

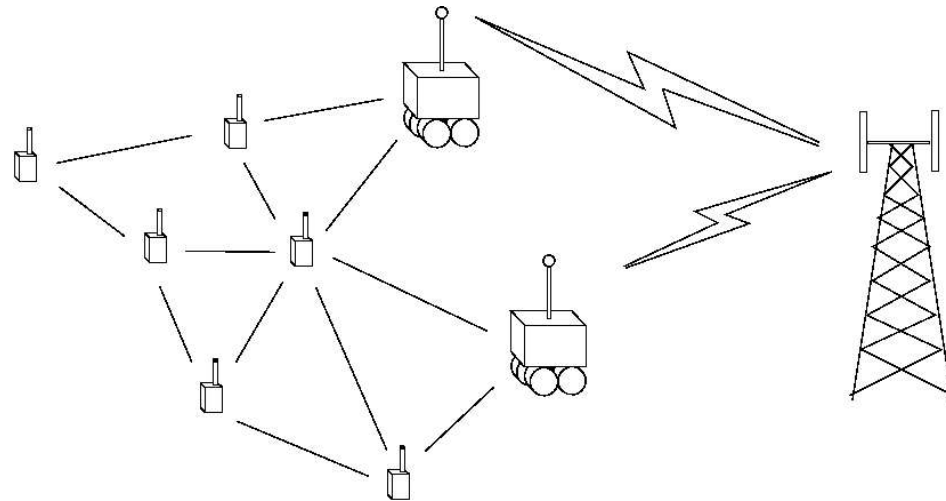
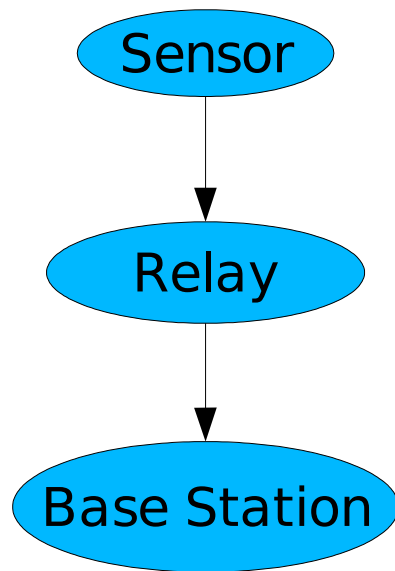
# Maneuverable Relays to Improve Energy Efficiency in Sensor Networks

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# Introduction

- One-sink data flow
- Energy efficiency targeted in many ways
  - Specialized routing algorithms
  - Hierarchical organization
  - Power management techniques
- Our approach
  - Hierarchy with specialized hardware

# Our Approach



- Sensors know their positions
- Relays are maneuverable, can be positioned to desired places, or are mobile

