

**HW 1**

Due: Monday, Sep 9.

Remember that you may (and should!) talk about the problems amongst yourselves, or discuss them with me, providing attribution for any good ideas you might get – but your final write-up should be your own.

**1: A problem of performance.** In addition to dense matrices, MATLAB supports a sparse matrix data structure, in which only the nonzero elements of the matrix are stored. For a variety of square matrix size  $n$  and sparsity levels  $s$  (where  $s$  is the fraction of the entries that are nonzero), compare the speed of dense matrix-vector multiply and sparse matrix-vector multiply. You can use `As = sparse(A)` to make a sparse version of a dense matrix  $A$ . What do you observe about the relative performance of these options?

*Note:* Your performance will vary depending on your machine and your version of MATLAB, of course.

**2: Identity plus low-rank.** Suppose  $u, v \in \mathbb{R}^n$  and  $A = I + uv^T$ . Write a function to compute  $A^k x$  in  $O(n)$  time. Your code should have the signature

```
function [Akx] = p2pow(u, v, k, x)
```

*Note:* If  $c$  is a scalar, you may assume  $c^k$  is computed in  $O(1)$  time.

**3: Structured perturbations.** In this problem, we consider the structure of the groups of orthogonal and symplectic matrices:

1. Suppose  $Q : \mathbb{R} \rightarrow \mathbb{R}^{n \times n}$  is a differentiable matrix-valued function and that  $Q(t)$  is orthogonal for all  $t$ . Show that  $S(t) = Q(t)^T \dot{Q}(t)$  is a skew-symmetric matrix, i.e.  $S(t) = -S(t)^T$ .
2. Suppose  $S : \mathbb{R} \rightarrow \mathbb{R}^{2n \times 2n}$  is a differentiable matrix-valued function and that  $S(t)$  is *symplectic* for all  $t$ , i.e.

$$S^T J S = J, \quad J = \begin{bmatrix} 0 & I \\ -I & 0 \end{bmatrix}.$$

Show that  $M = S^{-1} \dot{S}$  is *Hamiltonian*, i.e.  $JM + M^T J = 0$ .

**4: Norm!** Suppose  $W \in \mathbb{R}^{m \times n}$  is *nonnegative* (all the entries are nonnegative) and *row stochastic* (each row sums to one), and let  $p \geq 1$  be fixed. The function  $\alpha \mapsto \alpha^p$  is *convex*, so if  $y = Wx$ , then

$$|y_i|^p \leq \sum_j w_{ij} |x_j|^p.$$

Argue that

1. If  $u_j = \sum_i w_{ij}$  and  $v_j = |x_j|^p$ , then  $\|y\|_p^p \leq u \cdot v$ .
2. Therefore,  $\|y\|_p^p \leq \|u\|_\infty \|v\|_1$ .
3. Therefore,  $\|y\|_p^p \leq \|W\|_1 \|x\|_p^p$ .
4. Therefore,  $\|W\|_p \leq \|W\|_1^{1/p}$ .

*Note:*

1. Each of the steps above should be explained in a few words or symbols. Don't overthink!
2. The more general statement (which you do not need to prove) is that for any matrix  $A$ ,  $\|A\|_p \leq \|A\|_1^{1/p} \|A\|_\infty^{1/q}$ , where  $1/p + 1/q = 1$ . If you're curious, you can see the whole proof on pages 28–29 of Kato's *Perturbation Theory of Linear Operators*.
3. This question was inspired by a query by Bobby Kleinberg about two years ago, which I take as evidence that (i) this fact is interesting to someone beside me and (ii) it is not obvious!