



CS 3110

Lists and Pattern Matching

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Today's music: "Blank Space" by Taylor Swift

Review

Previously in 3110:

- **Functions:** definition, application, anonymous, higher-order
- **Variables:** bindings, scope

Today:

- **Lists:** OCaml's awesome built-in datatype
- **Pattern matching:** an awesome feature not found in most imperative languages

Lists: An introduction

```
let lst = [1;2;3]
```

```
let empty = []
```

```
let longer = 5::lst
```

```
let another = 5::1::2::3::[]
```

```
let rec sum xs =
```

```
  match xs with
```

```
  | [] -> 0
```

```
  | h::t -> h + sum t
```

```
let six = sum lst
```

```
let zero = sum empty
```

Lists: An introduction

```
let lst = [ "abc"; "def"; "ghi" ]
```

```
let rec concat ss =  
  match ss with  
  | [] -> ""  
  | s::ss' -> s ^ (concat ss')
```

```
let a_i = concat lst
```

Building lists

Syntax:

- `[]` is the empty list
- `e1::e2` prepends element **e1** to list **e2**
- `[e1; e2; ...; en]` is *syntactic sugar* for `e1::e2::...::en::[]`

`[]` is pronounced "nil"

`::` is pronounced "cons" (both from LISP)

Syntactic sugar: redundant kind of syntax that makes program "sweeter" or easier to write

Alan J. Perlis



“Syntactic sugar causes cancer of the semi-colon.”

First recipient of the Turing Award
*for his “influence in the area of advanced programming
techniques and compiler construction”*

1922-1990

Building lists

Evaluation:

- $[]$ is a value
- To evaluate $e1 :: e2$, evaluate $e1$ to a value $v1$, evaluate $e2$ to a (list) value $v2$, and return $v1 :: v2$

Consequence of the above rules:

- To evaluate $[e1; \dots; en]$, evaluate $e1$ to a value $v1$, ..., evaluate en to a value vn , and return $[v1; \dots; vn]$

Building lists

New types:

For any type t , the type $t \text{ list}$ describes lists where all elements have type t

- $[1;2;3] : \text{int list}$
- $[\text{true}] : \text{bool list}$
- $[[1+1;2-3];[3*7]] : \text{int list list}$

Nil:

$[] : 'a \text{ list}$

i.e., empty list has type $t \text{ list}$ for any type t

Cons:

If $e1 : t$ and $e2 : t \text{ list}$ then $e1::e2 : t \text{ list}$

With parens for clarity:

If $e1 : t$ and $e2 : (t \text{ list})$ then $(e1::e2) : (t \text{ list})$

Accessing lists

A list can only be:

- nil, or
- the cons of an element onto another list

Use **pattern matching** to access list in one of those ways:

```
let empty lst =  
  match lst with  
  | []      -> true  
  | h::t    -> false
```

Recursion!

Functions over lists are usually recursive: only way to “get to” all the elements

- What should the answer be for the empty list?
- What should the answer be for a non-empty list?
 - Typically in terms of the answer for the tail of the list



Example list functions

```
let rec sum xs =  
  match xs with  
  | [] -> 0  
  | h::t -> h + sum t
```

```
let rec length xs =  
  match xs with  
  | [] -> 0  
  | h::t -> 1 + length t
```

```
let rec append lst1 lst2 =  
  match lst1 with  
  | [] -> lst2  
  | h::t -> h::(append t lst2)
```

(* append is available as operator @ *)

Lists are immutable

- No way to *mutate* an element of a list
- Instead, build up new lists out of old
e.g., `::` and `@`

Match expressions

Syntax:

```
match e with  
| p1 -> e1  
| p2 -> e2  
| ...  
| pn -> en
```

the **pi** are *patterns*

the first pipe is optional

line breaks are optional

e.g.,

```
let empty lst =  
  match lst with [] -> true | h::t -> false
```

Patterns

Patterns have their own **syntax**

For now, a pattern can be any of these:

- a variable name (e.g., **x**)
- **[]**
- **p1 :: p2**
- an underscore **_**

As we learn more data structures, we'll learn more patterns

Patterns

Patterns **match** values

Intuition of matching is that pattern "looks like" the value, if variables in the pattern are replaced by pieces of the value

- `[]` looks like `[]`
- `h :: t` looks like `2 :: 3`
- `x` looks like `[1 ; 2 ; 3]`
- `_` looks like anything