# Goal

- Position relays so that the energy consumption is minimized

$$E_S = \sum_{i=1}^n f_i \cdot d_i \cdot E_U$$

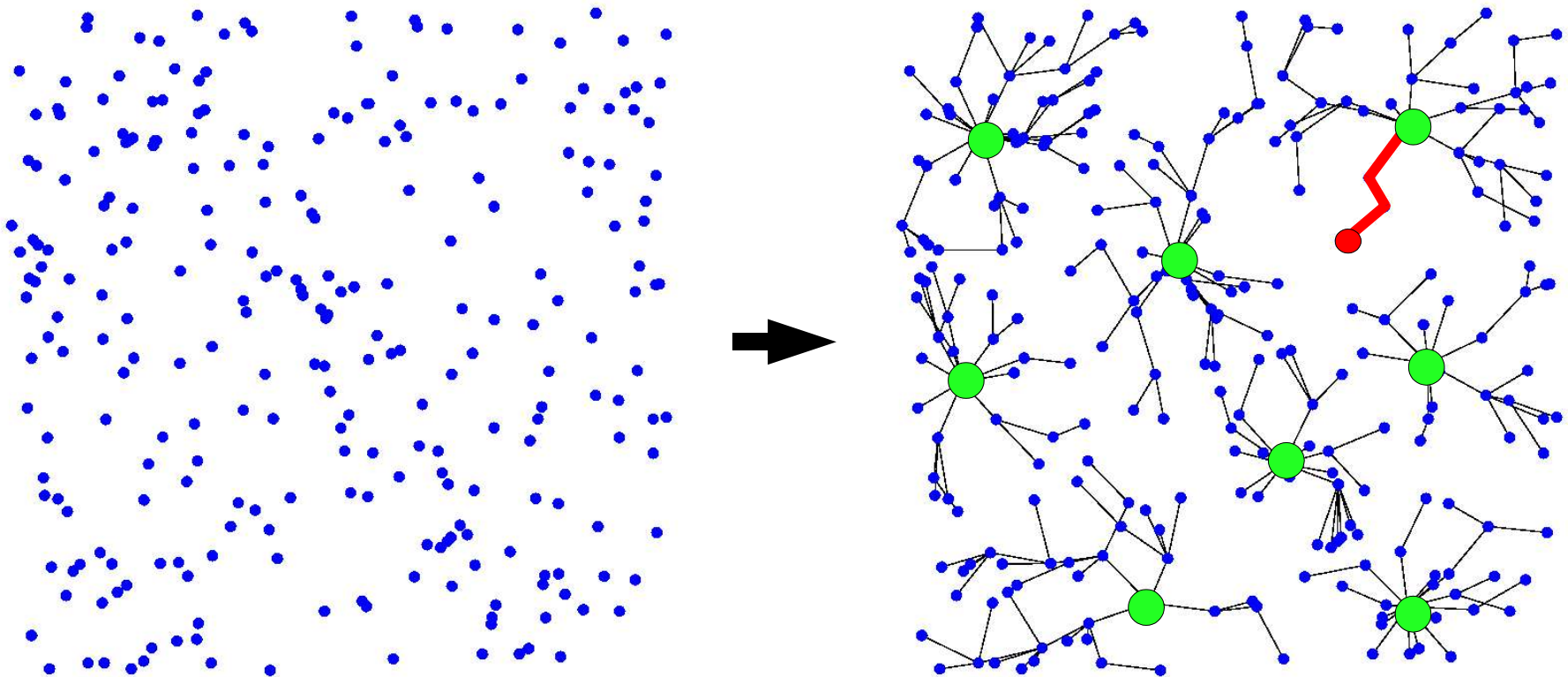
$f$  - # packets generated

$d$  - hop count to relay

$E_U$  - energy per one transmission

- $d$  is influenced by the relays positions, can be used to minimize  $E_S$

# Goal Example



$$E_S = \sum_{i=1}^n f_i \cdot d_i \cdot E_U$$

# Network Usage Algorithm

1. Deploy sensors, compute their positions
2. Gather sensor positions at base station
3. Compute suitable relay positions
4. Deploy relays
5. Relays inform sensors of their positions, sensors know where to send data
6. Data traffic routed along the shortest paths

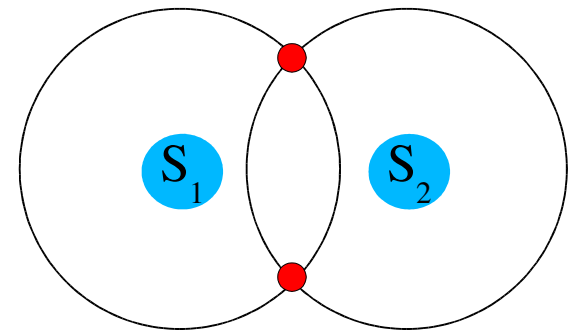
# Relay Positioning Problem

- Definition
  - Given  $n$  positions of sensors in the plane, find  $k$  relay positions such that the resulting energy usage  $E_s$  is minimum.
- Results
  - NP-complete
    - => heuristics needed for large instances
  - Empirical study of various heuristic approaches
    - => cost/quality trade-off

# NP-completeness of Relay Positioning

- Decision version: given energy threshold  $E_T$ , can we place the relays so that  $E_S \leq E_T$  ?

- in NP: optimum solution exists in which relays are placed at intersections of sensor radii circles

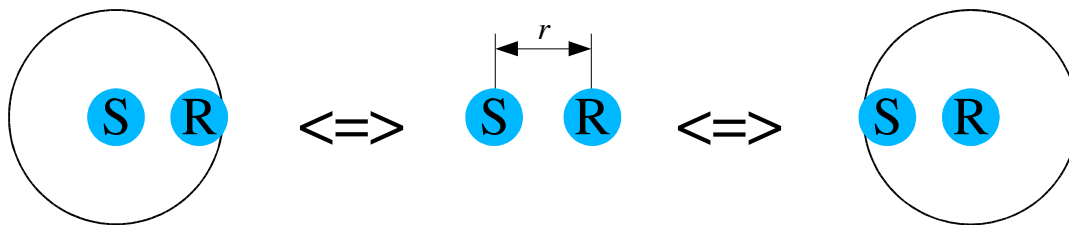


- Reduction from Geometric Covering by Discs
  - Instance: set  $P$  of integer-coordinate points in the plane, positive integers  $r$  and  $K$
  - Question: Can the points in  $P$  be covered by  $K$  discs of radius  $r$ .



