

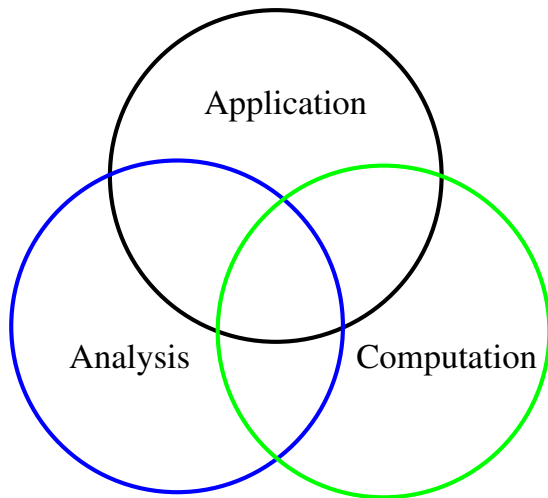
# From Network Tomography to Power Networks?

David Bindel

Department of Computer Science  
Cornell University

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# The Computational Science & Engineering Picture

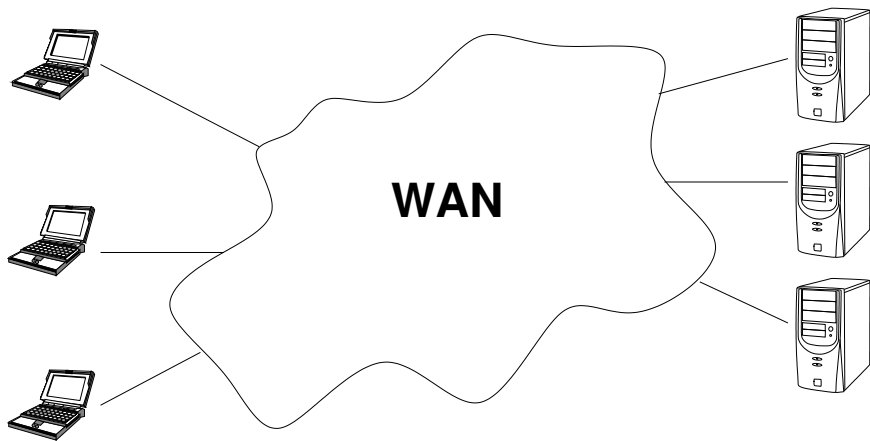


- Numerical methods
  - Structured and nonlinear eigenvalue solvers
  - Continuation for eigenvalue problems
  - Model reduction
  - Fast solvers for structured systems
  - Finite element methods
  - Network models
- Applications
- Software and computational infrastructure

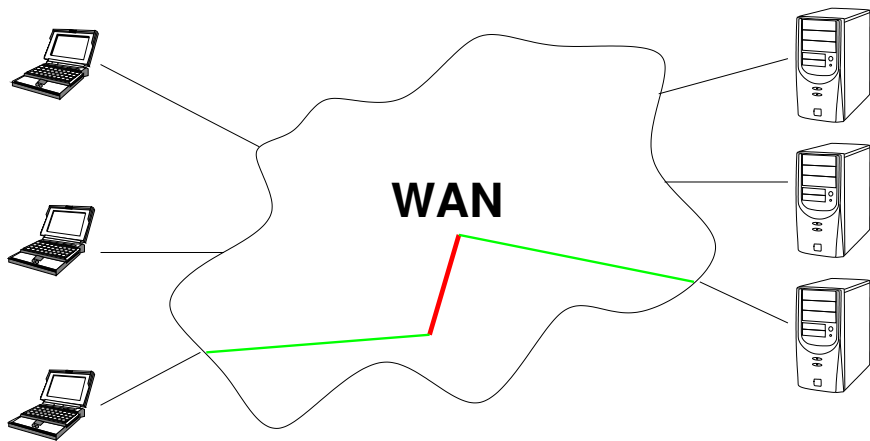
- Numerical methods
- Applications
  - MEMS simulation (system level and device level)
  - Computer network tomography
  - Social networks: community detection, opinion formation
  - Materials design
- Software and computational infrastructure

- Numerical methods
- Applications
- Software and computational infrastructure
  - Packages for individual algorithms
  - Simulators (SUGAR, HiQLab, MATFEAP, MatScat, deal2lab)
  - Infrastructure for cloud numerics (with Gehrke)