...we'll make this precise later

Match expressions

```
match e with
| p1 -> e1
| p2 -> e2
| ...
| pn -> en
```

Evaluation:

- Evaluate **e** to a value **v**
- If **p1** matches **v**, then evaluate **e1** to a value **v1** and return **v1**
- Else, if **p2** matches **v**, then evaluate **e2** to a value **v2** and return **v2**
- ...
- Else, if **pn** matches **v**, then evaluate **en** to a value **vn** and return **vn**
- Else, if no patterns match, raise an exception

When evaluating branch expression **ei**, any pattern variables that matched are in scope

Type checker will warn you if you write an *inexhaustive pattern match*

...so you can **prevent exceptions** from being raised at runtime by fixing your code when compiler warns you

Match expressions

```
match e with  
| p1 -> e1  
| p2 -> e2  
| ...  
| pn -> en
```

Type-checking:

If e and $p1 \dots pn$ have type ta
and $e1 \dots en$ have type tb
then entire match expression has type tb

Pattern matching

The pattern `[]` matches the value `[]` and nothing else

```
match [] with  
| []      -> 0  
| h::t    -> 1  (* evaluates to 0 *)
```

```
match [] with  
| h::t    -> 0  
| []      -> 1  (* evaluates to 1 *)
```

Pattern matching

The pattern **$h :: t$** matches any list with at least one element, and binds that element to **h** , and any remaining list to **t**

```
match [1;2;3] with
| []      -> 0
| h::t    -> h  (* evaluates to 1 *)
```

```
match [1;2;3] with
| []      -> [0]
| h::t    -> t  (* evaluates to [2;3] *)
```

A tricky pattern match

What's wrong with this code?

```
let rec drop_val v l =  
  match l with  
  | []      -> []  
  | v::l'   -> drop_val v l'  
  | h::t    -> drop_val v t
```

Hint: compiler warning (as configured in VM)

A tricky pattern match

What's wrong with this code?

```
let rec drop_val v l =  
  match l with  
  | []          -> []  
  | v::l'       -> drop_val v l'  
  | h::t        -> drop_val v t
```

*The **v** in the pattern shadows the argument **v***

A tricky pattern match

```
let rec drop_val v l =  
  match l with  
  | [] -> []  
  | h::t -> let t' = drop_val v t in  
             if h=v then t' else h::t'
```

Deep pattern matching

- Pattern **$a : : []$** matches all lists with exactly one element
- Pattern **$a : : b$** matches all lists with at least one element
- Pattern **$a : : b : : []$** matches all lists with exactly two elements
- Pattern **$a : : b : : c : : d$** matches all lists with at least three elements

Accessing lists, with poor style

- Two library functions that return head and tail
`List.hd, List.tl`
- **Not idiomatic** to apply directly to a list
 - Because they throw exceptions; *you can easily write buggy code*
 - Whereas pattern matching guarantees no exceptions when destructing list; *it's hard to write buggy code!*

Why pattern matching is AWESOME

1. You can't forget a case
(inexhaustive pattern-match warning)
2. You can't duplicate a case
(unused match case warning)
3. You can't get an exception
from forgetting to test the variant
(e.g., `hd []`)
4. Pattern matching leads to elegant, concise,
beautiful code

Functions that immediately match

Instead of

```
let f x =  
  match x with  
  | p1 -> e1  
  | ...  
  | pn -> en
```

can use another piece of syntactic sugar

```
let f = function  
| p1 -> e1  
| ...  
| pn -> en
```

Tail recursion

```
# length [0; 1; ...; 1_000_000];;  
Stack overflow during evaluation  
(looping recursion?).
```

Why?

```
let rec length xs =  
  match xs with  
  | [] -> 0  
  | h::t -> 1 + length t
```

Tail recursion

Solution: When recursive call is the *only thing left to do* in computation, compiler reuses the stack frame. Reduces space from $O(n)$ to $O(1)$.

```
let rec length_plus_n n = function
| [] -> n
| h::t -> length_plus_n (n+1) t

let length_tr = length_plus_n 0
```

Lists (recap)

- **Syntax:** `[] :: [a; b; c]`
- **Semantics:** building with `nil` and `cons`, accessing with pattern matching
- **Idioms:** recursive functions with pattern for `nil` and for `cons`, **function** syntactic sugar, tail recursion
- **Library:** awesome higher-order functions in OCaml standard library (next time)

Upcoming events

- [Monday and Tuesday] Recitations canceled because of Labor Day
- [next Thursday] A1 due

This is awesome.

THIS IS 3110