# NP-completeness of Relay Positioning

- Reduction:
  - Sensors placed at points from  $P$ , ranges =  $r$
  - $f_i, E_U$  set to 1  $\Rightarrow E_S = \sum d_i$  (hop-count path lengths)
  - $E_T = |P|$  (each sensor allowed only one hop to a relay) $\Rightarrow$  (we can place  $K$  relays so that  $E_S \leq E_T$ ) iff  
(we can cover  $P$  with  $K$  discs with radius  $r$ )

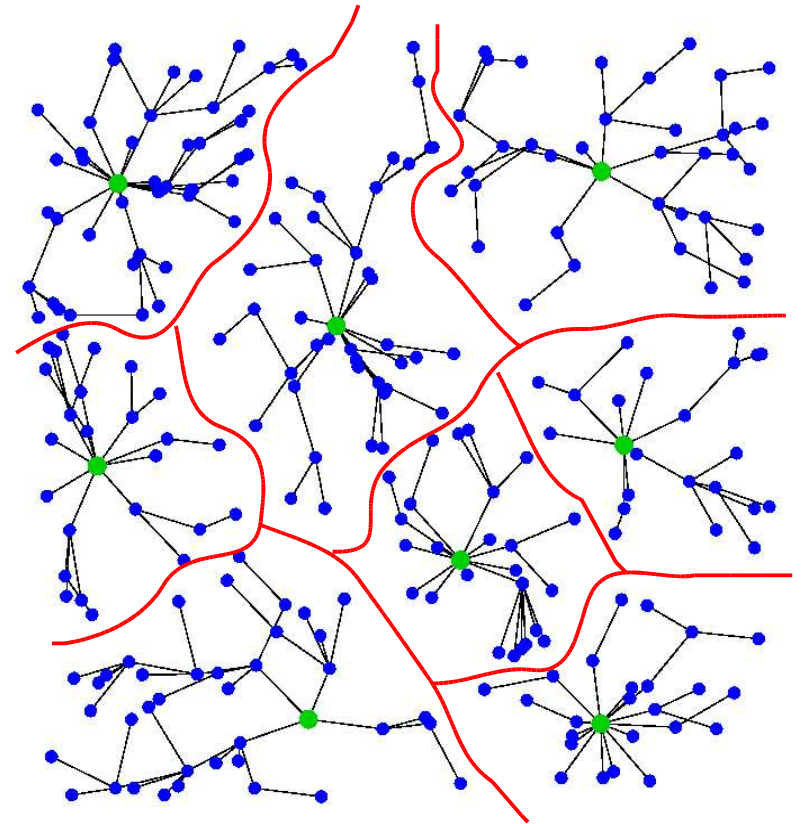


# Heuristic Approaches

- Main assumptions
  - Sensors: stationary
  - Data load: uniform

$$E_S = \sum_{i=1}^n d_i$$

- Positioning relays can be viewed as clustering sensors



# Clustering Algorithms

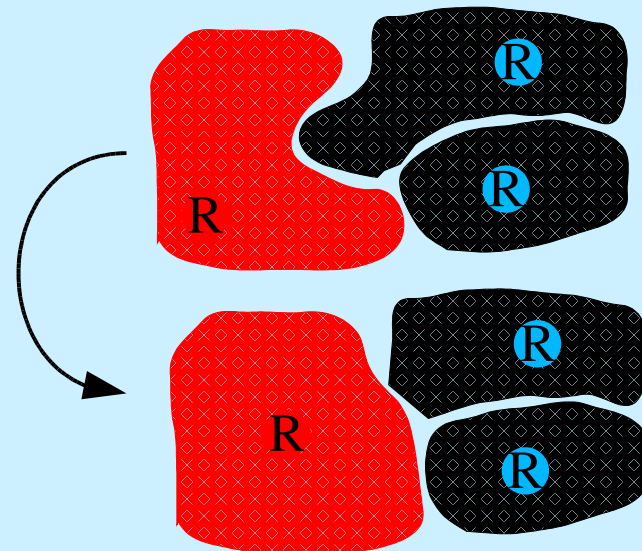
- Locally Improving

- **k-means**
- bottom-up

