simply combine P and Q because position of v not known initially. But we can just delete value v from Q:



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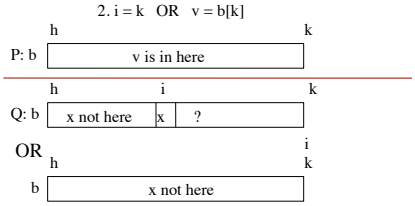
Linear search

Vague spec.: Find first occurrence of v in b[h..k-1].
Better spec.: Store an integer in i to truthify postcondition Q:

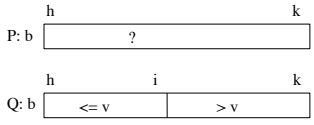
Q: 1. v is not in b[h..i-1]
2. i = k OR v = b[k]



8

Binary search: Vague spec: Look for v in **sorted** array segment b[h..k].
Better spec:
Precondition P: b[h..k] is sorted (in ascending order).
Store in i to truthify:
Postcondition Q: b[h..i] <= v and v < b[i+1..k]

Below, the array is in non-descending order:



Called binary search because each iteration of the loop cuts the array segment still to be processed in half

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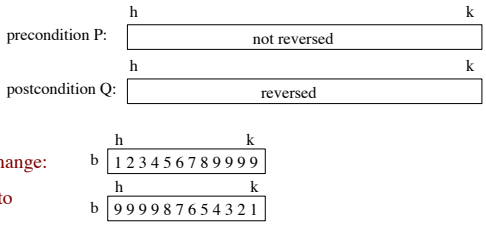
Reversal: Reverse the elements of array segment b[h..k].



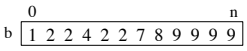
Change:

into

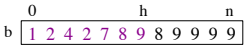
10

Remove adjacent duplicates

change:



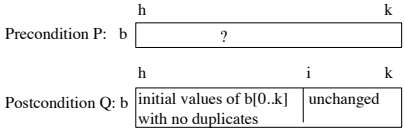
into



don't care what is in b[k+1..n]

Truthify:

b[0..h] = initial values in b[0..n] but with adj dups removed



11

Check whether two arrays are equal

```
/** = "b and c are equal" (both null or both contain
       arrays whose elements are the same) */
public static boolean equals(int[] b, int[] c) {




}
```

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