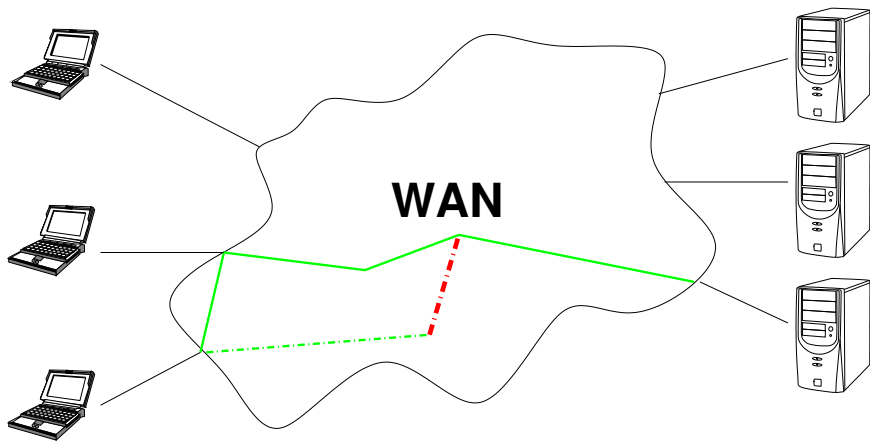
# Application: Computer Network Tomography



# A Possible Problem

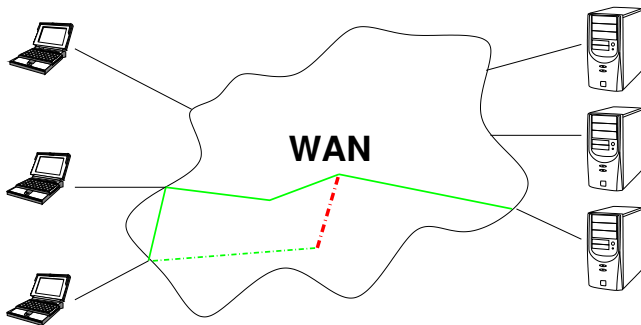


# Find and Fix or Route Around?





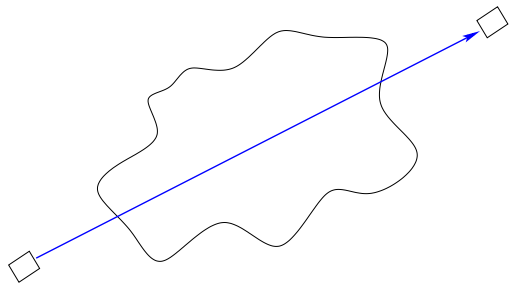
# Overlays and measurement



Measure a few paths to infer:

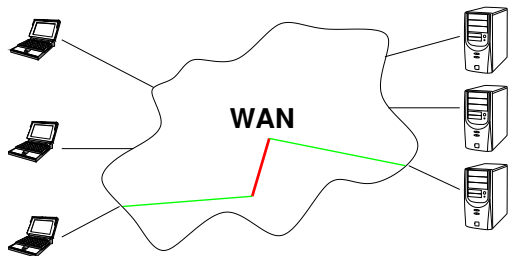
- *Path* properties (Chen, B., Song, Chavez, Katz: 2003, 2004, 2007)
- *Link* properties (Zhao, Chen, B.: 2006, 2009)
- Routing topology? (underway)

# Discrete Radon transform



Radon transform:

$$(Ru)(L) = \int_L u(\mathbf{x}) |d\mathbf{x}|$$



Discrete version:

$$(Gu)_i = \sum_{j \in \text{links}(i)} u_j$$

# Additive metrics and path matrices

For additive metrics ( $\log(P(\text{transmission}))$ , latency, jitter, ...):

$$Gu = b$$

where

- $b_i$  = property of  $i$ th end-to-end path
- $u_j$  = property of link  $j$
- $G_{ij} = \begin{cases} 1 & \text{if path } i \text{ uses link } j \\ 0 & \text{otherwise} \end{cases}$

**Goal:** Relate network structure to a structured decomposition of  $G$

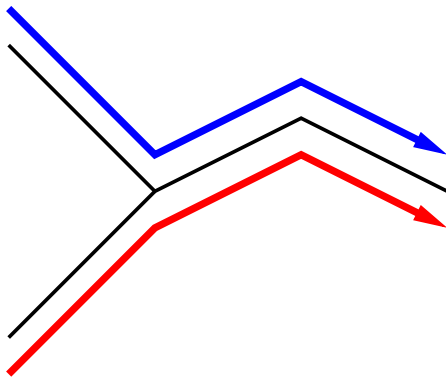
# Properties of $G$

Network	Matrix $G$
Routing path	Row of $G$
Short paths	Sparse $G$
Routing table updates	Low rank updates to $G$

$k = \text{rank}(G) < \# \text{ links} \ll \# \text{ paths}$  (for  $n$  sufficiently large).