- Globally Improving

- k-global
- greedy

- Place  $k$  relays randomly
- Repeat: center a relay in its cluster



# Clustering Algorithms

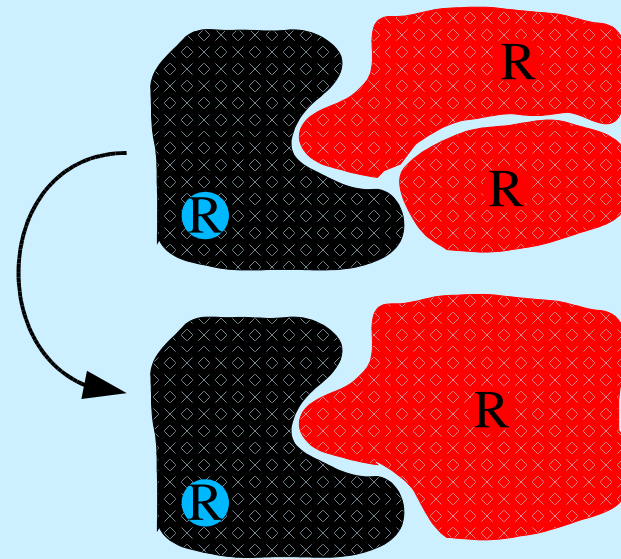
- Locally Improving

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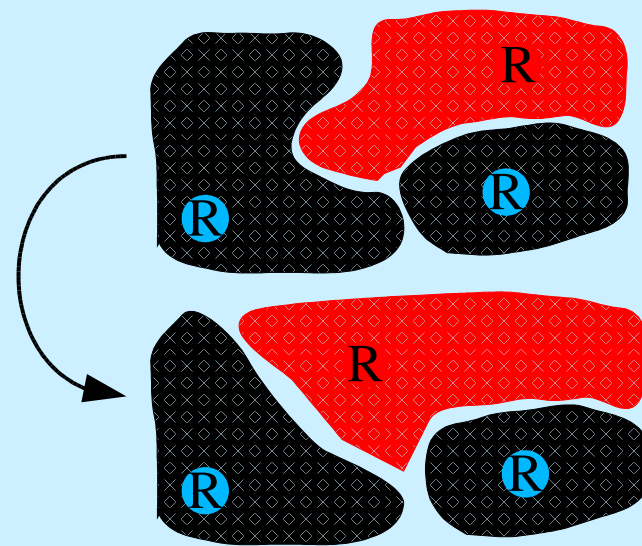
- Each sensor is a cluster
- Repeat: combine two closest clusters



# Clustering Algorithms

- Locally Improving
  - k-means
  - bottom-up
- Globally Improving
  - **k-global**
  - greedy

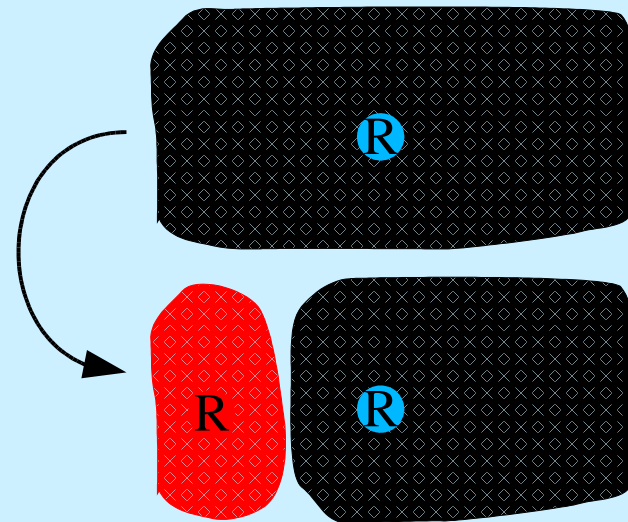
- Place  $k$  relays randomly
- Repeat: place a sensor to minimize overall cost



# Clustering Algorithms

- Locally Improving
  - k-means
  - bottom-up
- Globally Improving
  - k-global
  - greedy

