

Garbage collection; lightstick orientation

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Administrivia

- Assignment 3 is due Friday
- Prelim on 10/11
 - Review on Wed 10/10 in section (by Gurmeet)
 - Coverage through Wed 10/10



Manual storage reclamation

- Programmers always ask for a block of memory of a certain size
 - In C, explicitly declare when it is free
- Desirable but complex invariants:
 - Everything should be freed when it is no longer going to be used
 - And, it should be freed exactly once!
 - Also want to minimize fragmentation
- Any serious C programmer writes their own memory allocation system!

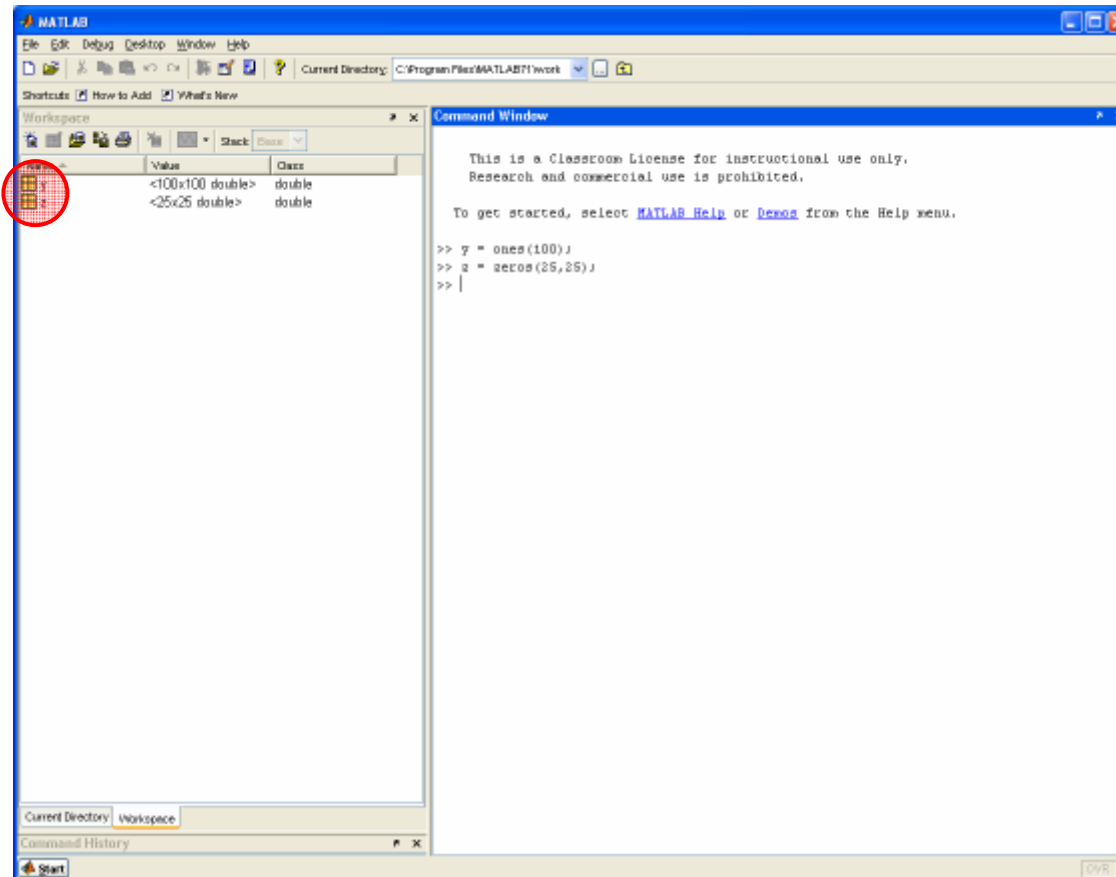


Automatic storage reclamation

- First challenge: figure out what memory locations are in use by the programmer
 - They are the ones that the programmer can get to (i.e., in a linked list currently in use)
 - Anything that has a name the programmer can get to (this is explicit in Matlab)
 - Start with something that has a name, then chase pointers as far as you can go
 - Some programs can access everything...
 - These languages don't allow such programs!



Root set in Matlab



How to garbage collect?

- Need to distinguish a value from a pointer
 - Garbage collector doesn't know anything about the programmer's data representation
- Usual solution: use highest-order bit
 - The pointer to location 3 is represented as $2^n + 3$, where n is usually 31 (word size - 1)
 - Some computers actually had hardware support to tell pointers from values
 - Nice idea that died out in the 90's



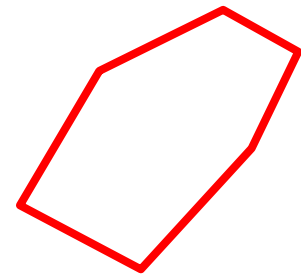
Simple algorithm: mark-sweep

- Chase the pointers from the root set until you are done
 - Modifying everything as you go (“mark”)
 - Typically, set a bit
- Everything not marked is garbage, and can go back on the free list
- Problem: typically a lot more garbage than non-garbage, and need to examine all the garbage with mark-sweep
 - Solutions exist, but are somewhat complex



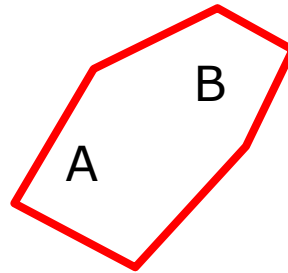
Back to lightstick shape

- We have a blob of red pixels
 - Now what?
- We can easily check which pixels are at the edge of the region
 - How does this help us figure out the shape?
- Easy solution: build a polygon
 - With a small number of vertices



Convex polygons

- A polygon is **convex** if, for any two points A, B in the polygon, all the points between A and B are also in the polygon



- Corresponds to your intuitive notion of convex
 - More or less...
- Convex things turn out to matter a lot!



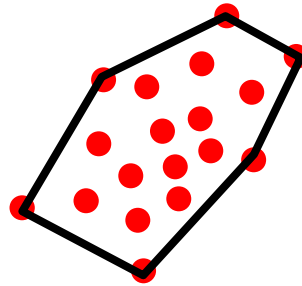
Testing convexity

- How can we test if a polygon P is convex?
 - The obvious way is pretty dumb
 - As is so often the case (in 100R, and in reality)
- Consider the smallest convex polygon containing P
 - Called the **CONVEX HULL**
 - What is the convex hull if P is convex?



Convex hull of point sets

- What is the smallest convex shape containing a bunch of points?
 - I.e., our red pixels?



- State of the art in lightstick orientation!
 - We still need to compute polygon orientation



Computing convex hull

- How do we actually compute the convex hull of a set of points?
 - Note that it is easy to identify at least a few points that are part of the convex hull
 - Which ones are they?
- More precisely, the bounding box is related to the convex hull
 - It touches at 4 points
 - Can it touch at more?



Gift-wrapping algorithm

- Convex hull: smallest convex polygon containing the points
- Gift-wrapping:
 - Start at lowest point
 - Find new point such that all the other points lie to the left of the line to it
 - I.e., the largest angle
 - Repeat

