User-defined Data Types

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Today’s music: *Pokémon Theme* by Jason Paige
Review

Yesterday:
• New idioms and library functions: Map, fold, and other higher-order functions

Today:
• Ways to define your own data types: records, tuples, variants, options
RECORDS
Record definition

- A **record** contains several named **fields**
- Before you can use a record, must **define** a record type:

  ```
  type time = {hour: int; min: int; ampm: string}
  ```

- To **build** a record:
  - Write a record expression:
    ```
    {hour=10; min=10; ampm="am"}
    ```
  - Order of fields doesn't matter:
    ```
    {min=10; hour=10; ampm="am"} is equivalent
    ```

- To **access** record's field: `r.hour`
Record expressions

• **Syntax:** \{f1 = e1; ...; fn = en\}

• **Evaluation:**
  – If \(e1\) evaluates to \(v1\), and ... \(en\) evaluates to \(vn\)
  – Then \(\{f1 = e1; ...; fn = en\}\) evaluates to \(\{f1 = v1, ...
  \hspace{1cm} ...; fn = vn\}\)
  – Result is a *record value*

• **Type-checking:**
  – If \(e1 : t1\) and \(e2 : t2\) and ... \(en : tn\),
  – and if \(t\) is a defined type of the form \(\{f1 : t1, ...
  \hspace{1cm} ...; fn : tn\}\)
  – then \(\{f1 = e1; ...; fn = en\} : t\)
Record field access

• **Syntax:** `e.f`

• **Evaluation:**
  – If `e` evaluates to `{ f = v, ... }`
  – Then `e.f` evaluates to `v`

• **Type-checking:**
  – If `e : t1`
  – and if `t1` is a defined type of the form `{ f : t2, ... }`
  – then `e.f : t2`
Evaluation notation

We keep writing statements like:
\[ \text{If } e \text{ evaluates to } \{ f = v, \ldots \} \text{ then } e \cdot f \text{ evaluates to } v \]

Let's introduce a shorthand notation:
- Instead of "e evaluates to v"
- Write "e ==> v"

So we can now write:
\[ \text{If } e \text{ ==> } \{ f = v, \ldots \} \text{ then } e \cdot f \text{ ==> } v \]
By name vs. by position

- Fields of record are identified **by name**
  - order we write fields in expression is irrelevant

- Opposite choice: identify **by position**
  - e.g., “Would the student named NN. step forward?”
  - vs. “Would the student in seat \( n \) step forward?”

- You’re accustomed to both:
  - Java object fields accessed by name
  - Java method arguments passed by position
    - (but accessed in method body by name)

- OCaml has something you might not have seen:
  - A kind of data accessed by position
PAIRS AND TUPLES
Pairs

A pair of data: two pieces of data glued together
e.g.,
• \((1,2)\)
• \((\text{true}, \text{"Hello"})\)
• \(([1;2;3], 0.5)\)

We need language constructs to build pairs and to access the pieces...
Pairs: building

• Syntax: \((e_1, e_2)\)

• Evaluation:
  – If \(e_1 \Rightarrow v_1\) and \(e_2 \Rightarrow v_2\)
  – Then \((e_1, e_2) \Rightarrow (v_1, v_2)\)
  – A pair of values is itself a value

• Type-checking:
  – If \(e_1 : t_1\) and \(e_2 : t_2\),
  – then \((e_1, e_2) : t_1*t_2\)
  – A new kind of type, the **product type**
Pairs: accessing

• **Syntax**: \texttt{fst e} and \texttt{snd e}
  
  *Projection functions*

• **Evaluation**:
  – If \( e \implies (v_1, v_2) \)
  – then \( \texttt{fst e} \implies v_1 \)
  – and \( \texttt{snd e} \implies v_2 \)

• **Type-checking**:
  – If \( e : \texttt{ta}*\texttt{tb} \),
  – then \( \texttt{fst e} \) has type \( \texttt{ta} \)
  – and \( \texttt{snd e} \) has type \( \texttt{tb} \)
Tuples

Actually, you can have *tuples* with more than two parts
  – A new feature: a generalization of pairs
  – Syntax, semantics are straightforward, except for projection...