– Place relays one by one, to a globally optimal positions



# Clustering Algorithms

- Distance definitions
  - Euclid distance
    - Easy to compute, may not approximate true cost well
  - Hop-count
    - Corresponds to the true path length, costly to compute
- Possible relay positions set
  - Sensor positions
    - Small  $O(n)$ , but may not contain any optimal solution
  - Plus radii intersection points
    - Large  $O(n^2)$ , but must contain an optimal solution

# Simulation Setup

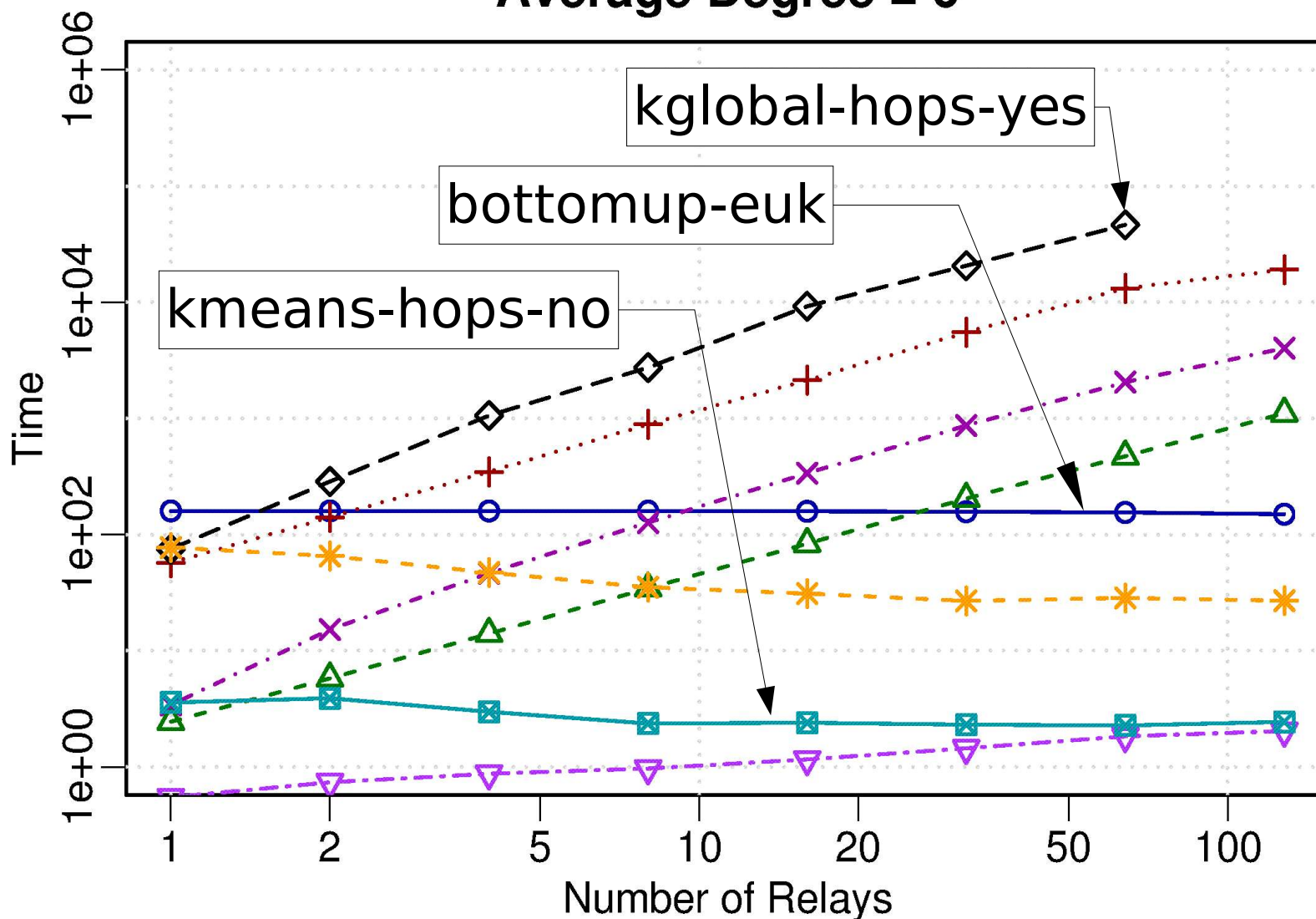
	Local		Global	
	k-means	bottom-up	k-global	greedy
Euclid	X	X		
Hops with	X		X	X
Hops without	X		X	X