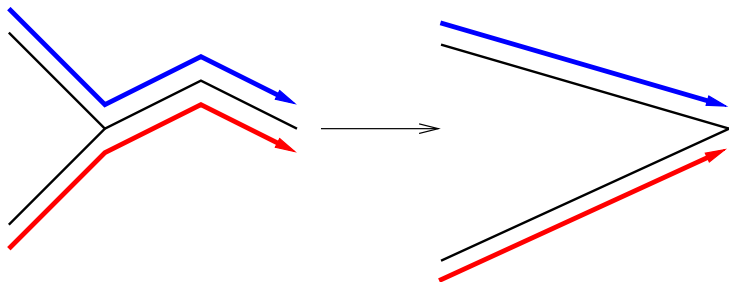
Ex: 500 nodes (of 100K),  $G$  is  $249500 \times 33237$ ,  $k = 9643$ .

# Identifiability



If both paths are flaky, what link is to blame?

# Network virtualization and matrix factorization



Factor out a zero-one “virtualization matrix”:

$$G(:, \text{fan}) = \begin{bmatrix} c_1 & c_2 & c_1 + c_2 & c_1 + c_2 \end{bmatrix} = \begin{bmatrix} c_1 & c_2 \end{bmatrix} \begin{bmatrix} 1 & 0 & 1 & 1 \\ 0 & 1 & 1 & 1 \end{bmatrix}$$

Even virtual links may be “unidentifiable.”

# Network virtualization and matrix factorization

Write network virtualization as

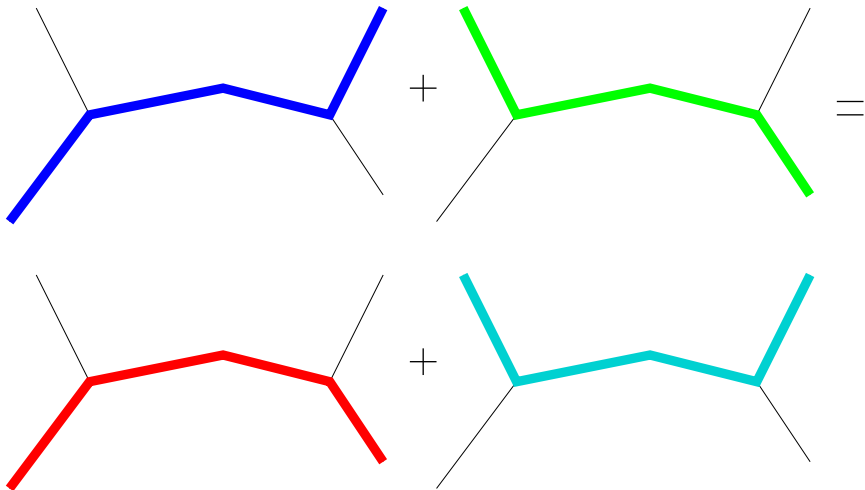
$$G = G^v S$$

where

- $G_{ij} = \begin{cases} 1 & \text{if path } i \text{ uses link } j \\ 0 & \text{otherwise} \end{cases}$
- $G_{ik}^v = \begin{cases} 1 & \text{if path } i \text{ uses virtual link } k \\ 0 & \text{otherwise} \end{cases}$
- $S_{kj} = \begin{cases} 1 & \text{if virtual link } k \text{ includes link } j \\ 0 & \text{otherwise} \end{cases}$

This handles (some) column dependencies. What about rows?

# Linear Algebra of Paths





# Matrix factorization perspective

Can combine with topology virtualization:

$$PG = \begin{bmatrix} I \\ T \end{bmatrix} \bar{G}^v S$$

where  $P$  is a permutation and  $S$  is a zero-one virtualization matrix.

Topics for offline discussion:

- How does one compute this factorization fast?
- How can you keep it up to date?
- How can you use it to infer
  - Path and link properties?
  - Missing topology?

# Other connections?

- Domain decomposition and subdomain model reduction
  - What should an aggregator send to be most useful for SE?
- Stability monitoring and novel eigenvalue problems
  - Electrical network + out-of-band control = DDAE stability?
- Latency-tolerant infrastructure
  - How to program bulk-synchronous methods on cloud (efficiently)?
  - Right tradeoffs between bulk-synchronous and asynchronous?