- \((e_1,e_2,\ldots,e_n)\)
- \(t_1 \times t_2 \times \ldots \times t_n\)
- \(\text{fst } e, \text{snd } e, \ldots\)

Instead of generalizing projection functions, use *pattern matching*...

New kind of pattern, the **tuple pattern**: \((p_1, \ldots, p_n)\)
Pattern matching tuples

```
match (1,2,3) with
  | (x,y,z) -> x+y+z

(* ==> 6 *)
```

```
let thrd t =
    match t with
      | (x,y,z) -> z

(* thrd : 'a*'b*'c -> 'c *)
```

Note: we never needed more than one branch in the match expression...
Pattern matching without match

(* OK *)
let thrd t =
  match t with
  | (x,y,z) -> z

(* good *)
let thrd t =
  let (x,y,z) = t in z

(* better *)
let thrd t =
  let (_,_,z) = t in z

(* best *)
let thrd (_,_,z) = z
Extended syntax for let

• Previously we had this syntax:
  – `let x = e1 in e2`
  – `let [rec] f x1 ... xn = e1 in e2`

• Everywhere we had a variable identifier x, we can really use a pattern!
  – `let p = e1 in e2`
  – `let [rec] f p1 ... pn = e1 in e2`

• Old syntax is just a special case of new syntax, since a variable identifier is a pattern
Pattern matching arguments

(* OK *)

let sum_triple t =
  let (x,y,z) = t
  in x+y+z

(* better *)

let sum_triple (x,y,z) = x+y+z

Note how that last version looks syntactically like a function in C/Java!
Question #3

What is the type of this expression?
let (x,y) = snd(“zar”,(“doz”,42))
in (42,y)

A. \{x:string; y:int\}
B. int*int
C. string*int
D. int*string
E. string*(string*int)
Question #3

What is the type of this expression?

```plaintext
let (x, y) = snd("zar", ("doz", 42))
in (42, y)
```

A. `{x:string; y:int}`
B. int*int
C. string*int
D. int*string
E. string*(string*int)
Unit

• Can actually have a tuple ( ) with no components whatsoever
  – Think of it as a degenerate tuple
  – Or, like a Boolean that can only have one value
• “Unit” is
  – a value written ( )
  – and a type written unit
• Might seem dumb now; will be useful later!
Pattern matching records

(* OK *)
let get_hour t =
  match t with
  | {hour=h; min=m; ampm=s} -> h

(* better *)
let get_hour t =
  match t with
  | {hour=h; min=_; ampm=_} -> h

(* better *)
let get_hour t =
  match t with
  | {hour; min; ampm} -> hour

New kind of pattern, the record pattern:
{f1[=p1]; ...; fn[=pn]}
By name vs. by position, again

How to choose between coding \((4, 7, 9)\) and \(\{ f=4 ; g=7 ; h=9 \}\)?

• Tuples are syntactically shorter
• Records are self-documenting
• For many (4? 8? 12?) fields, a record is usually a better choice
VARIANTS
Variant

```haskell
type day = Sun | Mon | Tue | Wed
    | Thu | Fri | Sat

let day_to_int d =
  match d with
    | Sun -> 1
    | Mon -> 2
    | Tue -> 3
    | Wed -> 4
    | Thu -> 5
    | Fri -> 6
    | Sat -> 7
```
Building and accessing variants

Syntax: type t = C1 | ... | Cn

the Ci are called constructors

Evaluation: a constructor is already a value

Type checking: Ci : t

Accessing: use pattern matching; constructor name is a pattern
Pokémon variant

<table>
<thead>
<tr>
<th>DEFENSE →</th>
<th>NOR</th>
<th>FIR</th>
<th>WAT</th>
<th>E</th>
</tr>
</thead>
<tbody>
<tr>
<td>NORMAL</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>FIRE</td>
<td>½</td>
<td>½</td>
<td></td>
<td></td>
</tr>
<tr>
<td>WATER</td>
<td>2</td>
<td>½</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>
**Pokémon variant**

```ocaml
type ptype = TNormal | TFire | TWater

type peff = ENormal | ENotVery | ESuper

let eff_to_float = function
  | ENormal    -> 1.0
  | ENotVery   -> 0.5
  | ESuper     -> 2.0

let eff_att_vs_def : ptype*ptype -> peff = function
  | (TFire,TFire)      -> ENotVery
  | (TWater,TWater)    -> ENotVery
  | (TFire,TWater)     -> ENotVery
  | (TWater,TFire)     -> ESuper
  | _                   -> ENormal
```
Argument order: records