- 1024 sensors, 1...128 relays, 6...24 average connectivity degree
- Greedy with intersections is guaranteed to find a solution within  $\sim 1.58$  of the optimum
  - => small differences in solution quality expected



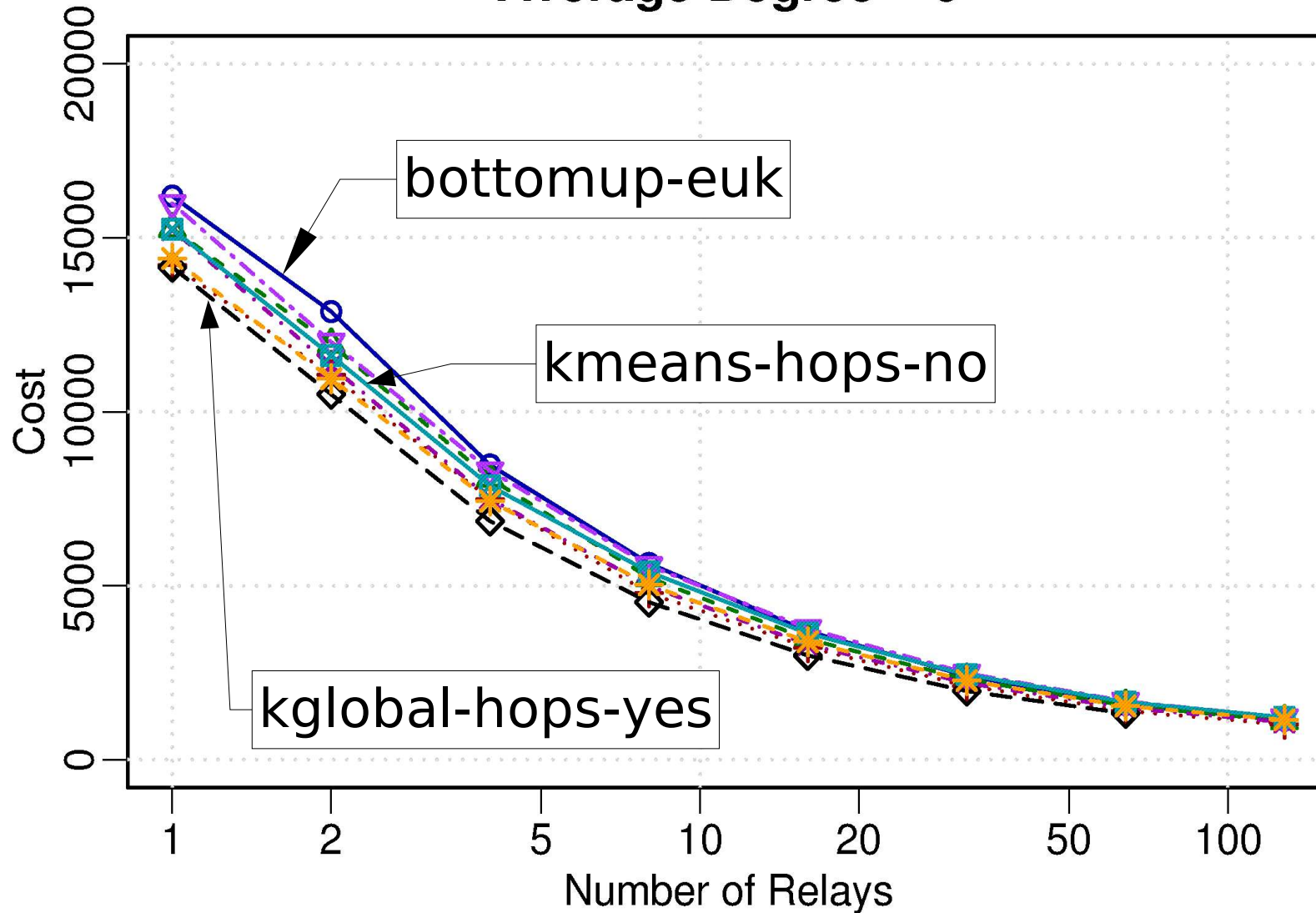
# Simulation Results

Average Degree = 6



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Average Degree = 6



# Conclusions

- Problem of positioning maneuverable relays presented
  - Shown to be NP-complete
  - Performance of various heuristics compared
- Computationally expensive methods bring some improvements
- Future work
  - Minimizing maximum energy use per node
  - Self-positioning relays moving towards regions with high traffic density