CS 4110

Programming Languages & Logics

Lecture 11
Weakest Preconditions

21 September 2016

Announcements

- Homework #3 due tonight at 11:59pm
- Homework #4 out now

Review: Decorating Programs

$\{\textbf{true}\}$

$$
\begin{array}{l}
\mathsf{x} := \mathsf{m}; \\
\mathsf{y} := 0; \\
\textbf{while } (\mathsf{n} < \mathsf{x}) \textbf{ do } (\\
\quad \mathsf{x} := \mathsf{x} - \mathsf{n}; \\
\quad \mathsf{y} := \mathsf{y} + 1 \\
)
\end{array}
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$\{ \qquad\qquad\qquad\qquad \}$

Review: Decorating Programs

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\{\mathsf{n} \times \mathsf{y} + \mathsf{x} = \mathsf{m}\}
\end{array}
$$

In other words, the program divides m by n, so y is the quotient and x is the remainder.

Generating Preconditions

To fill in a precondition:

$$\{\quad\} \; c \; \{Q\}$$

there are many possible preconditions—and some are more useful than others.

Weakest Preconditions

Intuition: The weakest liberal precondition for c and Q is the weakest assertion P such that $\{P\}\ c\ \{Q\}$ is valid.

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More formally...

Definition (Weakest Liberal Precondition)

P is a weakest liberal precondition of c and Q written $wlp(c, Q)$ if:

$$\forall \sigma, I.\ \sigma \vDash_I P \iff (\mathcal{C}[\![c]\!]\ \sigma) \text{ undefined } \vee (\mathcal{C}[\![c]\!]\sigma) \vDash_I Q$$

Weakest Preconditions

$$wlp(\mathbf{skip}, P) \quad = \quad P$$

Weakest Preconditions

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\begin{aligned}
wlp(\textbf{skip}, P) &= P \\
wlp(\mathsf{x} := a, P) &= P[a/\mathsf{x}]
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wlp(\textbf{skip}, P) & = & P \\
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wlp(\textbf{if } b \textbf{ then } c_1 \textbf{ else } c_2, P) & = & (b \implies wlp(c_1, P)) \land \\
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wlp(\textbf{while } b \textbf{ do } c, P) & = & \bigwedge_i F_i(P)
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$$

where

$$
\begin{array}{rcl}
F_0(P) & = & \textbf{true} \\
F_{i+1}(P) & = & (\neg b \implies P) \wedge (b \implies wlp(c, F_i(P)))
\end{array}
$$

Applications of Weakest Preconditions

Failing fast: avoid wasting work on bad inputs.

p := getPacket();
processPacket(p);
assert P_{safe}

Applications of Weakest Preconditions

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\begin{array}{l}
p := \mathrm{getPacket}(); \\
\mathrm{processPacket}(p); \\
\{P_{\mathrm{safe}}\}
\end{array}
$$

Applications of Weakest Preconditions

Failing fast: avoid wasting work on bad inputs.

p := getPacket();
$\{P_{\text{filter}}(p)\}$;
processPacket(p);
$\{P_{\text{safe}}\}$

Applications of Weakest Preconditions

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p := getPacket();
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Applications of Weakest Preconditions

Failing fast: avoid wasting work on bad inputs.

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assert $P_{\text{filter}}(p)$;
processPacket(p);

P_{filter} should be the *weakest* precondition to avoid ruling out legitimate inputs.

David Brumley, Hao Wang, Somesh Jha, and Dawn Song. "Creating Vulnerability Signatures Using Weakest Preconditions." In *Computer Security Foundations* (CSF), 2007.

Properties of Weakest Preconditions

Lemma (Correctness of Weakest Preconditions)

$$
\begin{aligned}
&\forall c \in \textbf{Com}, Q \in \textbf{Assn}. \\
&\qquad \models \{wlp(c, Q)\}\, c\, \{Q\} \;\; \textit{and} \\
&\qquad \forall R \in \textbf{Assn}.\;\; \models \{R\}\, c\, \{Q\}\; \textit{implies}\; (R \implies wlp(c, Q))
\end{aligned}
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Properties of Weakest Preconditions

Lemma (Correctness of Weakest Preconditions)

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Lemma (Provability of Weakest Preconditions)

$$\forall c \in \textbf{Com}, Q \in \textbf{Assn}.\; \vdash \{wlp(c, Q)\}\, c\, \{Q\}$$

Soundness and Completeness

Soundness: If we can prove it, then it’s actually true.

Completeness: If it’s true, then a proof exists.

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Definition (Soundness)

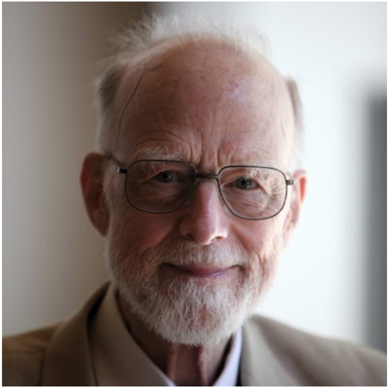
If $\vdash \{P\}\, c\, \{Q\}$ then $\models \{P\}\, c\, \{Q\}$.

Completeness: If it's true, then a proof exists.

Definition (Completeness)

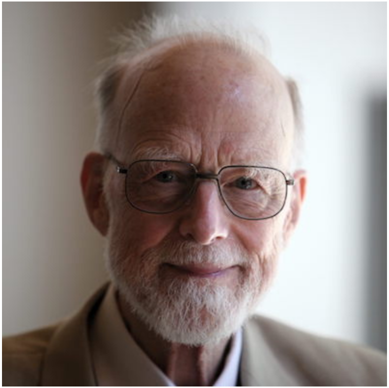
If $\models \{P\}\, c\, \{Q\}$ then $\vdash \{P\}\, c\, \{Q\}$.

vs.



Kurt Gödel

vs.



Sir Tony Hoare

Relative Completeness

Theorem (Cook (1974))

$\forall P, Q \in \textbf{Assn}, c \in \textbf{Com}.\ \vDash \{P\}\, c\, \{Q\}$ *implies* $\vdash \{P\}\, c\, \{Q\}$.

Relative Completeness

Theorem (Cook (1974))

$\forall P, Q \in \textbf{Assn}, c \in \textbf{Com}.\ \vDash \{P\}\ c\ \{Q\}$ *implies* $\vdash \{P\}\ c\ \{Q\}$.

Proof Sketch.

Let $\{P\}\ c\ \{Q\}$ be a valid partial correctness specification.

By the first Lemma we have $\vDash P \implies wlp(c, Q)$.

By the second Lemma we have $\vdash \{wlp(c, Q)\}\ c\ \{Q\}$.

We conclude $\vdash \{P\}\ c\ \{Q\}$ using the CONSEQUENCE rule. $\square$