If you are worried about clients of function forgetting which order to pass arguments in tuple, use a record:

```ocaml
type att_def = {att:ptype; def:ptype}

let eff_att_vs_def : att_def -> peff = function
    | {att=TFire;def=TFire}    -> ENotVery
    | {att=TWater;def=TWater}  -> ENotVery
    | {att=TFire;def=TWater}   -> ENotVery
    | {att=TWater;def=TFire}   -> ESuper
    | _                       -> ENormal
```
Argument order: labeled arguments

Or (though not quite as good) use labeled arguments:

```ocaml
def eff_att_vs_def att def = 
match (att, def) with 
| (TFire,TFire) -> ENotVery 
| (TWater,TWater) -> ENotVery 
| (TFire,TWater) -> ENotVery 
| (TWater,TFire) -> ESuper 
| _ -> ENormal 

def super = eff_att_vs_def att:TWater def:TFire 
def super = eff_att_vs_def def:TFire att:TWater 
def notvery = eff_att_vs_def TFire TWater
```
## Variants vs. records vs. tuples

<table>
<thead>
<tr>
<th></th>
<th>Define</th>
<th>Build/construct</th>
<th>Access/destroy</th>
</tr>
</thead>
<tbody>
<tr>
<td>Variant</td>
<td><code>type</code></td>
<td>Constructor name</td>
<td>Pattern matching</td>
</tr>
<tr>
<td>Record</td>
<td><code>type</code></td>
<td>Record expression with <code>{ ... }</code></td>
<td>Pattern matching OR field selection with dot operator .</td>
</tr>
<tr>
<td>Tuple</td>
<td>N/A</td>
<td>Tuple expression with <code>( ... )</code></td>
<td>Pattern matching OR <code>fst</code> or <code>snd</code></td>
</tr>
</tbody>
</table>

- **Variants**: *one-of types* aka *sum types*
- **Records, tuples**: *each-of types* aka *product types*
Question

Which of the following would be better represented with records rather than variants?

A. Coins, which can be pennies, nickels, dimes, or quarters
B. Students, who have names and id numbers
C. A plated dessert, which has a sauce, a creamy component, and a crunchy component
D. A and C
E. B and C
Question

Which of the following would be better represented with records rather than datatypes?

A. Coins, which can be pennies, nickels, dimes, or quarters
B. Students, who have names and NetIDs
C. A plated dessert, which has a sauce, a creamy component, and a crunchy component
D. A and C
E. B and C
OPTIONS
What is max of empty list?

```ocaml
let rec max_list = function
  | [] -> ???
  | h::t -> max h (max_list t)
```

How to fill in the ???

- `min_int` would be a reasonable choice...
- or could raise an exception...
- in Java, might return `null`...
- but OCaml gives us another option!
Options

Options:
• \( \text{t option} \) is a type for any type \( \text{t} \)  
  (much like \( \text{t list} \) is a type for any type \( \text{t} \))

Building and Type Checking and Evaluation:
• \( \text{None} \) has type '\( \text{a option} \)'
  – much like \( \text{[]} \) has type '\( \text{a list} \)'
  – \( \text{None} \) is a value
• \( \text{Some e : t option if e : t} \)
  – much like \( \text{e :: []} \) has type \( \text{t list if e : t} \)
  – If \( \text{e ===> v} \) then \( \text{Some e ===> Some v} \)

Accessing: \( \text{match e with} \)
  \( \text{None} \rightarrow \ldots \)
  \( \mid \text{Some x} \rightarrow \ldots \)
Again: What is max of empty list?

```ocaml
let rec max_list = function
  | [] -> None
  | h::t -> match max_list t with
            | None    -> Some h
            | Some x  -> Some (max h x)

(* max_list : 'a list -> 'a option *)
```

Very stylish!
...no possibility of exceptions
...no chance of programmer ignoring a “null return”
Recap: User-defined data types

• Records
• Tuples (pairs, unit)
• Variants
• Options